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## Minimal *k*-rankings and the rank number of $P_n^2$

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#### 1. Introduction

A *k*-ranking of a graph is a vertex labeling using integers between 1 and *k* inclusive such that any path between two vertices of the same rank contains a vertex of strictly larger rank. A ranking *f* has a *drop vertex x* if the labeling defined by g(v) = f(v) when  $v \neq x$  and g(x) < f(x) is still a ranking. It was shown by Jamison [9] that a ranking is minimal if and only if it contains no drop vertices. The *rank number* of a graph *G* is the smallest *k* such that *G* has a minimal *k*-ranking. The *arank number* of a graph *G* is the largest *k* such that *G* has a minimal *k*-ranking. When the value of *k* is unimportant, we will refer to a *k*-ranking simply as a ranking.

Recall that a vertex coloring of a graph is a vertex labeling in which no two adjacent vertices have the same label. Hence a *k*-ranking is a vertex coloring with an additional condition imposed. Then similar to the chromatic number, the *rank number*  $\chi_r(G)$  is the smallest *k* such that *G* has a

\* Corresponding author. E-mail addresses: snovotny@math.jhu.edu (S. Novotny), jpo208@lehigh.edu (J. Ortiz), dansma@rit.edu (D.A. Narayan). minimal *k*-ranking. The *arank number*  $\psi_r(G)$  is the largest *k* such that *G* has a minimal *k*-ranking.

The study of the rank number was motivated by applications including the design of very large scale integration (VLSI) layout and Cholesky factorizations associated with parallel processing [2,6,7,12-14], and [15]. Numerous related papers have since followed [1,8,9,3-5,10], and [11]. Ghoshal, Laskar, and Pillone were the first to investigate minimal k-rankings from a mathematical standpoint [6,7,13], and [14]. The determination of the rank number and the arank number was shown to be NP-complete [14]. The rank number was explored in [1] where they showed  $\chi_r(P_n) = \lfloor \log_2 n \rfloor + 1$ . Rank numbers are known for a few other graph families such as cycles, wheels, complete bipartite graphs, and split graphs (see [7] and [4]). Less is known about the arank number. Arank numbers for complete bipartite graphs and split graphs were established in [7]. Kostyuk, Narayan, and Williams [11] showed  $\psi_r(P_n) = |\log_2(n+1)| + |\log_2(n+1 - (2^{\lfloor \log_2 n \rfloor - 1}))|.$  Recently, the arank number of a cycle was investigated by Kostyuk and Narayan [10].

Throughout the paper  $P_n$  denotes the path on *n* vertices. We use  $P_2 \times P_n$  to denote the *Cartesian product* of  $P_2$ 

<sup>0020-0190/\$ –</sup> see front matter  $\hfill \mathbb{C}$  2008 Elsevier B.V. All rights reserved. doi:10.1016/j.ipl.2008.10.004

and  $P_n$ . The *k*th power of a path, denoted  $P_n^k$ , has vertices  $v_1, v_2, \ldots, v_n$  and edges  $(v_i, v_j)$  for all  $|i - j| \leq k$ .

In this paper we build upon known results and establish new rank numbers. Our main results are stated in the first two theorems. The first theorem gives the rank number of  $P_2 \times P_n$  and a relation to the arank number of a path.

**Theorem 1.**  $\chi_r(P_2 \times P_n) = \psi_r(P_n) + 1.$ 

An interesting result follows by a simple extension. We prove that the minimum k in a minimal k-ranking of  $P_n^2$  is the maximum k in a minimal k-ranking of  $P_n$ .

**Theorem 2.** We have  $\chi_r(P_n^2) = \psi_r(P_n)$ .

#### 2. The rank number of $P_n^2$

We begin by restating two elementary results of Ghoshal et al. [7].

**Lemma 3.** In any minimal ranking of a connected graph *G* the highest label must be unique.

**Proof.** Suppose there exist two vertices u and v which both have the highest label k. Then any path between u and v will not contain a vertex with a higher label. This is a contradiction.  $\Box$ 

The following lemma gives a monotonicity result involving the rank number.

