



Every 3-polytope with minimum degree 5 has a 6-cycle with maximum degree at most 11



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ABSTRACT

Let $\varphi_P(C_6)$ (respectively, $\varphi_T(C_6)$) be the minimum integer k with the property that every 3-polytope (respectively, every plane triangulation) with minimum degree 5 has a 6-cycle with all vertices of degree at most k . In 1999, S. Jendrol' and T. Madaras proved that $10 \leq \varphi_T(C_6) \leq 11$. It is also known, due to B. Mohar, R. Škrekovski and H.-J. Voss (2003), that $\varphi_P(C_6) \leq 107$.

We prove that $\varphi_P(C_6) = \varphi_T(C_6) = 11$.

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1. Introduction

The *degree* $d(x)$ of a vertex or face x in a plane graph G is the number of incident edges. A k -*vertex* (k -*face*) is a vertex (face) with degree k , a k^+ -*vertex* has degree at least k , etc. The minimum vertex degree of G is $\delta(G)$. We will drop the arguments whenever this does not lead to confusion.

A normal plane map (NPM) is a plane pseudograph in which loops and multiple edges are allowed, but $d(x) \geq 3$ for every vertex or face x . As proved by Steinitz [20], the 3-connected plane graphs are planar representations of the convex three-dimensional polytopes, called hereafter 3-polytopes.

In this note, we consider the class \mathbf{M}_5 of NPMs with $\delta = 5$ and its subclasses \mathbf{P}_5 of 3-polytopes and \mathbf{T}_5 of plane triangulations. A cycle on k vertices is denoted by C_k , and S_k stands for a k -star centered at a 5-vertex. (So, S_k is a subgraph of M_5 on a 5-vertex and k vertices adjacent to it, where $0 \leq k \leq 5$.)

In 1904, Wernicke [21] proved that $M_5 \in \mathbf{M}_5$ implies, in M_5 , the presence of a vertex of degree 5 adjacent to a vertex of degree at most 6. This result was strengthened by Franklin [8] in 1922 to the existence of a vertex of degree 5 with two neighbors of degree at most 6. In 1940, Lebesgue [15, p. 36] gave an approximate description of the neighborhoods of vertices of degree 5 in \mathbf{T}_5 .

The *weight* $w_M(H)$ is the maximum over $M_5 \in \mathbf{M}_5$ of the minimum degree-sum of the vertices of H over subgraphs H of M_5 . The *weights* $w_P(H)$ and $w_T(H)$ are defined similarly for \mathbf{P}_5 and \mathbf{T}_5 , respectively.

The bounds $w_M(S_1) \leq 11$ (Wernicke [21]) and $w_M(S_2) \leq 17$ (Franklin [8]) are tight. It was proved by Lebesgue [15] that $w_M(S_3) \leq 24$ and $w_M(S_4) \leq 31$, which was improved much later to the following tight bounds: $w_M(S_3) \leq 23$ (Jendrol')

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and Madaras [10]) and $w_M(S_4) \leq 30$ (Borodin and Woodall [6]). Note that $w_M(S_3) \leq 23$ readily implies $w_M(S_2) \leq 17$ and immediately follows from $w_M(S_4) \leq 30$ (it suffices to delete a vertex of maximum degree from a star of the minimum weight).

It follows from Lebesgue [15, p. 36] that $w_T(C_3) \leq 18$. In 1963, Kotzig [14] gave another proof of this fact and conjectured that $w_T(C_3) \leq 17$. (The bound 17 is easily shown to be tight.)

In 1989, Kotzig's conjecture was confirmed by Borodin [2] in a more general form, by proving $w_M(C_3) = 17$. Another consequence of this result is the confirming of a conjecture of Grünbaum [9] from 1975 that the cyclic connectivity (defined as the minimum number of edges to be deleted from a graph to obtain two components each containing a cycle) of every 5-connected planar graph is at most 11, which is tight (a bound of 13 was earlier obtained by Plummer [19]).

It also follows from Lebesgue [15, p. 36] that $w_T(C_4) \leq 26$ and $w_T(C_5) \leq 31$. In 1998, Borodin and Woodall [6] proved $w_T(C_4) = 25$ and $w_T(C_5) = 30$.

