



# Bounds on locating total domination number of the Cartesian product of cycles and paths



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## ABSTRACT

The problem of placing monitoring devices in a system in such a way that every site in the safeguard system (including the monitors themselves) is adjacent to a monitor site can be modeled by total domination in graphs. Locating-total dominating sets are of interest when the intruder/fault at a vertex precludes its detection in that location. A total dominating set  $S$  of a graph  $G$  with no isolated vertex is a locating-total dominating set of  $G$  if for every pair of distinct vertices  $u$  and  $v$  in  $V - S$  are totally dominated by distinct subsets of the total dominating set. The locating-total domination number of a graph  $G$  is the minimum cardinality of a locating-total dominating set of  $G$ . In this paper, we study the bounds on locating-total domination numbers of the Cartesian product  $C_m \square P_n$  of cycles  $C_m$  and paths  $P_n$ . Exact values for the locating-total domination number of the Cartesian product  $C_3 \square P_n$  are found, and it is shown that for the locating-total domination number of the Cartesian product  $C_4 \square P_n$  this number is between  $\lceil \frac{3n}{2} \rceil$  and  $\lceil \frac{3n}{2} \rceil + 1$  with two sharp bounds.

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## 1. Introduction

The location of monitoring devices, such as surveillance cameras or fire alarms, used to safeguard a system serves as the motivation for this work. The problem of placing monitoring devices in a system in such a way that every site in the system (including the monitors themselves) is adjacent to a monitor site can be modeled by total domination in graphs. Applications where it is also important that if there is a problem at a facility, the location can be uniquely identified by the set of monitors, can be modeled by a combination of total-domination and locating sets. Locating-total dominating set in graph was introduced by Haynes and Henning [5] and has been studied in [1,5–7] and elsewhere.

Graph theory terminology not presented here can be found in [3,4]. All graphs considered in this paper are simple without isolated vertices. Let  $G = (V, E)$  be a graph with vertex set  $V$  and edge set  $E$ . For any vertex  $v \in G$ , the *open neighborhood* of  $v$  is the set  $N(v) = \{u \in V | uv \in E\}$ , and its *closed neighborhood* is the set  $N[v] = N(v) \cup \{v\}$ . We denote the degree of a vertex  $v$  in  $G$  by  $d_G(v)$ , or simply by  $d(v)$  if the graph  $G$  is clear from the text. For any  $S \subseteq V$ ,  $N(S) = \cup_{v \in S} N(v)$ . Let  $\langle S \rangle$  denote the graph induced by  $S$ . For two vertices  $u, v \in V$ , the *distance* between  $u$  and  $v$  is  $d(u, v)$ . The distance between a vertex  $u$  and a set  $S$  of vertices in a graph is defined as  $d(u, S) = \min\{d(u, v) | v \in S\}$ . If  $S$  and  $T$  are two vertex disjoint subsets of  $V$ , then we denote the number of all edges of  $G$  that join a vertex of  $S$  and a vertex of  $T$  by  $e[S, T]$ . Throughout this paper, we use  $C_n$  and  $P_n$  to denote a cycle and a path of the order  $n$ , respectively.

For graphs  $G_1$  and  $G_2$ , the *Cartesian product*  $G_1 \square G_2$  is the graph with vertex set  $V(G_1) \times V(G_2)$  where two ver-

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tices  $(u_1, v_1)$  and  $(u_2, v_2)$  are adjacent if and only if either  $u_1 = u_2$  and  $v_1 v_2 \in E(G_2)$  or  $v_1 = v_2$  and  $u_1 u_2 \in E(G_1)$ .

Let  $\{v_{ij} | 1 \leq i \leq m, 1 \leq j \leq n\}$  be the vertex set of  $G = C_m \square P_n$  so that the subgraph induced by  $\mathcal{H}_i = \{v_{i1}, v_{i2}, \dots, v_{in}\}$  is isomorphic to the path  $P_n$  for each  $1 \leq i \leq m$  and the subgraph induced by  $\mathcal{V}_j = \{v_{1j}, v_{2j}, \dots, v_{mj}\}$  is isomorphic to the cycle  $C_m$  for each  $1 \leq j \leq n$ .

A subset  $S \subseteq V$  is a *total dominating set* (abbreviated, TDS) if every vertex in  $V$  has a neighbor in  $S$ . The *total domination number* of  $G$ , denoted by  $\gamma_t(G)$ , is the minimum cardinality of a total dominating set of  $G$ . Total domination was introduced by Cockayne et al. [2] and is now well-studied in graph theory [3,4].

