

Ore-type degree conditions for disjoint path covers in simple graphs

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ABSTRACT

A many-to-many k -disjoint path cover of a graph joining two disjoint vertex subsets S and T of equal size k is a set of k vertex-disjoint paths between S and T that altogether cover every vertex of the graph. It is classified as *paired* if each source in S is required to be paired with a specific sink in T , or *unpaired* otherwise. In this paper, we develop Ore-type sufficient conditions for the existence of many-to-many k -disjoint path covers joining arbitrary vertex subsets S and T . Also, an Ore-type degree condition is established for the one-to-many k -disjoint path cover, a variant derived by allowing to share a single source. The bounds on the degree sum are all best possible.

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1. Introduction

Let G be a simple undirected graph, whose vertex and edge sets are denoted by $V(G)$ and $E(G)$, respectively. The order of G is the number of vertices in G . If $(u, v) \in E(G)$, u is adjacent to v or u is a neighbor of v . A path from v_1 to v_m is a sequence of vertices, (v_1, v_2, \dots, v_m) , such that v_j is adjacent to v_{j-1} for every $j \in \{2, \dots, m\}$. A disjoint path cover (DPC for short) of G is a set of paths in G such that every vertex of G belongs to one and only one path.

Definition 1 (*Many-to-many k -disjoint path cover*). Given two disjoint vertex subsets $S = \{s_1, \dots, s_k\}$ and $T = \{t_1, \dots, t_k\}$, each representing k sources and sinks, the many-to-many k -disjoint path cover of G is a disjoint path cover of size k , each of whose paths joins a pair of a source and a sink.

The disjoint path cover is *paired* if every source s_i must be joined with a specific sink t_i . On the other hand, it is *unpaired* if, for some permutation σ on $\{1, \dots, k\}$, P_i is a path from s_i to $t_{\sigma(i)}$ for all $i \in \{1, \dots, k\}$. The sources and sinks are referred to as *terminals*. By these definitions, a paired k -disjoint path cover is always an unpaired k -disjoint path cover. An example of the paired many-to-many DPC is shown in Fig. 1(a).

Definition 2. A graph G is *paired* (resp. *unpaired*) *many-to-many k -coverable* if $|V(G)| \geq 2k$ and there exists a paired (resp. unpaired) many-to-many k -DPC for any disjoint source set S and sink set T of size k each.

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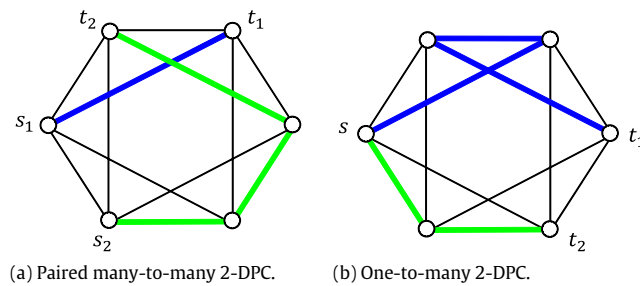


Fig. 1. Examples of disjoint path covers.

Two simpler variants of the many-to-many k -disjoint path cover can be derived by allowing to share a single source and/or a single sink. The first one is of *one-to-many* type with $S = \{s\}$ and $T = \{t_1, \dots, t_k\}$, in which all paths start from the unique source s . The second one is of *one-to-one* type with $S = \{s\}$ and $T = \{t\}$, where all internally disjoint paths connect the unique source and sink. Refer to Fig. 1(b) for an example of the one-to-many DPC.

Definition 3. A graph G is *one-to-many* (resp. *one-to-one*) k -coverable if $|V(G)| \geq k + 1$ and G has a one-to-many (resp. one-to-one) k -DPC for any disjoint source set $S = \{s\}$ and sink set $T = \{t_1, \dots, t_k\}$ (resp. $T = \{t\}$).

