



ELSEVIER

Contents lists available at ScienceDirect

Journal of Combinatorial Theory,  
Series B

www.elsevier.com/locate/jctb

Bipartitions of oriented graphs <sup>☆</sup>Jianfeng Hou <sup>a</sup>, Shufei Wu <sup>a,b</sup><sup>a</sup> Center for Discrete Mathematics, Fuzhou University, Fujian, 350116, China<sup>b</sup> School of Mathematics and Information Science, Henan Polytechnic University, Henan, 454000, China

## ARTICLE INFO

*Article history:*

Received 12 July 2016

Available online xxxx

*Keywords:*

Oriented graph

Partition

Semidegree

Tight component

## ABSTRACT

Let  $V(D) = X \cup Y$  be a bipartition of a directed graph  $D$ . We use  $e(X, Y)$  to denote the number of arcs in  $D$  from  $X$  to  $Y$ . Motivated by a conjecture posed by Lee, Loh and Sudakov (2016) [16], we study bipartitions of oriented graphs. Let  $D$  be an oriented graph with  $m$  arcs. In this paper, it is proved that if the minimum degree of  $D$  is  $\delta$ , then  $D$  admits a bipartition  $V(D) = V_1 \cup V_2$  such that  $\min\{e(V_1, V_2), e(V_2, V_1)\} \geq (\frac{\delta-1}{4\delta} + o(1))m$ . Moreover, if the minimum semidegree  $d = \min\{\delta^+(D), \delta^-(D)\}$  of  $D$  is at least 21, then  $D$  admits a bipartition  $V(D) = V_1 \cup V_2$  such that  $\min\{e(V_1, V_2), e(V_2, V_1)\} \geq (\frac{d}{2(2d+1)} + o(1))m$ . Both bounds are asymptotically best possible.

© 2018 Elsevier Inc. All rights reserved.

## 1. Introduction

An example of classical graph partitioning problems is the well-known *Max-Cut Problem*: Given a graph  $G$ , find the maximum bipartite subgraph of  $G$ . A simple average argument shows that every graph with  $m$  edges has a bipartite subgraph with at least  $m/2$  edges. In 1973, answering a question of Erdős, Edwards [5,6] improved this lower

<sup>☆</sup> This work is supported by the National Natural Science Foundation of China (No. 11671087).

*E-mail address:* shufeiwu@hotmail.com (S. Wu).

<https://doi.org/10.1016/j.jctb.2018.03.003>

0095-8956/© 2018 Elsevier Inc. All rights reserved.

bound to  $m/2 + (\sqrt{2m+1}/4 - 1/2)/4$ , which is essentially best possible as evidenced by complete graphs with odd orders. For more references concerning partitions of graphs or hypergraphs, see [3,4,7,8,11,13,18–24].

Partitioning problems of directed graphs (digraphs for short) tend to be more difficult than that of undirected graphs. In contrast to undirected graphs, where Max-Cut only needs to optimize a single parameter (the total number of crossing edges), the analogous problem in digraphs is to measure the size of the directed cut in each direction. For a digraph  $D$  and  $S, T \subseteq V(D)$ , we use  $e(S, T)$  to denote the number of arcs of  $D$  directed from  $S$  to  $T$ . In [21], Scott posed the following natural problem: What is the maximum constant  $c_d$  such that every digraph  $D$  with  $m$  arcs and minimum outdegree  $d$  admits a bipartition  $V(D) = V_1 \cup V_2$  such that

$$\min\{e(V_1, V_2), e(V_2, V_1)\} \geq c_d m?$$

For  $d = 1$ , consider the graph  $K_{1, n-1}$  and add a single edge inside the part of size  $n - 1$ . This graph can be oriented easily so that the minimum outdegree is 1 and  $\min\{e(V_1, V_2), e(V_2, V_1)\} \leq 1$  for every partition  $V(D) = V_1 \cup V_2$ . Hence,  $c_1 = 0$ . Lee, Loh and Sudakov [16] initiated the study of this problem and conjectured that

**Conjecture 1.1** (Lee, Loh and Sudakov [16]). *Let  $d$  be an integer satisfying  $d \geq 2$ . Every digraph  $D$  with  $m$  arcs and minimum outdegree at least  $d$  admits a bipartition  $V(D) = V_1 \cup V_2$  such that*

$$\min\{e(V_1, V_2), e(V_2, V_1)\} \geq \left(\frac{d-1}{2(2d-1)} + o(1)\right)m.$$

In the same paper, Lee, Loh and Sudakov verified Conjecture 1.1 for  $d = 2, 3$ , and noted that their method is not adequate for  $d \geq 4$ .

Even for oriented graphs, Conjecture 1.1 seems difficult to prove. It is not clear to us whether or not the minimum outdegree condition alone is sufficient for Conjecture 1.1 with large  $d$ . However, if we relax it to minimum degree or semidegree for oriented graphs, then we have the following results.

**Theorem 1.2.** *Every oriented graph  $D$  with  $m$  arcs and minimum degree  $\delta \geq 1$  admits a bipartition  $V(D) = V_1 \cup V_2$  such that*

$$\min\{e(V_1, V_2), e(V_2, V_1)\} \geq \left(\frac{\delta-1}{4\delta} + o(1)\right)m.$$

**Theorem 1.3.** *Let  $d$  be an integer with  $d \geq 21$ . Every oriented graph  $D$  with  $m$  arcs and minimum semidegree  $d$  admits a bipartition  $V(D) = V_1 \cup V_2$  such that*

$$\min\{e(V_1, V_2), e(V_2, V_1)\} \geq \left(\frac{d}{2(2d+1)} + o(1)\right)m.$$

Download English Version:

<https://daneshyari.com/en/article/8903837>

Download Persian Version:

<https://daneshyari.com/article/8903837>

[Daneshyari.com](https://daneshyari.com)