Contents lists available at ScienceDirect

## **Discrete Mathematics**

journal homepage: www.elsevier.com/locate/disc

## On the diameter of Kronecker graphs

### Justyna Banaszak, Tomasz Łuczak\*

Faculty of Mathematics and Computer Science, Adam Mickiewicz University, Poznań, Poland

#### ARTICLE INFO

#### ABSTRACT

Article history: Received 4 May 2017 Received in revised form 5 June 2018 Accepted 26 July 2018 We prove that a.a.s. as soon as a Kronecker graph becomes connected it has a finite diameter. © 2018 Elsevier B.V. All rights reserved.

*Keywords:* Random graphs Kronecker graph Diameter Connectivity

#### 1. Introduction

A Kronecker graph is a random graph with vertex set  $V = \{0, 1\}^n$ , where the probability that two vertices  $u, v \in V$  are adjacent strongly depends on the structure of the vectors  $u = (u_1, ..., u_n)$ , and  $v = (v_1, ..., v_n)$ . More specifically, let **P** be a symmetric matrix

$$\mathbf{P} = \frac{1}{0} \begin{pmatrix} 1 & 0 \\ \alpha & \beta \\ \beta & \gamma \end{pmatrix},$$

where zeros and ones are labels of rows and columns of **P**,  $\alpha$ ,  $\beta$ ,  $\gamma \in [0, 1]$ , and  $\alpha \geq \gamma$ . In the Kronecker graph  $\mathcal{K}(n, \mathbf{P})$  two vertices  $u = (u_1, \ldots, u_n), v = (v_1, \ldots, v_n) \in V = \{0, 1\}^n$  are adjacent with probability

$$p_{u,v} = \prod_{i=1}^{n} \mathbf{P}[u_i, v_i]$$

independently for each such pair.

Kronecker graphs were introduced by Leskovec, Chakrabarti, Kleinberg and Faloutsos in [2] to model some real world networks (see also [1,3,7]). Since then they have been studied by several authors but their properties are still far from being well understood (see [4] and references therein). In particular, Radcliffe and Young [9] determined the exact threshold for the property that  $\mathcal{K}(n, \mathbf{P})$  is connected, supplementing a slightly weaker result of Mahdian and Xu [8].

\* Corresponding author.

https://doi.org/10.1016/j.disc.2018.07.026 0012-365X/© 2018 Elsevier B.V. All rights reserved.





E-mail addresses: tabor@amu.edu.pl (J. Banaszak), tomasz@amu.edu.pl (T. Łuczak).

#### Theorem 1.

$$\lim_{n \to \infty} \mathbb{P}(\mathcal{K}(n, \mathbf{P}) \text{ is connected}) = \begin{cases} 0 & \text{if } \beta + \gamma = 1, \ \beta \neq 1 \\ 0 & \text{if } \beta = 1, \ \alpha = \gamma = 0 \\ 1 & \text{if } \beta = 1, \ \alpha > 0 \text{ and } \gamma = 0 \\ 1 & \text{if } \beta + \gamma > 1. \end{cases}$$

The main result of this work states that as soon as  $\mathcal{K}(n, \mathbf{P})$  becomes connected its diameter is bounded by a constant.

**Theorem 2.** If either  $\beta + \gamma > 1$ , or  $\beta = 1$ ,  $\alpha > 0$  and  $\gamma = 0$ , then there exists a constant  $a = a(\alpha, \beta, \gamma)$  such that a.a.s.  $diam(\mathcal{K}(n, \mathbf{P})) \leq a$ .