Now let $\varphi_M(H)$ ($\varphi_P(H)$, $\varphi_T(H)$) be the minimum integer k with the property that every normal plane map (3-polytope, plane triangulation) with minimum degree 5 has a copy of H with all vertices of degree at most k .

It follows from Franklin [8] that $\varphi_M(S_2) = 6$. From $w_M(C_3) = 17$ (Borodin [2]), we have $\varphi_M(C_3) = 7$. In 1996, Jendrol' and Madaras [10] proved $\varphi_M(S_4) = 10$ and $\varphi_T(C_4) = \varphi(C_5) = 10$. R. Soták (personal communication; see surveys of Jendrol' and Voss [12,13]) proved $\varphi_P(C_4) = 11$ and $\varphi_P(C_5) = 10$.

In 1999, Jendrol' et al. [11] obtained the following bounds: $10 \leq \varphi_T(C_6) \leq 11$, $15 \leq \varphi_T(C_7) \leq 17$, $15 \leq \varphi_T(C_8) \leq 29$, $19 \leq \varphi_T(C_9) \leq 41$, and $\varphi_T(C_p) = \infty$ whenever $p \geq 11$. Madaras and Soták [17] proved $20 \leq \varphi_T(C_{10}) \leq 415$.

For the broader class \mathbf{P}_5 (an easy induction proof shows that every planar triangulation on at least four vertices is 3-connected), it is known that $10 \leq \varphi_P(C_6) \leq 107$ due to Mohar et al. [18] (in fact, this bound is proved in [18] for all 3-polytopes with $\delta \geq 4$ in which no 4-vertex is adjacent to a 4-vertex), and $\varphi_P(C_7) \leq 359$ is due to Madaras et al. [16].

The purpose of our note is to prove that $\varphi_P(C_6) = \varphi_T(C_6) = 11$. This answers a question raised by Jendrol' et al. [11].

Theorem 1. *Every 3-polytope with minimum degree 5 has a 6-cycle such that each of its vertices has degree at most 11, and this bound is tight.*

Other structural results on \mathbf{M}_5 , some of which have application to coloring, can be found in the papers already mentioned and in [3,4,7,16–18].

One of the ideas used in our proof is to look for a suitable 6-cycle not in the whole graph but in a carefully chosen portion of it. A similar approach to coloring problems on plane graphs is described in a survey [5, pp. 520–521], and it has been used by us several times, beginning with [1].

2. Proving the tightness of Theorem 1

We transform the octahedron (the 4-regular plane triangulation on six vertices) to a plane triangulation in which every 6-cycle goes through a vertex of degree at least 11, replacing each of the eight 3-faces of the octahedron by the configuration shown in Fig. 1.

More specifically, half of the image of every edge (partly invisible) of the octahedron starts at an “angular” 12-vertex, goes through an 11-vertex, cuts an edge between two 5-vertices, then goes through two 5-vertices, cuts another edge between two 5-vertices, and ends in a 12-vertex, the mid-point of the image of the edge. The graph obtained has only 5-, 11-, and 12-vertices. Furthermore, every 5-vertex belongs to a blue (shadowed) triangle. It is easily seen that the subgraph on 5-vertices does not contain 6-cycles.

Note that we could use instead of the octahedron any plane triangulation with $\delta \geq 4$ to obtain a plane triangulation with the desired property.

3. Proving the upper bound in Theorem 1

Suppose G' is a counterexample to the main statement of Theorem 1. Thus G' is a 3-polytope with $\delta = 5$ in which no 6-cycle avoids a 12⁺-vertex.

By Euler's formula $|V'| - |E'| + |F'| = 2$ for G' , we have

$$\sum_{v \in V'} (d(v) - 4) + \sum_{f \in F'} (d(f) - 4) = -8. \quad (1)$$

This implies that G' has a 3-face. So we may assume that the external face of G' is bounded by a 3-cycle with the vertex set T' .

A special triangle $T^* = t_1 t_2 t_3$ of G' is a 3-cycle of G' with the fewest vertices inside. We define G to be the subgraph of G' induced by the vertices inside T^* . The vertices of G are *internal*, and the vertices t_1 , t_2 , and t_3 are *special*.

By G^{**} we denote the subgraph of G' induced by the vertices of $G \cup T^*$. In particular, $T^* = T'$ when $G^{**} = G'$. In both cases, T^* is the boundary $\partial(f_\infty)$ of the external face f_∞ of G^{**} .

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