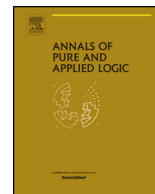




Contents lists available at ScienceDirect

Annals of Pure and Applied Logic

www.elsevier.com/locate/apal



## Atomic polymorphism and the existence property

Gilda Ferreira <sup>a,b,\*</sup>,<sup>1</sup><sup>a</sup> Departamento de Matemática, Faculdade de Ciências da Universidade de Lisboa, Campo Grande, Ed. C6, 1749-016, Lisboa, Portugal<sup>b</sup> Universidade Aberta, Rua Braamcamp, 90, 1250-052, Lisboa, Portugal

## ARTICLE INFO

## Article history:

Available online xxxx

## MSC:

03F07

03F03

03B20

03F05

## Keywords:

Predicative polymorphism  
 Intuitionistic predicate calculus  
 Existence property  
 Natural deduction  
 Normalization  
 Faithfulness

## ABSTRACT

We present a purely proof-theoretic proof of the existence property for the full intuitionistic first-order predicate calculus, via natural deduction, in which commuting conversions are not needed. Such proof illustrates the potential of an atomic polymorphic system with only three generators of formulas – conditional and first and second-order universal quantifiers – as a tool for proof-theoretical studies.

© 2018 Elsevier B.V. All rights reserved.

## 1. Introduction

In [5] it was shown that there is a purely proof-theoretic proof of the disjunction property for the full intuitionistic propositional calculus (**IPC**), via natural deduction without the need of commuting conversions. Such a proof relies in a sound and faithful embedding of **IPC** into an atomic polymorphic system with only two generators of formulas – conditional and second-order universal quantifier – denoted by  $\mathbf{F}_{\text{at}}$ .

Is it possible to give an alternative proof of the existence property for the full intuitionistic first-order predicate calculus (**IQC**) in the above *commuting conversion-free* manner? We show in this paper that the answer is yes. The atomic polymorphic system used to interpret **IQC**, which we denote by  $\mathbf{F}_{\text{at}}^{\mathbf{Q}}$ , was

\* Correspondence to: Departamento de Matemática, Faculdade de Ciências da Universidade de Lisboa, Campo Grande, Ed. C6, 1749-016, Lisboa, Portugal.

E-mail address: [gferreira@fc.ul.pt](mailto:gferreira@fc.ul.pt).

<sup>1</sup> The author acknowledges the support of Fundação para a Ciência e a Tecnologia [UID/MAT/04561/2013, UID/CEC/00408/2013 and grant SFRH/BPD/93278/2013] and is also grateful to Centro de Matemática, Aplicações Fundamentais e Investigação Operacional and Large-Scale Informatics Systems Laboratory.

introduced in [3] under the designation of atomic QSOL<sup>i</sup>. Such system, which we review in the section below, has three generators of formulas, the conditional and the first and second-order universal quantifiers.

Using atomic polymorphic systems to avoid complications of the natural deduction calculus is not a novelty. The defects of the latter calculus were eloquently exposed in [8]. According to Jean-Yves Girard *et al.* (page 74) ‘The elimination rules [for ⊥, ∨ and ∃] are very bad.’ and proceed ‘Indeed, we cannot decently work with the full fragment without identifying *a priori* different deduction... the need to add these supplementary rules [referring to the commuting conversions for ⊥, ∨ and ∃] reveals an inadequacy of the syntax.’ Fernando Ferreira coined a simple and elegant way of avoiding the ‘inadequacy’ of the natural deduction calculus in the context of **IPC** presenting in [2] an embedding of **IPC** into **F<sub>at</sub><sup>2</sup>** – a predicative system without *bad* connectives. Several studies in **IPC** taking advantage of its embedding into **F<sub>at</sub>** followed. See for example [4–7]. F. Ferreira and the author showed in [3] that the previous embedding could be extended to the full intuitionistic first-order predicate calculus. More precisely, it is possible to embed **IQC** into an atomic polymorphic calculus (**F<sub>at</sub><sup>Q</sup>**) with implication and first and second-order universal quantifiers, i.e., still without *bad* connectives. In the present paper, having in view to show that the above embedding into **F<sub>at</sub><sup>Q</sup>** can be a useful tool for **IQC**, we prove the faithfulness of such embedding and explore some proof-theoretic properties of **F<sub>at</sub><sup>Q</sup>**. Namely, we show that if the translation of ∃x.A is derivable in **F<sub>at</sub><sup>Q</sup>** then there is a first-order term *t* such that the translation of A[t/x] is derivable in **F<sub>at</sub><sup>Q</sup>**. Note that in order to obtain *t* we do not rely on commuting conversions. As an illustration of the embedding’s potential we show that the existence property for **IQC** can be easily derived from the homologous property in the atomic polymorphic context. Not surprisingly, in this alternative proof of the existence property for **IQC** commuting conversions are absent.