#### 2. The idea of the proof

In order to sketch our argument let us recall how one shows that the diameter is bounded from above for the binomial random graph model G(N, p), and for many other random graph models. Typically, since random graphs are good expanders, it is proven first that for some small k the k-neighbourhood of each vertex is much larger than  $\sqrt{N}$ . Then, in the second part of the proof, one argues that since two random subsets of vertices of size larger than  $\sqrt{N}$  intersect with large probability, each pair of vertices is a.a.s. connected by a path of length at most 2k. However, in our case this procedure fails completely. The main reason is that most neighbourhood of a given vertex v have a similar structure, and so the events ' $x \sim v$ ' and ' $y \sim v$ ' are strongly correlated. Thus, the k-neighbourhood of a given vertex is very far from being a random subset, which is crucial for the second step of the procedure. Even more importantly, we do not understand expanding properties of  $\mathcal{K}(n, \mathbf{P})$  and it is hard to control how fast the k-neighbourhoods of a vertex  $\mathcal{K}(n, \mathbf{P})$  grow, which in most of the other random graph models is quite easy to investigate.

In [8], the diameter of  $\mathcal{K}(n, \mathbf{P})$  is studied for  $\gamma \leq \beta \leq \alpha$ . For this specific range of parameters the probability of appearance of an edge of  $\mathcal{K}(n, \mathbf{P})$  grows with the weights of its ends, i.e. for every two vertices u, v the probability that there exists an edge uv is always greater than the probability of an edge uv', whenever v has greater weight than v'. Using this fact the authors of [8] bounded from above the diameter of  $\mathcal{K}(n, \mathbf{P})$  using well known bounds for the diameter of binomial random graphs.

To handle the difficulties related to the dependence of edges in  $\mathcal{K}(n, \mathbf{P})$  in the general case we use the following approach. We consider two vertices, v and u which are 'similar' to each other (more specifically, we choose both of them from the middle layer of the *n*-cube and assume that they have small Hamming distance from each other). Then we generate their neighbourhoods at the same time until, for some k, we observe that the k-neighbourhood of v does not expand according to its expected rate. This is because many, in fact most, candidates for (k + 1)-neighbours of v have already been placed in the *i*-neighbourhood of v for some  $i \leq k$ . However, the chance that a vertex x is in the *i*-neighbourhood of v is roughly the same as the probability that x is in the *i*-neighbourhood of u so, if most potential (k + 1)-neighbours of v are already in its kth neighbourhood, many of them are also in the kth neighbourhood of u. Consequently, there is a path of length at most 2k joining v and u.

The structure of the paper is the following. First we treat a special 'pathological' case  $\beta = 1$ . Then we present the crucial part of our argument showing that the subgraph induced in  $\mathcal{K}(n, \mathbf{P})$  by its middle layer has a.a.s. a small diameter. Finally, we complete the proof showing that a.a.s. each vertex of  $\mathcal{K}(n, \mathbf{P})$  is connected to the middle layer by a short path.

#### 3. Case $\beta = 1$

In this section we show that if  $\beta = 1$ ,  $\alpha > 0$ , and  $\gamma = 0$ , then the diameter of  $\mathcal{K}(n, \mathbf{P})$  is a.a.s. bounded by a constant. This set of parameters  $\alpha$ ,  $\beta$ ,  $\gamma$  is somewhat special as it is the only case, when  $\beta + \gamma = 1$  and still  $\mathcal{K}(n, \mathbf{P})$  is a.a.s. connected.

We introduce some notation, which we shall use throughout the paper. By d(v, u) we denote the Hamming distance between two vertices v and u and w(v) stands for the weight of a vertex  $v = (v_1, \ldots, v_n)$ , i.e. the number of ones in its label that is

$$w(v) = \sum_{i=1}^n v_i.$$

For a vertex  $v = (v_1, ..., v_n)$ , we set  $\bar{v} = (1 - v_1, ..., 1 - v_n)$ .

Now let us go back to the case  $\beta = 1$ . Note that

$$\mathbb{P}(v \sim \bar{v}) = \beta^n = 1,$$

and observe that either v or  $\bar{v}$  has weight at least n/2. Thus, to show the assertion it is enough to verify that a.a.s. there exists a path of bounded length between every pair of vertices in R defined as

$$R = \{v \in V : w(v) \ge n/2\}.$$

Download English Version:

# https://daneshyari.com/en/article/8902839

Download Persian Version:

https://daneshyari.com/article/8902839

Daneshyari.com