

Weak connectedness of tensor product of digraphs

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ABSTRACT

In this paper we give a complete characterization for digraphs whose tensor product is weakly connected, which solves an open problem posed by Harary and Trauth (1966).

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1. Introduction

For the notation and terminology below on digraphs, see [1,4,9]. A digraph D is a pair $D = (V(D), E(D))$ with $E(D) \subseteq V(D) \times V(D)$. For an arc $(u, v) \in E(D)$, u is said to be the *initial vertex* of the arc and v the *end vertex*. The *out-degree* (resp. *in-degree*) of a vertex $v \in V(D)$ is defined to be the number of arcs having v as an initial (resp. end) vertex. A *source* (resp. *sink*) is a vertex with in-degree (resp. out-degree) equal to 0.

The *directed cycle* C_n is the digraph with vertex set $[n] = \{1, 2, \dots, n\}$ and arc set $\{(1, 2), (2, 3), \dots, (n, 1)\}$. In particular, C_1 consists of a single vertex with a loop. The *directed path* P_n is C_n with the arc $(n, 1)$ removed. The *length* of a directed cycle or path is defined to be the number of arcs it contains.

A digraph D is called *in-functional* (resp. *out-functional*) if the out-degree (resp. in-degree) of each vertex in D equals to 1. Note that the concept of in-functional digraphs here coincides with the concept of functional digraphs in [5].

For $k \geq 1$, a *route* of length k in D from v_1 to v_{k+1} is a pair $L = (S, \sigma)$, where $S = (v_1, v_2, \dots, v_{k+1})$ is a sequence of $k+1$ vertices in D , and $\sigma = (\sigma_1, \sigma_2, \dots, \sigma_k)$ is a sequence consisting of 1 and -1 , such that

$$\sigma_i = 1 \implies (v_i, v_{i+1}) \in E(D), \quad \text{and} \quad \sigma_i = -1 \implies (v_{i+1}, v_i) \in E(D).$$

We call v_1 the *initial vertex* of L , and v_{k+1} the *end vertex*. L is said to be *closed* if $v_1 = v_{k+1}$, and *alternating* if $\sigma_i = (-1)^i$ for $1 \leq i \leq k$. The type and weight of L are defined to be σ and $\sum_{i=1}^k \sigma_i$ respectively. The greatest common divisor (abbreviated GCD) of the weights of all closed routes in D is called the *weight* of D , where we define $\text{GCD}(0, 0) = 0$. Let $w(D)$ denote the weight of D . Then we have $w(D) = 0$ if all closed routes in D have weight 0.

Let $L = (S, \sigma)$ be a route with $S = (v_1, \dots, v_{k+1})$ and $\sigma = (\sigma_1, \dots, \sigma_k)$. Let $R = (T, \tau)$ be another route with $T = (u_1, \dots, u_{l+1})$ and $\tau = (\tau_1, \dots, \tau_l)$. If $v_{k+1} = u_1$, then the *composite route* LR is defined to be $LR = ((v_1, \dots, v_{k+1}, u_2, \dots, u_{l+1}), (\sigma_1, \dots, \sigma_k, \tau_1, \dots, \tau_l))$.

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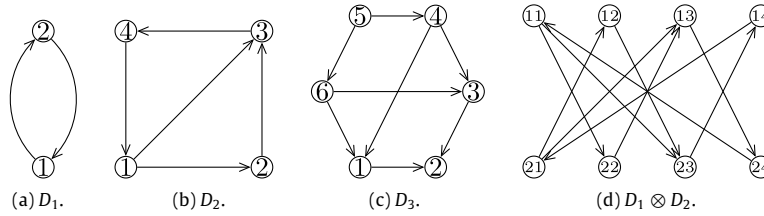


Fig. 1. Examples of weights and tensor product.

A *walk* is a route with type $(1, 1, \dots, 1)$. There are three kinds of connectedness for digraphs defined as follows. A digraph is said to be *strong* (strongly connected) if for each pair u, v of its vertices, there is a walk from u to v and a walk from v to u . It is *unilateral* (unilaterally connected) if for each pair u, v of its vertices, there is a walk from u to v or a walk from v to u . It is *weak* (weakly connected) if for each pair u, v of its vertices, there is a route from u to v . Note that the digraph P_1 is not weak.

If D is a weak digraph with $w(D) = 0$, then for any $u, v \in V(D)$, all routes in D from u to v have the same weight, which we denote by $w(u, v)$, and we define the *diameter* of D to be $d(D) = \max\{w(u, v) | u, v \in V(D)\}$.

