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## Note Element deletion changes in dynamic coloring of graphs

ABSTRACT

respectively.

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#### 1. Introduction

In this paper, all graphs G = (V, E) are finite, simple and undirected. For  $v \in V$ ,  $N_G(v)$  is the set of vertices adjacent to v, and the degree of v, denoted by  $d_G(v)$ , is  $|N_G(v)|$ . We use  $\Delta(G)$  and  $\delta(G)$  to denote the maximum degree and minimum degree of G, respectively. When the graph G is understood from the context, we often omit the subscript G, and use  $\delta$ ,  $\Delta$  for  $\delta(G), \Delta(G)$ , respectively. If  $uv \in E$ , then u is a **neighbor** of v. For  $W \subseteq V, G - W$  denotes the graph obtained from G by deleting the vertices in W together with their incident edges. If  $W = \{w\}$ , we often write G - w for  $G - \{w\}$ . If  $U \subseteq V$ , then G[U] denotes the graph on U whose edges are precisely the edges of G with both ends in U. Let  $C_n$  and  $P_n$  denote a cycle and a path on *n* vertices, respectively. In a graph G, an **elementary subdivision** of an edge  $e = uv \in E(G)$  is the operation of replacing e with a path  $uv_e v$  through a new vertex  $v_e$ . A graph H is a **subdivision** of a graph G if H can be obtained from G by a sequence of elementary subdivisions. For a real number x, we use  $\lceil x \rceil$  to denote the least integer no less than x.

For an integer k > 0, let  $\overline{k} = \{1, 2, \dots, k\}$ . If  $S \subseteq V(G)$  is a subset and  $c : V(G) \mapsto \overline{k}$  is a mapping, then define  $c(S) = \{c(x) : x \in S\}$ . A **dynamic** *k*-coloring of a graph *G* is a mapping  $c : V(G) \mapsto \overline{k}$  satisfying both of the following:

(C1) If  $uv \in E(G)$ , then  $\varphi(u) \neq \varphi(v)$ , and

(C2) for each vertex  $v \in V(G)$ ,  $|c(N(v))| \ge \min\{2, d_G(v)\}$ .

The **dynamic chromatic number**  $\chi_d(G)$  is the smallest integer k such that G has a dynamic k-coloring. Dynamic coloring was first introduced in [12,9], and is a special case of the r-hued colorings [8,7,13] when r = 2. The study of dynamic coloring has drawn lots of attention, as seen in [1-6,8,9,12,10,11,13,14], among others.

Unlike classic colorings, a subgraph of a graph G may have a bigger dynamic chromatic number than G. A natural problem is to investigate the differences between  $\chi_d(G)$  and  $\chi_d(G-e)$ , and between  $\chi_d(G)$  and  $\chi_d(G-v)$ . This motivates the current study. In Section 2, we will investigate the best possible bounds for the differences between  $\chi_d(G-e)$  and  $\chi_d(G)$ , and between  $\chi_d(G - v)$  and  $\chi_d(G)$ .

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A proper vertex k-coloring of a graph G is dynamic if for every vertex v with degree at

least 2, the neighbors of v receive at least two different colors. The smallest integer k such

that G has a dynamic k-coloring is the dynamic chromatic number  $\chi_d(G)$ . In this paper the

differences between  $\chi_d(G)$  and  $\chi_d(G-e)$ , and between  $\chi_d(G)$  and  $\chi_d(G-v)$  are investigated

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#### 2. Comparisons between $\chi_d(G)$ and $\chi_d(G - e)$ , and between $\chi_d(G)$ and $\chi_d(G - v)$

It is well known that if *H* is a subgraph of a graph *G*, then  $\chi(G) \ge \chi(H)$ . However, there exist graphs *G* with a subgraph *H* such that  $\chi_d(H) > \chi_d(G)$ . For example, let *G* be the 5-cycle with one chord, and let *H* be the 5-cycle, then it is routine to verify that  $\chi_d(G) = 4$  but  $\chi_d(H) = 5$ .

In this section, we investigate tight bounds for the change of the dynamic chromatic number when an edge or a vertex is being removed. We start with a lemma, which follows from definition immediately.

**Lemma 2.1.** If G is a connected graph on at least 2 vertices, then  $\chi_d(G) \leq 2$  is and only if  $G \in \{K_1, K_2\}$ .

#### Theorem 2.1. Each of the following holds.

(i) Let G be a connected graph with  $|V(G)| \ge 3$ . Then for any edge  $e = uv \in E(G)$ ,

$$\chi_d(G) - 2 \le \chi_d(G - e) \le \chi_d(G) + 2.$$

(ii) There exists a graph G such that  $\chi_d(G - e) = \chi_d(G) + 2$  for at least one edge  $e \in E(G)$ .

(iii) If a connected graph G satisfies that  $\chi_d(G-e) = \chi_d(G) - 2$  for at least one edge e in G, then  $G = C_5$ .

