



# A Kraft–McMillan inequality for free semigroups of upper-triangular matrices



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## ABSTRACT

We prove a version of the Kraft–McMillan inequality for free semigroups of upper-triangular integer matrices of any dimension.

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## 1. Introduction

Let  $S$  be a semigroup and let  $X$  be a subset of  $S$ . Then  $X$  is called a *code* if for any integers  $m, n \geq 1$  and any elements  $x_1, \dots, x_m, y_1, \dots, y_n \in X$  the equation

$$x_1 x_2 \dots x_m = y_1 y_2 \dots y_n$$

implies that

$$m = n \quad \text{and} \quad x_i = y_i \quad \text{for } 1 \leq i \leq m.$$

A semigroup  $S$  is called *free* if there exists a code  $X \subseteq S$  such that  $X$  generates  $S$ . For a general introduction to questions concerning freeness of various semigroups see [3].

Let now  $\Sigma$  be an alphabet and let  $\Sigma^*$  be the free monoid generated by  $\Sigma$ . The free submonoids of  $\Sigma^*$  and the codes generating them have been extensively studied. For an excellent introduction to this theory, sometimes called the theory of variable-length codes, see [1].

Most results concerning codes in word monoids do not extend to more general monoids. As an example, consider the monoid  $\text{Tri}(3, \mathbb{N})$  of upper-triangular  $3 \times 3$  matrices having nonnegative integer entries. While in a word monoid a two-element set is a code if and only if the two words are not powers of a common word, the matrices

$$A = \begin{pmatrix} 2 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix} \quad \text{and} \quad B = \begin{pmatrix} 1 & 0 & 0 \\ 0 & 3 & 0 \\ 0 & 0 & 1 \end{pmatrix}$$

are not powers of a common matrix but the set  $\{A, B\}$  still fails to be a code because we have  $AB = BA$ .

Also, in a word monoid it is easy to decide whether or not a given finite set is a code, while in  $\text{Tri}(3, \mathbb{N})$  this problem is undecidable (see [1,2]).

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Let again  $\Sigma$  be an alphabet and assume that  $\pi : \Sigma^* \rightarrow \mathbb{R}_+$  is a Bernoulli distribution. (Here  $\mathbb{R}_+$  is the multiplicative monoid of nonnegative real numbers.) In other words,  $\pi$  is a monoid morphism such that

$$\sum_{\sigma \in \Sigma} \pi(\sigma) = 1.$$

If  $C \subseteq \Sigma^*$ , define

$$\pi(C) = \sum_{c \in C} \pi(c).$$

Now, if  $C$  is a code, then the Kraft–McMillan inequality says that

$$\pi(C) \leq 1$$

(see [1,4,5]). This result can be applied to show that a given subset  $C$  of  $\Sigma^*$  is not a code and, more importantly, it implies that if a code  $C$  satisfies  $\pi(C) = 1$  then  $C$  is a maximal code. Recall that a code  $C$  is *maximal* if there does not exist a code  $D$  such that  $C$  is a proper subset of  $D$ .

The purpose of this paper is to prove a version of the Kraft–McMillan inequality for codes in the monoid of upper-triangular integer matrices of any dimension. We will not use Bernoulli distributions. Instead, as a measure of a matrix we will use its spectral radius. The proof of the inequality will use heavily special properties of upper-triangular matrices. It turns out that the applicability of our inequality is greatly enhanced by the fact that the spectral radius  $\rho(MN)$  of the product of two upper-triangular matrices  $M$  and  $N$  is often less than the product of  $\rho(M)$  and  $\rho(N)$ .

We now outline the contents of the paper. In Section 2 we recall the basic definitions. In Section 3 we state and prove the Kraft–McMillan inequality for codes consisting of upper-triangular integer matrices. In Section 4 we present various examples and explain the connection of our inequality to the classical Kraft–McMillan inequality.

We will assume some familiarity with elementary linear algebra, but otherwise the paper is self-contained.

## 2. Definitions

As usual,  $\mathbb{Z}$  is the set of integers.

If  $k$  is a positive integer,  $\mathbb{Z}^{k \times k}$  is the set of  $k \times k$  matrices having integer entries and  $\text{Tri}(k, \mathbb{Z})$  is the subset of  $\mathbb{Z}^{k \times k}$  consisting of upper-triangular matrices.

If  $M$  is a  $k \times k$  matrix, then the  $(i, j)$ -entry of  $M$  is denoted by  $M_{ij}$  for  $1 \leq i, j \leq k$  and the *diagonal* of  $M$ , denoted by  $\text{diag}(M)$ , is the vector  $(M_{11}, M_{22}, \dots, M_{nn})$ .

If  $M \in \mathbb{Z}^{k \times k}$  and the characteristic values of  $M$  are  $\lambda_1, \dots, \lambda_k$ , then the *spectral radius* of  $M$  is defined by

$$\rho(M) = \max\{|\lambda_1|, \dots, |\lambda_k|\},$$

where  $|\lambda|$  is the absolute value of  $\lambda$ . If  $M \in \text{Tri}(k, \mathbb{Z})$ , then the characteristic values of  $M$  are the diagonal entries  $M_{ii}$  for  $1 \leq i \leq k$  and

$$\rho(M) = \max\{|M_{11}|, |M_{22}|, \dots, |M_{kk}|\}.$$

A matrix  $M \in \mathbb{Z}^{k \times k}$  is called *nilpotent* if there is a positive integer  $n$  such that  $M^n$  is the zero matrix. A matrix  $M \in \text{Tri}(k, \mathbb{Z})$  is nilpotent if and only if its diagonal is the zero vector, or equivalently, if and only if  $\rho(M) = 0$ .

Suppose  $\mathcal{M} \subseteq \mathbb{Z}^{k \times k}$ . If  $n$  is a nonnegative integer then  $\mathcal{M}^n$  is the set defined by

$$\mathcal{M}^n = \{M_1 M_2 \dots M_n \mid M_1, \dots, M_n \in \mathcal{M}\}$$

and  $\mathcal{M}^*$  is the set

$$\mathcal{M}^* = \bigcup_{n \geq 0} \mathcal{M}^n.$$

A set  $\mathcal{M} \subseteq \mathbb{Z}^{k \times k}$  is called a *code* if for all  $s, t \geq 1$  and  $M_1, \dots, M_s, N_1, \dots, N_t \in \mathcal{M}$  the condition

$$M_1 M_2 \dots M_s = N_1 N_2 \dots N_t$$

implies that

$$s = t \quad \text{and} \quad M_i = N_i \quad \text{for all } 1 \leq i \leq s.$$

If  $k$  is a positive integer, a code  $\mathcal{M} \subseteq \text{Tri}(k, \mathbb{Z})$  is called *maximal* if there does not exist a code  $\mathcal{N} \subseteq \text{Tri}(k, \mathbb{Z})$  such that  $\mathcal{M}$  is a proper subset of  $\mathcal{N}$ . (Observe that we consider maximality in the monoid  $\text{Tri}(k, \mathbb{Z})$ .)

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