The paper is organized as follows: in Section 2 we review the atomic polymorphic system **F<sub>at</sub><sup>Q</sup>** and the sound translation of **IQC** into **F<sub>at</sub><sup>Q</sup>** and we prove some properties used in following sections such as strong normalization and faithfulness for **F<sub>at</sub><sup>Q</sup>**; in Section 3, we prove that **F<sub>at</sub><sup>Q</sup>** enjoys the existence property and in Section 4 we present an alternative proof of the existence property for **IQC**, via natural deduction, without commuting conversions.

## 2. Atomic polymorphism for IQC

The atomic polymorphic calculus used in this paper is **F<sub>at</sub><sup>Q</sup>**, a calculus described in [3],<sup>3</sup> which has first and second-order constants, first-order variables *x, y, z, ...*, sentential variables *X, Y, Z, ...* and the connectives: → (conditional), ∀<sub>1</sub> (first-order universal quantifier) and ∀<sub>2</sub> (second-order universal quantifier). *Atomic formulas* are either second-order constants and variables or expressions of the form P(*t*<sub>1</sub>, ..., *t*<sub>*n*</sub>), where *P* is a *n*-ary relational symbol and *t*<sub>1</sub>, ..., *t*<sub>*n*</sub> are first-order terms. The class of *formulas* is built in the usual way from atomic formulas, closing the class under →, ∀<sub>1</sub> and ∀<sub>2</sub>. I.e. if *A* and *B* are formulas, *x* is a first-order variable and *X* is a second-order variable then *A* → *B*, ∀<sub>1</sub>*x.A* and ∀<sub>2</sub>*X.A* are also formulas. The subscripts in ∀<sub>1</sub> and ∀<sub>2</sub> are usually omitted. In what follows, the language of **F<sub>at</sub><sup>Q</sup>** will be denoted by  $\mathcal{L}$ .

The logic of **F<sub>at</sub><sup>Q</sup>** is *intuitionistic* logic that we formulate in the natural deduction calculus through the usual introduction and elimination rules for →, ∀<sub>1</sub> and ∀<sub>2</sub>:

$$\frac{[A] \quad \vdots \quad B}{A \rightarrow B} \rightarrow\text{I} \qquad \frac{\vdots \quad A}{\forall x.A} \forall_1\text{I} \qquad \frac{\vdots \quad A}{\forall X.A} \forall_2\text{I}$$

where *x* and *X* do not occur free in any undischarged hypothesis (respectively),

<sup>2</sup> In paper [2] **F<sub>at</sub>** was designated by atomic PSOL<sup>i</sup>.

<sup>3</sup> In [3], system **F<sub>at</sub><sup>Q</sup>** was denoted by atomic QSOL<sup>i</sup>.

Download English Version:

<https://daneshyari.com/en/article/11010144>

Download Persian Version:

<https://daneshyari.com/article/11010144>

[Daneshyari.com](https://daneshyari.com)