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## Longest cycles in almost regular 3-partite tournaments

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#### Abstract

If D is a digraph, then we denote by V(D) its vertex set. A multipartite or c-partite tournament is an orientation of a complete c-partite graph. The global irregularity of a digraph D is defined by

$$i_g(D) = \max{\{\max(d^+(x), d^-(x)) - \min(d^+(y), d^-(y)) \mid x, y \in V(D)\}}.$$

If  $i_g(D) = 0$ , then D is regular, and if  $i_g(D) \le 1$ , then D is called almost regular. In 1997, Yeo has shown that each regular multipartite tournament is Hamiltonian. This remains valid for almost all almost regular c-partite tournaments with  $c \ge 4$ . However, there exist infinite families of almost regular 3-partite tournaments without any Hamiltonian cycle. In this paper we will prove that every vertex of an almost regular 3-partite tournament D is contained in a directed cycle of length at least |V(D)| - 2. Examples will show that this result is best possible.

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#### 1. Terminology

A *c-partite* or *multipartite tournament* is an orientation of a complete *c*-partite graph. By a *cycle* (*path*) we mean a directed cycle (directed path).

In this paper all digraphs are finite without loops or multiple arcs. The vertex set and the arc set of a digraph D are denoted by V(D) and E(D), respectively. If xy is an arc of a digraph D, then we write  $x \to y$  and say x dominates y. If X and Y are two disjoint subsets of V(D) or subdigraphs of D such that every vertex of X dominates every vertex of Y, then we say that X dominates Y, denoted by  $X \to Y$ . Furthermore,  $X \leadsto Y$  denotes the property that there is no arc from Y to X.

The *out-neighborhood*  $N_D^+(x) = N^+(x)$  of a vertex x is the set of vertices dominated by x, and the *in-neighborhood*  $N_D^-(x) = N^-(x)$  is the set of vertices dominating x. The numbers  $d_D^+(x) = d^+(x) = |N^+(x)|$  and  $d_D^-(x) = d^-(x) = |N^-(x)|$  are the *outdegree* and *indegree* of x, respectively. Let

$$\delta^{-}(D) = \delta^{-} = \min\{d^{-}(x) \mid x \in V(D)\},\$$

$$\delta^{+}(D) = \delta^{+} = \min\{d^{+}(x) \mid x \in V(D)\},\$$

and  $\delta(D) = \delta = \min{\{\delta^+, \delta^-\}}$  be the *minimum degree* of the digraph D.

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The *global irregularity* of a digraph D is defined by

$$i_g(D) = \max\{\max(d^+(x), d^-(x)) - \min(d^+(y), d^-(y)) \mid x, y \in V(D)\},\$$

and the local irregularity by

$$i_l(D) = \max_{x \in V(D)} |d^+(x) - d^-(x)|.$$

If  $i_g(D) = 0$ , then D is regular and if  $i_g(D) \le 1$ , then D is called almost regular.

A cycle of length m is an m-cycle. A cycle in a digraph D is Hamiltonian if it includes all the vertices of D. A cycle-factor of a digraph D is a spanning subdigraph of D consisting of disjoint cycles.

For a vertex set X of D, we define D[X] as the subdigraph induced by X. A set  $X \subseteq V(D)$  of vertices is *independent* if the induced subdigraph D[X] has no arcs. The *independence number*  $\alpha(D) = \alpha$  is the maximum size among the independent sets of vertices of D.

A digraph D is *strongly connected* or *strong* if, for each pair of vertices u and v, there is a path from u to v in D. A digraph D with at least k+1 vertices is k-connected if for any set A of at most k-1 vertices, the subdigraph D-A obtained by deleting A is strong. The *connectivity* of D, denoted by  $\kappa(D)$ , is then defined to be the largest value of k such that D is k-connected.

#### 2. Introduction and preliminary results

In 2002, Tewes et al. [4] have made the following observation.

**Lemma 2.1** (Tewes et al. [4]). If D is a c-partite tournament with the partite sets  $V_1, V_2, \ldots, V_c$ , then  $||V_i|| \leq |V_i|| \leq 2i_g(D)$  for  $1 \leq i \leq j \leq c$ .

The next result is a part of the main theorem in Yeo's paper [9].

**Theorem 2.2** (Yeo [9]). Let  $V_1, V_2, \ldots, V_c$  be the partite sets of a c-partite tournament D such that  $|V_1| \leq |V_2| \leq \cdots \leq |V_c|$ . If

$$i_g(D) \leqslant \frac{|V(D)| - |V_{c-1}| - 2|V_c| + 2}{2},$$

then D is Hamiltonian.

In the case that D is an almost regular c-partite tournament with the partite sets  $V_1, V_2, \ldots, V_c$  such that  $|V_1| \le |V_2| \le \cdots \le |V_c| = r$ , it follows from Lemma 2.1 that  $|V_1| \ge r - 2$ . If D is an almost regular c-partite tournament with  $c \ge 4$ , then Theorem 2.2 shows that D is Hamiltonian with exception of the case that  $1 = |V_1| = |V_2| < |V_3| = 2 < |V_4| = 3$ . In this case there really exists the following example which is not Hamiltonian.

**Example 2.3.** Let  $V_1 = \{x\}$ ,  $V_2 = \{y\}$ ,  $V_3 = \{u_1, u_2\}$ ,  $V_4 = \{v_1, v_2, v_3\}$  be the partite sets of a 4-partite tournament such that  $u_1 \to \{x, y\} \to u_2 \to V_4 \to u_1, x \to y \to \{v_2, v_3\} \to x \to v_1 \to y$ . Then it is straightforward to verify that D is almost regular without a Hamiltonian cycle.

We have shown above that almost all almost regular c-partite tournaments with  $c \ge 4$  are Hamiltonian. However, the next examples will show that this is no longer true for almost regular 3-partite tournaments.

**Example 2.4.** Let *D* be a 3-partite tournament with the partite sets  $V_1$ ,  $V_2$ ,  $V_3$  such that  $|V_1| = r - 1$  and  $|V_2| = |V_3| = r$  and  $V_1 \rightarrow V_2 \rightarrow V_3 \rightarrow V_1$ . Then *D* is almost regular without a Hamiltonian cycle.

The longest cycle in this infinite family of almost regular 3-partite tournaments has length |V(D)| - 2. In this paper we will prove that every vertex of an almost regular 3-partite tournament D is contained in a directed cycle of length at least |V(D)| - 2. In special cases we can even give better bounds.

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