Let D_1 and D_2 be digraphs. Their tensor product $D_1 \otimes D_2$ is defined to be the digraph with vertex set $V_1 \times V_2$, and $((u_1, v_1), (u_2, v_2))$ is an arc in $D_1 \otimes D_2$ precisely if (u_1, u_2) and (v_1, v_2) are arcs in D_1 and D_2 respectively. See [8] for more details. Note that we also use the symbol $\otimes$ for the matrix tensor in this paper.

Example 1. We consider the weights of digraphs in Fig. 1. Let $W(D)$ denote the set consisting of weights of closed routes in D with no repeated vertices, other than the repetition of the initial and end vertices. Then we have $W(D_1) = \{0, \pm 2\}$, $W(D_2) = \{0, \pm 1, \pm 3, \pm 4\}$ and $W(D_3) = \{0\}$. Thus $w(D_1) = 2$, $w(D_2) = 1$ and $w(D_3) = 0$. Since $L = ((5, 6, 3, 2), (1, 1, 1))$ is one of the routes in D_3 with maximal weight, we have $d(D_3) = 3$. The digraph in Fig. 1(d) is the tensor product of D_1 and D_2 , where we use ij to denote the vertex (i, j) .

Considering the connectedness of the tensor product of digraphs, three problems arise naturally: characterize digraphs whose tensor product has each of the three kinds of connectedness. McAndrew [8] solved one of the problems by characterizing digraphs whose tensor product is strong. Harary and Trauth [6] solved the problem for a unilaterally connected tensor product. The problem for a weakly connected tensor product was introduced and analyzed in [6]. Hartfiel and Maxson [7] gave a partial answer to this problem by showing that $D_1 \otimes D_2$ is weak for any weak digraph D_1 if and only if D_2 is chainable. Blondel et al. [2] noticed that this problem is equivalent to characterizing digraphs whose similarity matrices have only positive entries. Despite these contributions on this problem, a convenient characterization for a weakly connected tensor product of digraphs is, to our knowledge, still open.

In this paper, we solve the above problem posed by Harary and Trauth. Let D_1 and D_2 be weak digraphs. In Section 2, we consider the case when D_1 and D_2 are both in-functional or out-functional, and show that $D_1 \otimes D_2$ is weak if and only if the GCD of their weights is 1. In Section 3, we consider the case when D_2 is a directed path of length l , and show that $D_1 \otimes D_2$ is weak if and only if D_1 is l -chainable. In Section 4, we show that the connectedness of the tensor product of two general weak digraphs can be reduced to one of the above two cases, and give a complete characterization for a weakly connected tensor product of digraphs in Theorem 15 and Corollary 16.

2. Tensor product of functional digraphs

Let $L = ((u_1, \dots, u_k), \sigma)$ and $R = ((v_1, \dots, v_k), \sigma)$ be routes with the same type in D_1 and D_2 respectively. Then L and R determine bijectively a route $L_0 = ((u_1, v_1), \dots, (u_k, v_k)), \sigma$ in $D_1 \otimes D_2$, which we denote by $(u_1, v_1) \xrightarrow{(L,R)} (u_k, v_k)$. We call L and R the projections of L_0 to D_1 and D_2 respectively.

Example 2. Let $L = ((1, 2, 1, 2), (1, 1, -1))$ and $R = ((1, 2, 3, 1), (1, 1, -1))$ be routes in D_1 and D_2 respectively in Fig. 1. Then we have $(1, 1) \xrightarrow{(L,R)} (2, 1) = (((1, 1), (2, 2), (1, 3), (2, 1)), (1, 1, -1))$ being a route from $(1, 1)$ to $(2, 1)$ in $D_1 \otimes D_2$.

The following result gives a necessary condition for a weak tensor product.

Lemma 3. Let D_1 and D_2 be digraphs with $w(D_1) = d_1$ and $w(D_2) = d_2$. If $D_1 \otimes D_2$ is weak, then $\text{GCD}(d_1, d_2) = 1$.

Proof. Let $u \in V(D_1)$ and $(v_1, v_2) \in E(D_2)$. Since $D_1 \otimes D_2$ is weak, there is a route $(u, v_1) \xrightarrow{(L,R)} (u, v_2)$ from (u, v_1) to (u, v_2) for some routes L and R . Let w be the weight of L . Then L is a closed route in D_1 with weight w , and $R\tilde{v}_2v_1$ is a closed route in D_2 with weight $w - 1$, where $\tilde{v}_2v_1$ denotes the route $((v_2, v_1), (-1))$. Since $\text{GCD}(w, w - 1) = 1$, the result follows. $\square$

A *directed in-rooted* (resp. *out-rooted*) *tree* is a weak digraph with the out-degree (resp. in-degree) of one distinguished vertex, called the root, equal to 0, and the out-degree (resp. in-degree) of other vertices equal to 1. By [5, Theorem 1], a weak in-functional digraph consists of a directed cycle C_p , called the underlined cycle of the digraph, and some directed in-rooted

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