A total dominating set  $S$  in a graph  $G = (V, E)$  is a *locating-total dominating set* (abbreviated, LTDS) of  $G$  if for every pair of distinct vertices  $u$  and  $v$  in  $V - S$ ,  $N(u) \cap S \neq N(v) \cap S$ . The minimum cardinality of a locating-total dominating set is the *locating-total domination number*  $\gamma_t^L(G)$ . We call a  $\gamma_t^L(G)$ -set a *locating-total dominating set* in  $G$  of cardinality  $\gamma_t^L(G)$ .

In [5], Haynes et al. gave a lower bound on the locating-total domination number of a tree in terms of order and characterized the extremal tree as achieving equality in the lower bound. In [1], Chen and Sohn established lower and upper bounds on the locating-total domination number of trees and constructively characterized the extremal trees achieving the bounds. In [6], Henning and Löwenstein showed that the locating-total domination number of a claw-free cubic graph is at most one-half its order and characterized the graphs that achieved this bound. In [7], Henning and Rad gave lower and upper bounds on the locating-total domination number of a graph, showed that the locating-total domination number and total domination number of a connected cubic graph can differ significantly, and investigated the locating-total domination number of grid graph  $P_m \square P_n$  for small  $m$ .

In this paper, we establish upper bounds on locating total domination numbers of the Cartesian product  $C_m \square P_n$  of cycles  $C_m$  and paths  $P_n$ . In particular, we prove that for any positive integer  $n$ ,  $\gamma_t^L(C_3 \square P_n) = n + 1$  and  $\lceil \frac{3n}{2} \rceil \leq \gamma_t^L(C_4 \square P_n) \leq \lceil \frac{3n}{2} \rceil + 1$ , and these bounds are sharp.

**2. Bounds of locating-total domination number of  $C_m \square P_n$**

In this section, we present upper and lower bounds on the locating-total domination number of the Cartesian product of cycles  $C_m$  and paths  $P_n$ .

**Theorem 2.1.** For any positive integers  $m, n$  such that  $m \equiv 0 \pmod{3}$  and  $n \geq 3$ ,  $\gamma_t^L(C_m \square P_n) \leq \frac{1}{3}m(n + 1)$ .

**Proof.** Let  $G \cong C_m \square P_n$ , where  $m = 3t$  for a positive integer  $t$ . For any integers  $i, j$  such that  $1 \leq i \leq 3$  and  $1 \leq j \leq n$ , let  $D_{ij} = \mathcal{V}_j - \cup_{l=0}^{t-1} \{v_{(3l+i)j}\}$ .

If  $n = 3$ , then  $S = D_{11} \cup D_{23}$  is a LTDS of order  $4t = \frac{1}{3}m(n + 1)$  in  $G$ . If  $n = 4$ , then  $S = D_{11} \cup D_{23} \cup (\cup_{i=0}^{t-1} \{v_{(3i+1)4}\})$  is a LTDS of order  $5t = \frac{1}{3}m(n + 1)$  in  $G$ . If  $n = 5$ , then  $S = D_{11} \cup D_{23} \cup D_{35}$  is a LTDS of order  $6t = \frac{1}{3}m(n + 1)$  in  $G$ .

Assume that  $n \geq 6$ . Let  $n = 6k + r$ , where  $k \geq 1$  and  $0 \leq r \leq 5$ . Let  $S_0 = \cup_{j=0}^{k-1} (D_{1(6j+1)} \cup D_{2(6j+3)} \cup D_{3(6j+5)})$ . Then  $|S_0| = 6kt$ .

If  $r = 0$ , then  $S = S_0 \cup (\cup_{i=0}^{t-1} \{v_{(3i+2)n}\})$  is a LTDS of order  $tn + t = \frac{1}{3}m(n + 1)$  in  $G$ . If  $r = 1$ , then  $S = S_0 \cup D_{1n}$  is a LTDS of order  $6kt + 2t = \frac{1}{3}m(n + 1)$  in  $G$ . If  $r = 2$ , then  $S = S_0 \cup D_{1(n-1)} \cup (\cup_{i=0}^{t-1} \{v_{(3i+2)n}\})$  is a LTDS of order  $6kt + 3t = \frac{1}{3}m(n + 1)$  in  $G$ . If  $r = 3$ , then  $S = S_0 \cup D_{1(n-2)} \cup D_{2n}$  is a LTDS of order  $6kt + 4t = \frac{1}{3}m(n + 1)$  in  $G$ . If  $r = 4$ , then  $S = S_0 \cup D_{1(n-3)} \cup D_{2(n-1)} \cup (\cup_{i=0}^{t-1} \{v_{(3i+1)n}\})$  is a LTDS of order  $6kt + 5t = \frac{1}{3}m(n + 1)$  in  $G$ . If  $r = 5$ , then  $S = S_0 \cup D_{1(n-4)} \cup D_{2(n-2)} \cup D_{3n}$  is a LTDS of order  $6(k + 1)t = \frac{1}{3}m(n + 1)$  in  $G$ .