**Lemma 4.** Let *H* be a subgraph of a graph *G*. Then  $\chi_r(H) \leq \chi_r(G)$ .

**Proof.** The proof is straightforward. Suppose  $\chi_r(H) > \chi_r(G)$ . Then we could relabel the vertices of *H* using the corresponding labels used in the ranking of *G*. This produces a ranking with fewer labels, and hence a contradiction.  $\Box$ 

#### 2.1. The ladder graph $L_n$

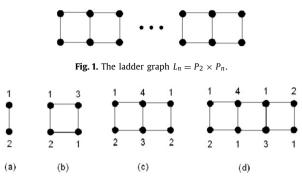
Before investigating the rank number of  $P_n^2$  we describe a family of graphs built using the *Cartesian product*.

**Definition 5.** The *Cartesian product* of *G* and *H* written  $G \times H$  is the graph with vertex set  $v(G) \times V(H)$  specified by putting  $\{u, v\}$  adjacent to (u', v') if and only if (i) u = u' and  $(v, v') \in E(H)$  or (ii) v = v' and  $(u, u') \in E(G)$ .

We define the ladder graph,  $L_n = P_2 \times P_n$  as shown in Fig. 1.

**Example 6.** We give labelings showing  $\chi_r(L_n)$  for  $n \leq 4$  in Fig. 2.

All four cases are easy to prove. The strategy in all of the cases will be to establish a lower bound and then construct a k-ranking where k equals that lower bound. We



**Fig. 2.** Minimal  $\chi_r$ -rankings of small ladders.

have  $\chi_r(L_1) = 2$  since we cannot have two adjacent vertices labeled with a 1 and a minimal 2-ranking is shown in Fig. 2(a). We next consider the case where  $n \ge 2$ . If  $\chi_r(L_2) = 2$  then we can apply Lemma 3 to conclude that the vertex labeled 2 must be unique, and the remaining labels are all labeled 1. This is impossible since we would be forced to have adjacent vertices both with labels of 1. Hence  $\chi_r(L_2) \ge 3$ , and since we have a minimal 3-ranking in Fig. 2(b) it follows that  $\chi_r(L_2) = 3$ . Next we consider the third case where n = 3. By Lemma 4 we have  $\chi_r(L_3) \ge \chi_r(L_2) \ge 3$ . If  $\chi_r(L_3) = 3$  then we can conclude by Lemma 3 that the vertex labeled 3 must be unique. This means that the remaining five vertices would all receive labels that are either 2 or 1. We are guaranteed to have either two adjacent vertices labeled 1 or a P<sub>3</sub> with vertices labeled (in order) 2, 1 and 2. Both cases violate the ranking condition, and hence  $\chi_r(L_3) \ge 4$ . Since a minimal 4-ranking is shown in Fig. 2(c) it follows that  $\chi_r(L_3) = 4$ . Finally for the last case we apply Lemma 4 to conclude that  $\chi_r(L_4) \ge \chi_r(L_3) = 4$ . The given minimal 4-ranking in Fig. 2(d) shows that this bound is tight.

We now seek to extend the minimal  $\chi_r$ -rankings of small ladders to minimal  $\chi_r$ -rankings of large ladders. In our construction we will start with two copies of  $L_s$  and join them with a bridge that consists of either a single 'vertical edge' or a pair of 'vertical edges' along with an appropriate set of edges that connects the pieces to form a larger ladder. These are demonstrated in Figs. 3 and 4. We call these connecting subgraphs, 1-bridges and 2-bridges, respectively. We next present two lemmas which will be used to build minimal  $\chi_r$ -rankings of large ladders. In Lemma 7 we let  $G = L_{2s+1}$  and in Lemma 8 we let  $G = L_{2s+2}$ . Each of the lemmas will show that  $\chi_r(G) \ge \chi_r(L_s) + 2$ . The idea is to first note that since  $L_s$ is a subgraph of G we can apply Lemma 4 to conclude that  $\chi_r(G) \ge \chi_r(L_s)$ . The insertion of the bridge adds either two or four vertices that push the rank number up by at least 2.

**Lemma 7** (First Ladder Lemma). Let *G* be the union of two copies of  $L_s$  connected by a 1-bridge. Then  $\chi_r(G) \ge \chi_r(L_s) + 2$ .

**Proof.** Let the two added vertices be labeled *a* and *b*. We consider cases for different minimal rankings of *G*. We will show in each case there is a vertex with a label greater than or equal to  $\chi_r(L_s) + 2$ .

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