



On the boolean partial derivatives and their composition

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ABSTRACT

The main goal of this work is to introduce the relation between the partial boolean derivatives of an n -variable boolean function and their directional boolean derivatives.

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1. Introduction and preliminaries

Let $\mathbb{F}_2^n$ be the n -dimensional vector space over the Galois field $\mathbb{F}_2 = \{0, 1\}$, and set $\{e_1, \dots, e_n\}$ as its standard basis, that is,

$$e_1 = (1, 0, \dots, 0), e_2 = (0, 1, 0, \dots, 0), \dots, e_n = (0, \dots, 0, 1). \quad (1)$$

For two vectors $x = (x_1, \dots, x_n) \in \mathbb{F}_2^n$ and $y = (y_1, \dots, y_n) \in \mathbb{F}_2^n$, we can define the XOR addition operation as follows:

$$x \oplus y = (x_1 \oplus y_1, \dots, x_n \oplus y_n) \in \mathbb{F}_2^n. \quad (2)$$

An n -variable boolean function is a map of the form $f : \mathbb{F}_2^n \rightarrow \mathbb{F}_2$. The set of all n -variable boolean functions is denoted by $\mathcal{BF}_n$ and its cardinality is $|\mathcal{BF}_n| = 2^{2^n}$. The vector

$$t_f = (f(v_0), f(v_1), \dots, f(v_{2^n-1})) \in \mathbb{F}_2^{2^n}, \quad (3)$$

where $v_0 = (0, \dots, 0)$, $v_1 = (0, \dots, 0, 1)$, $\dots$, $v_{2^n-1} = (1, \dots, 1)$, is called the truth table of f . Note that for $1 \leq i \leq 2^n - 1$, v_i is the binary representation of i written as a vector of length 2^n .

The usual representation of a boolean function f is by means of its algebraic normal form (ANF for short) which is the n -variable polynomial representation over $\mathbb{F}_2$, that is,

$$f(x_1, \dots, x_n) = a_0 \oplus \bigoplus_{\substack{1 \leq k \leq n \\ 1 \leq i_1, i_2, \dots, i_k \leq n}} a_{i_1, i_2, \dots, i_k} x_{i_1} x_{i_2} \dots x_{i_k}, \quad (4)$$

where $a_0, a_{i_1, \dots, i_k} \in \mathbb{F}_2$. The degree of the ANF is the algebraic degree of the function. The simplest boolean functions, considering their ANF, are the affine boolean functions: $f(x_1, \dots, x_n) = a_0 \oplus a_1 x_1 \oplus a_2 x_2 \oplus \dots \oplus a_n x_n$, where $a_0, a_1, \dots, a_n \in \mathbb{F}_2$. If $a_0 = 0$, we have the linear boolean functions and they are denoted by $l_a(x)$ with $a = (a_1, \dots, a_n) \in \mathbb{F}_2^n$.

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The Hamming weight of a boolean vector x is denoted by $wt(x)$ and is defined as the number of ones in the vector x . In this sense, the Hamming weight of a boolean function f is the Hamming weight of its truth table t_f . An n -variable boolean function f is said to be balanced if its weight is exactly 2^{n-1} , that is, if the number of ones equals the number of zeros of its truth table.

2. The partial derivative of a boolean function

The notion of the boolean derivative was introduced by Vichniac (see [1]) and it is defined as follows:

Definition 1. The partial derivative of an n -variable boolean function f with respect to the i th variable x_i is another n -variable boolean function, $D_i f$, defined as follows:

$$\begin{aligned} D_i f: \mathbb{F}_2^n &\rightarrow \mathbb{F}_2 \\ x &\mapsto D_i f(x) = f(x) \oplus f(x \oplus e_i), \end{aligned} \quad (5)$$

that is,

$$D_i f(x) = f(x_1, \dots, x_i, \dots, x_n) \oplus f(x_1, \dots, x_i \oplus 1, \dots, x_n). \quad (6)$$

The notion of the boolean derivative is very important and useful in, for example, cryptography (see [2]).

Example 1. Let us consider the four-variable boolean function whose ANF is $f(x_1, x_2, x_3, x_4) = 1 \oplus x_3 \oplus x_1 x_2 \oplus x_2 x_3 \oplus x_1 x_2 x_3 \oplus x_2 x_3 x_4$; then, as simple CALCULATIONS show, $D_1 f(x_1, x_2, x_3, x_4) = x_2 \oplus x_2 x_3$.

This definition allows one to state a derivation rule similar to the derivation rule for multivariate polynomials over real numbers:

Lemma 1. Let f be an n -variable boolean function whose ANF is (4). Then for each variable x_i we have

$$f(x) = g_i(x \oplus x_i e_i) \oplus x_i h_i(x \oplus x_i e_i), \quad (7)$$

where h_i and g_i are $(n-1)$ -variable boolean functions which do not depend on the variable x_i . Moreover, if f does not depend on the variable x_i then $h_i = 0$.

Proof. Set $1 \leq i \leq n$; then the n -variable boolean function f can be factored by taking the common factor x_i , and consequently f can be written as follows:

$$f(x) = g_i(x_1, \dots, \widehat{x_i}, \dots, x_n) \oplus x_i h_i(x_1, \dots, \widehat{x_i}, \dots, x_n), \quad (8)$$

where g_i and h_i are $(n-1)$ -variable boolean functions which do not depend on x_i . For the sake of simplicity we set $x \oplus x_i e_i = (x_1, \dots, \widehat{x_i}, \dots, x_n)$; then,

$$f(x) = g_i(x \oplus x_i e_i) \oplus x_i h_i(x \oplus x_i e_i), \quad (9)$$

thus finishing the proof. $\square$

Example 2. Let us consider the four-variable boolean function whose ANF is $f(x_1, x_2, x_3, x_4) = x_1 \oplus x_2 x_3 \oplus x_3 x_4 \oplus x_2 x_3 x_4$. Then,

$$\begin{aligned} f(x_1, x_2, x_3, x_4) &= x_2 x_3 \oplus x_3 x_4 \oplus x_2 x_3 x_4 \oplus x_1 \\ &= x_1 \oplus x_3 x_4 \oplus x_2 (x_3 \oplus x_3 x_4) \\ &= x_1 \oplus x_3 (x_2 \oplus x_4 \oplus x_2 x_4) \\ &= x_1 \oplus x_2 x_3 \oplus x_4 (x_3 \oplus x_2 x_3). \end{aligned} \quad (10)$$

Proposition 1. Let f be an n -variable boolean function. Then,

$$D_i f(x) = h_i(x \oplus x_i e_i). \quad (11)$$

Proof. By definition, $D_i f(x) = f(x) \oplus f(x \oplus e_i)$, and taking into account the last lemma, this yields

$$D_i f(x) = g_i(x \oplus x_i e_i) \oplus x_i h_i(x \oplus x_i e_i) \oplus g_i(x \oplus x_i e_i) \oplus (x_i \oplus 1) h_i(x \oplus x_i e_i) = h_i(x \oplus x_i e_i), \quad (12)$$

thus finishing the proof. $\square$

As a consequence, the partial derivative (with respect to one variable) reduces the algebraic degree of the boolean function by 1.

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