Therefore,  $\gamma_t^L(G) \leq |S| = \frac{1}{3}m(n + 1)$ . This completes the proof.

**Theorem 2.2.** For any positive integers  $m, n$  such that  $m \equiv 1 \pmod{3}$  and  $n \geq 3$ ,

$$\gamma_t^L(C_m \square P_n) \leq \begin{cases} 11, & m = 4, n = 7; \\ 2m, & n = 5; \\ \frac{1}{3}(m - 1)(n + 1) + \lceil \frac{n}{2} \rceil, & \text{otherwise.} \end{cases}$$

**Proof.** Let  $G \cong C_m \square P_n$ , where  $m = 3t + 1$  for a positive integer  $t$ . For any integers  $i, j$  such that  $1 \leq i \leq 4$  and  $1 \leq j \leq n$ , let  $D_{ij} = \mathcal{V}_j - \cup_{l=0}^{t-1} \{v_{(3l+i)j}\}$  and let  $\lambda = \frac{1}{3}(m - 1)(n + 1) + \lceil \frac{n}{2} \rceil$ .

If  $n = 3$ , then  $S = D_{11} \cup D_{23}$  is a LTDS of order  $4t + 2 = \lambda$  in  $G$ . If  $n = 4$ , then  $S = D_{11} \cup D_{23} \cup (\cup_{i=0}^{t-1} \{v_{(3i+1)4}\})$  is a LTDS of order  $5t + 2 = \lambda$  in  $G$ . If  $n = 5$ , then  $S = \mathcal{V}_2 \cup \mathcal{V}_4$  is a LTDS of order  $2m$  in  $G$ . If  $m = 4$  and  $n = 7$ , then  $S = D_{42} \cup \mathcal{V}_6 \cup \{v_{11}, v_{34}, v_{44}, v_{15}\}$  is a LTDS of order 11 in  $G$ .

Assume that  $n \geq 6$ . Let  $n = 6k + r$ , where  $k \geq 1$  and  $0 \leq r \leq 5$ . Let  $S_0 = \cup_{j=0}^{k-1} (D_{1(6j+1)} \cup D_{2(6j+3)} \cup D_{3(6j+5)})$ . Then  $|S_0| = 6kt + 3k$ .

If  $r = 0$ , then  $S = S_0 \cup (\cup_{i=0}^{t-1} \{v_{(3i+2)n}\})$  is a LTDS of order  $6kt + 3k + t = \lambda$  in  $G$ . If  $r = 1$ , then  $S = S_0 \cup D_{1n}$  is a LTDS of order  $6kt + 3k + 2t + 1 = \lambda$  in  $G$ . If  $r = 2$ , then  $S = S_0 \cup D_{1(n-1)} \cup (\cup_{i=0}^{t-1} \{v_{(3i+2)n}\})$  is a LTDS of order  $6kt + 3k + 3t + 1 = \lambda$  in  $G$ . If  $r = 3$ , then  $S = S_0 \cup D_{1(n-2)} \cup D_{2n}$  is a LTDS of order  $6kt + 3k + 4t + 2 = \lambda$  in  $G$ . If  $r = 4$ , then  $S = S_0 \cup D_{1(n-3)} \cup D_{2(n-1)} \cup (\cup_{i=0}^{t-1} \{v_{(3i+1)n}\})$  is a LTDS of order  $6kt + 3k + 5t + 2 = \lambda$  in  $G$ . If  $r = 5$ , then  $S = S_0 \cup D_{1(n-4)} \cup D_{2(n-2)} \cup D_{3n}$  is a LTDS of order  $6kt + 3k + 6t + 3 = \lambda$  in  $G$ . This completes the proof.

**Theorem 2.3.** For any positive integers  $m, n$  such that  $m \equiv 2 \pmod{3}$  and  $n \geq 3$ ,

$$\gamma_t^L(C_m \square P_n) \leq \begin{cases} \frac{2}{5}mn, & n = 5 \text{ or } m = 5, \\ & n \equiv 0 \pmod{5}; \\ \frac{1}{3}(m + 1)(n + 1) - 1, & \text{otherwise.} \end{cases}$$

**Proof.** Let  $G \cong C_m \square P_n$ , where  $m = 3t + 2$  for a positive integer  $t$ . For any integers  $i, j$  such that  $1 \leq i \leq 5$  and  $1 \leq j \leq n$ , let  $D_{ij} = \mathcal{V}_j - \cup_{l=0}^{t-1} \{v_{(3l+i)j}\}$  and let  $\mu = \frac{1}{3}(m + 1)(n + 1) - 1$ .

If  $n = 3$ , then  $S = (D_{11} - \{v_{m1}\}) \cup D_{33}$  is a LTDS of order  $4t + 3 = \mu$  in  $G$ . If  $n = 4$ , then  $S = D_{11} \cup D_{23} \cup$

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