

Contents lists available at ScienceDirect

### **Discrete Mathematics**

journal homepage: www.elsevier.com/locate/disc



# Clique number of the square of a line graph



## Małgorzata Śleszyńska-Nowak

Faculty of Mathematics and Information Science, Warsaw University of Technology, Warsaw, Poland

#### ARTICLE INFO

Article history:
Received 31 July 2015
Received in revised form 31 December 2015
Accepted 5 January 2016
Available online 29 January 2016

Keywords: Strong chromatic index Clique number of the square of a line graph Fractional strong chromatic index

#### ABSTRACT

We prove that the clique number of the square of a line graph of a graph G is at most  $1.5\Delta_G^2$  and that the fractional strong chromatic index of G is at most  $1.75\Delta_G^2$ .

An *edge coloring* of a graph G is strong if each color class is an induced matching of G. The *strong chromatic index* of G, denoted by  $\chi_s'(G)$ , is the minimum number of colors for which G has a strong edge coloring. The strong chromatic index of G is equal to the chromatic number of the square of the line graph of G. The chromatic number of the square of the line graph of G is greater than or equal to the clique number of the square of the line graph of G, denoted by G(G).

In this note we prove that  $\omega(L) \leq 1.5 \Delta_G^2$  for every graph G. Our result allows to calculate an upper bound on the fractional strong chromatic index of G, denoted by  $\chi_{fS}'(G)$ . We prove that  $\chi_{fS}'(G) \leq 1.75 \Delta_G^2$  for every graph G.

© 2016 Elsevier B.V. All rights reserved.

#### 1. Introduction

A *strong edge coloring* of a graph G is an edge coloring in which every color class is an induced matching, that is, any two vertices that belong to distinct edges of the same color are not adjacent (in particular, no two edges of the same color intersect). The strong chromatic index of G, denoted by  $\chi'_s(G)$ , is the minimum number of colors in any strong edge coloring of G.

Concept of the strong edge coloring was introduced around 1985 by Erdős and Nešetřil [8]. They conjectured that for every graph G, with maximum degree  $\Delta_G$ ,  $\chi_s'(G) \leq \frac{5}{4}\Delta_G^2$ . The example of graph obtained from the cycle of length five by replacing each vertex by an independent set of size  $\frac{\Delta}{3}$  shows that this bound, if true, is tight.

replacing each vertex by an independent set of size  $\frac{\Delta}{2}$  shows that this bound, if true, is tight.

The trivial bound on the strong chromatic index of G is  $2\Delta_G^2 - 2\Delta_G + 1$ . Molloy and Reed [11] proved that  $\chi_s'(G) \leq (2-\epsilon)\Delta_G^2$  for  $\Delta_G$  sufficiently large, where  $\epsilon$  is a small constant around  $\frac{1}{50}$ . Recently Bruhn and Joos [4] improved it to the  $1.93\Delta_G^2$  for  $\Delta_G$  sufficiently large.

We approach this problem from a different angle. A *line graph of G* is a graph whose each vertex represents an edge of G and two vertices are adjacent if and only if their corresponding edges are incident in G. A *square of a graph H* is a graph with the same set of vertices as H, in which two vertices are adjacent when their distance in H is at most G. The strong chromatic index of the graph G is equal to the chromatic number of the square of the line graph of G (say G). The chromatic number of G is greater than or equal to the clique number of G (denoted by G), so finding bounds on G0 is also an interesting problem. It is not known if the clique number of G1 is bounded by G2 [1]. Chung et al. [6] proved that if G3 is a clique, then G4 has at most G5 edges. Faudree et al. [9] proved that if G6 is a bipartite graph, then the clique number of G4 is at most G5. The example of G6 with G7 shows that this bound is tight. Recently Bruhn and Joos [4] proved that for G7 with G8 equals G9 the clique number of G8 is at most G9. We improve this result.

In this note, we focus on two topics related with the strong chromatic index of graphs. The first problem is finding the upper bound on the clique number of the square of the line graph of *G*. In our main theorem we prove an upper bound for the general case:

**Theorem 1.** Let G be a simple graph and L be a square of the line graph of G. Then the clique number of L is at most  $1.5\Delta_C^2$ .

