



All roots spectral methods: Constraints, floating point arithmetic and root exclusion



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ARTICLE INFO

Article history:

Received 14 October 2016

Received in revised form 29

November 2016

Accepted 30 November 2016

Available online 7 December 2016

Keywords:

Chebyshev polynomials

Nonlinear ordinary differential equations

Two-point boundary value problem

Lemniscate elliptic function

Computer algebra

ABSTRACT

The nonlinear two-point boundary value problem (TPBVP for short)

$$u_{xx} + u^3 = 0, u(0) = u(1) = 0,$$

offers several insights into spectral methods. First, it has been proved *a priori* that $\int_0^1 u(x) dx = \pi/\sqrt{2}$. By building this constraint into the spectral approximation, the accuracy of $N + 1$ degrees of freedom is achieved from the work of solving a system with only N degrees of freedom. When N is small, generic polynomial system solvers, such as those in the computer algebra system Maple, can find all roots of the polynomial system, such as a spectral discretization of the TPBVP. Our second point is that floating point arithmetic in lieu of exact arithmetic can double the largest practical value of N . (Rational numbers with a huge number of digits are avoided, and eliminating M symbols like $\sqrt{2}$ and π reduces $N + M$ -variate polynomials to polynomials in just the N unknowns.) Third, a disadvantage of an “all roots” approach is that the polynomial solver generates many roots – $(3^N - 1)$ for our example – which are genuine solutions to the N -term discretization but spurious in the sense that they are not close to the spectral coefficients of a true solution to the TPBVP. We show here that a good tool for “root-exclusion” is calculating $\rho \equiv \sqrt{\sum_{n=1}^N b_n^2}$; spurious roots have ρ larger than that for the physical solution by at least an order of magnitude. The ρ -criterion is suggestive rather than infallible, but root exclusion is very hard, and the best approach is to apply multiple tools with complementary failings.

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1. Introduction

The two-point boundary-value problem

$$u_{xx} + u^3 = 0, \quad u(0) = u(1) = 0 \quad (1)$$

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has the exact solution

$$u = 2\mathcal{K}\text{cn}(2\mathcal{K}(x - 1/2), 1/\sqrt{2}), \quad \mathcal{K} = 1.85407 \quad (2)$$

where “cn” is the usual Jacobian elliptic function and $\mathcal{K} = K(1/\sqrt{2})$ is the complete elliptic integral for an elliptic modulus $k = 1/\sqrt{2}$. Boyd showed that the problem is “cryptoperiodic” in the sense the most efficient spectral basis is a Fourier basis [1] because the elliptic function is periodic. Here, we apply a polynomial basis partly for comparison and partly and more importantly because polynomial approximations are more widely used than the Fourier basis.

Although difficult problems require N very large where N is the number of degrees of freedom, spectral methods are so powerful that they are often accurate even for very small N . This makes it possible to bypass Newton’s iteration and its ilk by applying a polynomial system solver to find all its zeros. This system of N equations in N unknowns is the spectral discretization of the TPBVP and the unknowns are the coefficients of the Chebyshev polynomial approximation [1,2].

2. Imposing an integral constraint

N can be minimized by incorporating symmetry and boundary conditions into the form of the approximation. Here $u(x)$ is symmetric with respect to reflection about the midpoint of the interval [1]; this symmetry can be imposed by including only basis functions with the same symmetry, that is, only even degree (shifted) Chebyshev polynomials, yielding

$$u_N(x) = x(1-x) \left\{ b_0 + \sum_{n=1}^N b_n T_{2n}(2x-1) \right\}. \quad (3)$$

Independent of the coefficients b_n , $u_N(0) = u_N(1) = 0$ and $u_N(1-x) = u_N(x)$.

Gheorghiu and Trif [3] observed that

$$\int_0^1 u(x) dx = \pi/\sqrt{2}. \quad (4)$$

This can be used to express b_0 in terms of the other b_n . We assume that the two boundary conditions will be incorporated into the algorithm also. Thus, although the approximation contains only N degrees of freedom, it is a polynomial of degree $(N+2)$ and thus a polynomial with $(N+3)$ coefficients.

Table 1 shows that for this smooth problem, albeit a cubically nonlinear problem, the four degree of freedom approximation has a maximum pointwise error which is only 1/14,600 of the maximum of $u(x)$. The complete Maple listing is given in Table 2.

The theory and practice of low degree polynomial spectral methods are described in Chapter 20 of [4], Chapter 2 of [5], Chapters 1 through 6 of [6], the review [7] and [1,7,8].

3. The downside of exact arithmetic

In our Maple 15 system running in Windows 7, the largest feasible number of degrees of freedom was limited by system crashes to $N_{feasible} = 2$ when exact arithmetic was used as far as possible in the computation. $N_{feasible}$ was *doubled* by altering a single line:

$$r[m]- > \text{evalf}(r[m]) \quad (5)$$

where $r[m]$ is the m th equation of Galerkin system and where “evalf” is the Maple command to force replacement of symbols like $\sqrt{2}$, π and $BesselJ(0)$ by their floating equivalents. Why does “evalf” make such a difference?

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