**Proof.** (i) Let  $k_1 = \chi_d(G - e)$ , and let  $c_1 : V(G - e) \mapsto \overline{k}_1$  be a dynamic  $k_1$ -coloring of G - e. Obtain a new coloring  $c'_1$  from  $c_1$  by defining

$$c'_{1}(z) = \begin{cases} c_{1}(z) & \text{if } z \notin \{u, v\} \\ k_{1} + 1 & \text{if } z = u \\ k_{1} + 2 & \text{if } z = v. \end{cases}$$

By definition,  $c'_1 : V(G) \mapsto \overline{k_1 + 2}$  is a dynamic  $(k_1 + 2)$ -coloring of *G*, and so  $\chi_d(G) \le \chi_d(G - e) + 2$ .

Now let  $k_2 = \chi_d(G)$  and  $c_2 : V(G) \mapsto \overline{k_2}$  be a dynamic  $k_2$ -coloring of G. Since  $|V(G)| \ge 3$  and since G is connected, there exists  $x \in N_G(u) - \{v\}$  or  $y \in N_G(v) - \{u\}$ . Choose such x and y so that  $|\{x, y\}|$  is maximized. If  $|\{x, y\}| = 1$ , then by the maximality of  $|\{x, y\}|$ , and since G is connected, we must have  $d_G(u) \le 2$  and  $d_G(v) \le 2$ . In this case, we have  $\chi_d(G) = \chi_d(G-e)$ , and so  $\chi_d(G) \le \chi_d(G-e) + 2$ . Hence we assume that  $x \ne y$ . Obtain a new coloring  $c'_2$  from  $c_2$  by defining

$$c'_{2}(z) = \begin{cases} c_{2}(z) & \text{if } z \notin \{x, y\} \\ k_{2} + 1 & \text{if } z = x \\ k_{2} + 2 & \text{if } z = y. \end{cases}$$

By definition,  $c'_2 : V(G - e) \mapsto \overline{k_2 + 2}$  is a dynamic  $(k_2 + 2)$ -coloring of G - e, and so  $\chi_d(G - e) \leq \chi_d(G) + 2$ . This proves (i). (ii) For an integer  $r \geq 4$ , let H be a complete r-partite graph with partite sets  $V_1, V_2, \ldots, V_r$ , such that  $|V_i| \geq 2$  for each i with  $1 \leq i \leq r$ , and let u and v be two new vertices. Let G be the graph obtained from H by adding a new edge uv to H and by joining u to every vertex in  $V_1$  and joining v to every vertex in  $V_2$ . It is routine to verify that  $\chi_d(G) = \chi(G) = r$ , and that  $\chi_d(G - uv) = r + 2$ , since the vertices in each of  $V_1$  and  $V_2$  must be colored with at least two colors.

(iii) Let *G* be a connected graph with at least one edge such that  $\chi_d(G - e) = \chi_d(G) - 2$  for some edge  $e = uv \in E(G)$ , and let  $k = \chi_d(G - e)$ . If  $\chi_d(G - e) \le 2$ , then by Lemma 2.1,  $G \in \{K_2, P_3\}$ , contrary to the assumption that  $\chi_d(G - e) = \chi_d(G) - 2$  for some  $e \in E(G)$ . Hence we assume that  $k = \chi_d(G - e) \ge 3$ .

Let  $c : V(G - e) \mapsto \overline{k}$  be a dynamic *k*-coloring. Assume without loss of generality, that  $d_G(u) \ge d_G(v)$ . If  $d_G(v) = 1$ , then v is an isolated vertex of G - e. As  $k \ge 3$ , we can pick a vertex  $u' \in N_G(u) - \{v\}$  and redefine  $c(v) \in \overline{k} - \{c(u), c(u')\}$  to obtain a *k*-coloring of G, contrary to the assumption that  $\chi_d(G - e) = \chi_d(G) - 2$ . If  $d_G(u) \ge 3$ , then by  $k \ge 3$ , we can redefine c(u) = k + 1 to obtain a (k + 1)-coloring of G, contrary to the assumption that  $\chi_d(G - e) = \chi_d(G) - 2$ . If  $d_G(u) \ge 3$ , then by  $k \ge 3$ , we can may assume that  $d_G(u) = d_G(v) = 2$ . Let  $N_G(u) = \{v, u'\}$ ,  $N_G(v) = \{u, v'\}$ . We have the following claims.

Claim 1.  $u' \neq v'$ .

If u' = v', then obtain a new coloring c' from c by defining

$$c'(z) = \begin{cases} c(z) & \text{if } z \neq u \\ k+1 & \text{if } z = u \end{cases}$$

By definition,  $c' : V(G) \mapsto \overline{k+1}$  is a dynamic (k+1)-coloring of *G*, contrary to the assumption that  $\chi_d(G-e) = \chi_d(G) - 2$ . Thus Claim 1 must hold.

**Claim 2.** 
$$c(u) = c(v') \neq c(u') = c(v)$$

If  $c(u) \neq c(v')$ , then obtain a new coloring c'' from c by defining

$$c''(z) = \begin{cases} c(z) & \text{if } z \neq v \\ k+1 & \text{if } z = v. \end{cases}$$

(1)

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