We also present (Theorem 4) a new proof of known bounds for bipartite graphs ( $\Delta_G^2$ ). The second subject is a fractional strong chromatic index, denoted by  $\chi'_{fs}(G)$ . Using Theorem 1 we prove an upper bound on  $\chi'_{fs}(G)$ :

**Theorem 2.** Let G be a simple graph. Then the fractional strong chromatic index of G is at most  $1.75\Delta_G^2$ .

#### 2. Upper bound on the clique number of the square of the line graph of G

For any two different edges from a graph G we define  $dist_G(e, f)$  as the number of edges in the shortest path between e and f plus 1, i.e.  $dist_G(e, f) = 1$  iff e and f intersect.

**Remark 3.** Let *G* be a simple graph, *L* be a square of the line graph of *G* and *H* be a subgraph of *G*, such that for each  $e, f \in E(H)$  we have  $dist_G(e, f) \le 2$ . Then  $\omega(L) = |E(H)|$ .

#### 2.1. Bipartite graphs

To better understand our method, first we present a new proof of a known bound for bipartite graphs. We use the same technique in the proof of our main theorem (which is more complicated).

**Theorem 4.** Let G be a simple bipartite graph and L be a square of the line graph of G. Then the clique number of L is at most  $\Delta_G^2$ .

**Proof.** Let H be a subgraph of G, such that for each  $e, f \in E(H)$  we have  $dist_G(e, f) \leq 2$ . We show that H has at most  $\Delta_G^2$  edges.

Consider a vertex  $v \in V(H)$  of degree  $\Delta_H$ . Edges of the graph H can be divided into following sets (see Fig. 1):

- 1.  $A = \{e \in E(H) : v \in e\}, |A| = \Delta_H.$ 
  - A is the set of edges of H which are incident to v.
  - Notice that all other edges of the graph H are in a distance at most 2 to all edges from A.
- 2.  $B = \{e \in E(H) : e \notin A \land \exists_{f \in A} e \text{ and } f \text{ are adjacent}\}, |B| \leq \Delta_H(\Delta_H 1).$ 
  - B is the set of edges of H which are adjacent to edges from A and are not contained in A.
- 3.  $C = \{e \in E(H) : \exists_{f \in E(G)} (v \in f \land f \notin E(H) \land e \text{ and } f \text{ are adjacent})\}, |C| \le (\Delta_G \Delta_H)\Delta_H.$ 
  - C is the set of edges of H which are adjacent to edges of G H which are incident to v.
- 4.  $D = \{e \in E(H) : e \notin C \land \forall_{f \in A} dist_G(e, f) = 2\}.$ 
  - D is the set of edges of H which are at the distance 2 to all edges from A and are not contained in C.
  - Let *S* be the subgraph of *H* induced by D.

We define a super vertex as a vertex which is adjacent in G to all neighbors of v from H, and it is not v. Because G is bipartite, each edge from S contains exactly one super vertex. We have at most  $\Delta_G - 1$  super vertices in G. Furthermore each super vertex is incident in G to at most G from neighbors of G from G from neighbors of G from G from G from neighbors of G from neighbors of G from G fr

$$(\Delta_G - 1)(\Delta_G - \Delta_H)$$
.

Now we can sum up the number of edges in H

$$|E(H)| \leq \Delta_H + \Delta_H(\Delta_H - 1) + (\Delta_G - \Delta_H)\Delta_H + (\Delta_G - 1)(\Delta_G - \Delta_H) = \Delta_G^2 - \Delta_G + \Delta_H \leq \Delta_G^2. \quad \Box$$

#### 2.2. General case

In our proof we use the following lemma.

**Lemma 5.** Let G be a graph with maximum degree  $\Delta$ , p and w be integers such that  $\Delta \leq p \leq w$  and  $\Delta > w - p$ . Consider p vertex covers of G (not necessarily different) such that each vertex cover contains at most w vertices. Moreover assume that for each vertex  $v \in V(G)$  we have  $\deg(v) \leq w - a$ , where a is a number of vertex covers which contain v. Then the graph G has at most  $w^2 - \frac{pw}{2}$  edges.

## Download English Version:

# https://daneshyari.com/en/article/4647128

Download Persian Version:

https://daneshyari.com/article/4647128

<u>Daneshyari.com</u>