The disjoint path cover problem is strongly related to the well-known Hamiltonian problem, which is a fundamental one in graph theory. Actually, a Hamiltonian path joining a pair of vertices in a graph forms a many-to-many, one-to-many, and one-to-one 1-disjoint path covers of the graph. A graph of order $n \geq 3$ is one-to-many 2-coverable if and only if it is Hamiltonian-connected. Moreover, a graph of order $n \geq 3$ is one-to-one 2-coverable if and only if it is Hamiltonian.

One of the core subjects in Hamiltonian graph theory is to develop sufficient conditions for a graph to have a Hamiltonian path/cycle (refer to [11] for a survey). The approaches taken to develop sufficient conditions usually involve degree conditions for providing enough edges to overcome any obstacle to the existence of a Hamiltonian path/cycle. Dirac [4] proved that a graph G of order $n \geq 3$ is Hamiltonian if $d_G(v) \geq n/2$ for every vertex v of G , where $d_G(v)$ denotes the degree of a vertex v in G . Ore [14,15] improved Dirac’s condition as follows:

Theorem 1 (Ore [14,15]). (a) A graph G of order $n \geq 3$ is Hamiltonian if $d_G(u) + d_G(v) \geq n$ for all distinct nonadjacent vertices u and v .

(b) A graph G of order $n \geq 2$ is Hamiltonian-connected if $d_G(u) + d_G(v) \geq n + 1$ for all distinct nonadjacent vertices u and v .

The close relationship between the disjoint path cover problem and the Hamiltonian problem motivates the study of developing degree conditions for a graph to have disjoint path covers. In this paper, we establish Ore-type conditions for the existence of k -disjoint path covers in a simple graph as follows:

1. A graph G of order $n \geq 2k$, where $k \geq 1$, is unpaired many-to-many k -coverable if $d_G(u) + d_G(v) \geq n + k$ for every pair of nonadjacent vertices u and v .
2. A graph G of order $n \geq 2k$, where $k \geq 2$, is paired many-to-many k -coverable if $d_G(u) + d_G(v) \geq n + 3k - 4$ for every pair of nonadjacent vertices u and v .
3. A graph G of order $n \geq k + 1$, where $k \geq 2$, is one-to-many k -coverable if $d_G(u) + d_G(v) \geq n + k - 1$ for every pair of nonadjacent vertices u and v .

Also, we show that all the above three bounds on the degree sum $d_G(u) + d_G(v)$ are the minimum possible. Note that the results for unpaired many-to-many and one-to-many disjoint path covers are generalizations of Ore’s theorem, Theorem 1(b). For the one-to-one disjoint path cover problem, an Ore-type condition was derived by Lin et al. [12]: A graph G of order $n \geq k + 1$, where $k \geq 2$, is one-to-one k -coverable if $d_G(u) + d_G(v) \geq n + k - 2$ for every pair of nonadjacent vertices u and v . Moreover, the bound on the degree sum is tight. A one-to-one k -coverable graph is also known as a k^* -connected graph.

2. Related works and definitions

The disjoint path cover problem finds applications in many areas such as software testing, database design, and code optimization [1,13]. In addition, the problem is concerned with applications where full utilization of network nodes is important [18]. It has been studied with respect to various graphs such as hypercubes [3,5,6], recursive circulants [8,9,18,19], hypercube-like graphs [7,10,18,19], cubes of connected graphs [16,17], k -ary n -cubes [20,21], and spanning connected graphs [12]. Unfortunately, it is NP-complete to determine whether or not there exists a many-to-many, one-to-many, or one-to-one k -DPC for a given pair of terminal sets in a general graph for any fixed $k \geq 1$ [18,19].

A *Hamiltonian path* of a graph G is a path that contains all the vertices of G . A graph is said to be *Hamiltonian-connected* if every pair of distinct vertices are joined by a Hamiltonian path. A *cycle* is a closed path of three or more vertices. A *Hamiltonian cycle* of G is a closed Hamiltonian path, i.e., a cycle that contains all the vertices of G . A graph is called *Hamiltonian* if it has

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