



# Graphs with vertex-coloring and detectable 2-edge-weighting

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## Abstract

For a connected graph  $G$  of order  $|V(G)| \geq 3$  and a  $k$ -edge-weighting  $c : E(G) \rightarrow \{1, 2, \dots, k\}$  of the edges of  $G$ , the code,  $code_c(v)$ , of a vertex  $v$  of  $G$  is the ordered  $k$ -tuple  $(\ell_1, \ell_2, \dots, \ell_k)$ , where  $\ell_i$  is the number of edges incident with  $v$  that are weighted  $i$ . (i) The  $k$ -edge-weighting  $c$  is *detectable* if every two adjacent vertices of  $G$  have distinct codes. The minimum positive integer  $k$  for which  $G$  has a detectable  $k$ -edge-weighting is the *detectable chromatic number*  $det(G)$  of  $G$ . (ii) The  $k$ -edge-weighting  $c$  is a *vertex-coloring* if every two adjacent vertices  $u, v$  of  $G$  with codes  $code_c(u) = (\ell_1, \ell_2, \dots, \ell_k)$  and  $code_c(v) = (\ell'_1, \ell'_2, \dots, \ell'_k)$  have  $1\ell_1 + 2\ell_2 + \dots + k\ell_k \neq 1\ell'_1 + 2\ell'_2 + \dots + k\ell'_k$ . The minimum positive integer  $k$  for which  $G$  has a vertex-coloring  $k$ -edge-weighting is denoted by  $\mu(G)$ . In this paper, we have enlarged the known families of graphs with  $det(G) = \mu(G) = 2$ .

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## 1. Introduction

For graph-theoretical terminology and notation, we in general follow [1]. In this paper, we assume that the graphs  $G$  in discussion are finite, connected, undirected and simple with order  $|V(G)| \geq 3$ .

Let  $c : E(G) \rightarrow \{1, 2, \dots, k\}$  be a  $k$ -edge-weighting of  $G$ , where  $k$  is a positive integer. The *color code* of a vertex  $v$  of  $G$  is the ordered  $k$ -tuple  $code_c(v) = (\ell_1, \ell_2, \dots, \ell_k)$ , where  $\ell_i$  is the number of edges incident with  $v$  that are weighted  $i$  for  $i \in \{1, 2, \dots, k\}$ . Therefore,  $\ell_1 + \ell_2 + \dots + \ell_k = d_G(v)$ , the degree of  $v$  in  $G$ . It follows that for  $u, v \in V(G)$  if  $d_G(u) \neq d_G(v)$ , then  $code_c(u) \neq code_c(v)$ . The  $k$ -edge-weighting  $c$  of  $G$  is called *detectable* if every two adjacent vertices of  $G$  have distinct color codes. The *detectable chromatic number*  $det(G)$  of  $G$  is the minimum positive integer  $k$  for which  $G$  has a detectable  $k$ -edge-weighting.

Any  $k$ -edge-weighting  $c : E(G) \rightarrow \{1, 2, \dots, k\}$  induces a vertex-weighting  $f_c : V(G) \rightarrow \mathbb{N}$  defined by  $f_c(v) = \sum_{e \text{ is incident with } v} c(e)$ . An edge-weighting  $c$  is a *vertex-coloring* if  $f_c(u) \neq f_c(v)$  for any edge  $uv$ . Denote by  $\mu(G)$  the minimum  $k$  for which  $G$  has a vertex-coloring  $k$ -edge-weighting.

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If a graph has an edge as a component, then it neither has a detectable edge-weighting nor has a vertex-coloring edge-weighting. So in this paper, we only consider graphs without a  $K_2$  component and such graphs are called *nice graphs*. As the graph  $G$  in discussion is connected and as  $|V(G)| \geq 3$ ,  $G$  is nice.

Karoński et al. [2] initiated the study of vertex-coloring  $k$ -edge-weighting and they posed the following conjecture:

**Conjecture 1.1** (1-2-3-Conjecture). *Every nice graph admits a vertex-coloring 3-edge-weighting.*

Consider a vertex-coloring  $k$ -edge-weighting  $c$  of  $G$ . For  $uv \in E(G)$ , let  $\ell_i, \ell'_i$ , respectively, be the number of edges incident with  $u, v$  that are weighted  $i$  in  $c$ . Then  $1\ell_1 + 2\ell_2 + \dots + k\ell_k \neq 1\ell'_1 + 2\ell'_2 + \dots + k\ell'_k$  and hence  $(\ell_1, \ell_2, \dots, \ell_k) \neq (\ell'_1, \ell'_2, \dots, \ell'_k)$ . So  $c$  is a detectable  $k$ -edge-weighting. Consequently,  $\det(G) \leq \mu(G)$ .

**Proposition 1.1.**  $\det(G) \leq \mu(G)$ .

**Proposition 1.2.** *For every nice graph  $G$ , following three conditions are equivalent:*

- (i)  $\det(G) = 1$ ,
- (ii)  $\mu(G) = 1$ ,
- (iii)  $G$  has no adjacent vertices with the same degree.

**Proposition 1.3.** *If  $\mu(G) = 2$ , then  $\det(G) = 2$ .*

If  $c$  is a detectable 2-edge-weighting of a  $k$ -regular graph  $G$  with  $k \geq 3$ , then  $c$  is a vertex-coloring 2-edge-weighting. This follows from the fact that  $\ell_1 + \ell_2 = k = \ell'_1 + \ell'_2$  and  $(\ell_1, \ell_2) \neq (\ell'_1, \ell'_2)$  imply  $1\ell_1 + 2\ell_2 \neq 1\ell'_1 + 2\ell'_2$ .

**Proposition 1.4.** *Let  $G$  be a  $k$ -regular graph with  $k \geq 3$ . If  $\det(G) = 2$ , then  $\mu(G) = 2$ .*

In [2], Karoński et al. proved that: (i)  $\det(G) \leq 183$ , and (ii) if  $d_G(v) \geq 10^{99}$  for every  $v \in V(G)$ , then  $\det(G) \leq 30$ .

In [3], Addario-Berry et al. proved that: (i)  $\det(G) \leq 4$ , (ii) if  $d_G(v) \geq 1000$  for every  $v \in V(G)$ , then  $\det(G) \leq 3$ , and (iii) if  $\chi(G) \leq 3$ , then  $\det(G) \leq 3$ .

In [4], among other results, Escudro et al. proved that: (i)  $\det(K_{n_1, n_2, \dots, n_k}) = 1$  if  $n_1 < n_2 < \dots < n_k$ ,  $\det(K_{n_1, n_2, \dots, n_k}) = 3$  if  $n_1 = n_2 = \dots = n_k = 1$  and  $\det(K_{n_1, n_2, \dots, n_k}) = 2$  otherwise, where  $K_{n_1, n_2, \dots, n_k}$  is the complete  $k$ -partite graph with partite sizes  $n_1, n_2, \dots, n_k$  ( $k \geq 3$  and  $n_1 \leq n_2 \leq \dots \leq n_k$ ), (ii)  $\det(C_3 \square K_2) = 3$ ,  $\det(C_5 \square K_2) = 3$  and if  $n \geq 7$  is an odd integer, then  $\det(C_n \square K_2) = 2$ , where  $\square$  denotes the Cartesian product, and (iii) if  $G$  is a unicyclic graph that is not a cycle, then  $\det(G) \leq 2$ .

See Fig. 5 of [4]; detectable 3-edge-weighting of  $C_3 \square K_2$  and that of  $C_5 \square K_2$ , in the figure, are vertex-coloring 3-edge-weightings. Hence,  $\mu(C_3 \square K_2) = 3$  and  $\mu(C_5 \square K_2) = 3$ . If  $n \geq 7$  is an odd integer, then it follows from  $\det(C_n \square K_2) = 2$  and Proposition 1.4 that  $\mu(C_n \square K_2) = 2$ .

**Theorem 1.1.**  $\det(C_3 \square K_2) = \mu(C_3 \square K_2) = 3$ ,  $\det(C_5 \square K_2) = \mu(C_5 \square K_2) = 3$  and if  $n \geq 7$  is an odd integer, then  $\det(C_n \square K_2) = \mu(C_n \square K_2) = 2$ .

From [5,6], and [4], we have:

**Theorem 1.2.** *For the path  $P_n$  on  $n$  vertices,  $\det(P_3) = \mu(P_3) = 1$  and  $\det(P_n) = \mu(P_n) = 2$  if  $n \geq 4$ .*

**Theorem 1.3.** *For the cycle  $C_n$  on  $n$  vertices,  $\det(C_n) = \mu(C_n) = 2$  if  $n \equiv 0 \pmod{4}$  and  $\det(C_n) = \mu(C_n) = 3$  if  $n \equiv 1, 2$  or  $3 \pmod{4}$ .*

**Theorem 1.4.** *For the complete graph  $K_n$  on  $n \geq 3$  vertices,  $\det(K_n) = \mu(K_n) = 3$ .*

**Theorem 1.5.** *For  $r + s \geq 3$ ,  $\det(K_{r,s}) = \mu(K_{r,s}) = 1$  if  $r \neq s$  and  $\det(K_{r,s}) = \mu(K_{r,s}) = 2$  if  $r = s$ , where  $K_{r,s}$  is the complete bipartite graph with partite sizes  $r$  and  $s$ .*

The *theta graph*  $\theta(\ell_1, \ell_2, \dots, \ell_r)$  is the graph obtained from  $r$  disjoint paths  $P_1(u_1, v_1), P_2(u_2, v_2), \dots, P_r(u_r, v_r)$  of lengths  $\ell_1, \ell_2, \dots, \ell_r$ , respectively, by identifying their end-vertices  $u := u_1 = u_2 = \dots = u_r$  and  $v := v_1 = v_2 = \dots = v_r$ , where  $P_i(u_i, v_i)$  is a path of length  $\ell_i$  with origin  $u_i$  and terminus  $v_i$ . Note that  $\theta(\ell_1) = P_{\ell_1+1}$  and  $\theta(\ell_1, \ell_2) = C_{\ell_1+\ell_2}$ .

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