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Bounds for separating hash families

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ABSTRACT

This paper aims to present new upper bounds on the size of separating hash families. These bounds improve previously known bounds for separating hash families.

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1. Introduction

Let h be a function from a set A to a set B and let $C_1, C_2, \dots, C_t \subseteq A$ be t pairwise disjoint subsets. We say that h separates C_1, C_2, \dots, C_t if $h(C_1), h(C_2), \dots, h(C_t)$ are pairwise disjoint. Let $|A| = n$ and $|B| = m$. We call a set \mathcal{H} of N functions from A to B an $(N; n, m)$ -hash family. We say that \mathcal{H} is an $(N; n, m, \{w_1, w_2, \dots, w_t\})$ separating hash family, and we shall also write as an SHF($N; n, m, \{w_1, w_2, \dots, w_t\}$), if for all pairwise disjoint subsets $C_1, C_2, \dots, C_t \subseteq A$ with $|C_i| = w_i$, for $i = 1, 2, \dots, t$, there exists at least one function $h \in \mathcal{H}$ that separates C_1, C_2, \dots, C_t . The multi-set $\{w_1, w_2, \dots, w_t\}$ is the type of the separating hash family. Obviously, we have $2 \leq t \leq m$ and $\sum_{i=1}^t w_i \leq n$. Separating hash family with $t = 2$ was introduced in [13] and the general case in [16]. It is worth remarking that various well-known combinatorial objects may be viewed as special cases of separating hash families. For example, if $w_1 = w_2 = \dots = w_t = 1$, an SHF($N; n, m, \{1, 1, \dots, 1\}$) is called a perfect hash family which is usually denoted by PHF($N; n, m, t$). Perfect hash families have been studied extensively, see for instance, [1,3,5,9,10,12,18]. A w -frameproof code is a separating hash family of type $\{1, w\}$ [4,6,11] and a w -secure frameproof code is a separating hash family of type $\{w, w\}$ [13]. Further, a w -IPP code (code with identifiable parent property) [7,11,17], is necessarily a PHF with $t = w + 1$ and an SHF of type $\{w, w\}$.

An SHF($N; n, m, \{w_1, w_2, \dots, w_t\}$) can be depicted as an $N \times n$ array \mathcal{A} in which the columns are labeled by the elements of A , the rows by the functions $h_i \in \mathcal{H}$ and the (i, j) -entry of the array is the

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value $h_i(j)$. Thus, an SHF($N; n, m, \{w_1, w_2, \dots, w_t\}$) is equivalent to an $N \times n$ array with entries from a set of m symbols such that for all disjoint sets of columns C_1, C_2, \dots, C_t of \mathcal{A} with $|C_i| = w_i$, for $i = 1, 2, \dots, t$, there exists at least one row r of \mathcal{A} such that

$$\{\mathcal{A}(r, x): x \in C_i\} \cap \{\mathcal{A}(r, y): y \in C_j\} = \emptyset,$$

for all $i \neq j$. We call \mathcal{A} the *array representation* or *matrix representation* of the hash family.

In general, for given $N, m, \{w_1, w_2, \dots, w_t\}$ we want to maximize n . The determination of bounds for n has been subject of much research recently [2,8,11,14–16].

The best known upper bounds on n for separating hash families of type $\{w_1, w_2\}$ are the following.

Theorem 1. (See [5,11].) *Suppose there exists an SHF($N; n, m, \{1, w\}$) with $w \geq 2$. Then $n \leq w(m^{\lceil \frac{N}{w} \rceil} - 1)$.*

Theorem 2. (See [16].) *Suppose there is an SHF($N; n, m, \{2, 2\}$). Then $n \leq 4m^{\lceil \frac{N}{3} \rceil} - 3$.*

For the special case $\{w_1, w_2, w_3\} = \{1, 1, 2\}$ we have the following strong bound.

Theorem 3. (See [16].) *Suppose there is an SHF($N; n, m, \{1, 1, 2\}$). Then $n \leq 3m^{\lceil \frac{N}{3} \rceil} + 2 - 2\sqrt{3m^{\lceil \frac{N}{3} \rceil} + 1}$.*

A general bound for SHF of type $\{w_1, \dots, w_t\}$ has been obtained by Stinson and Zaverucha in [14]. In [2] Blackburn, Etzion, Stinson and Zaverucha introduce a new method to establish a significant bound for SHF of type $\{w_1, \dots, w_t\}$, which considerably improves the bound in [14], when $w_i \geq 2$ for all $i = 1, \dots, t$. We record this bound for SHF of type $\{w_1, \dots, w_t\}$ in the following theorem.

Theorem 4. (See [2].) *Suppose an SHF($N; n, m, \{w_1, \dots, w_t\}$) exists. Let $u = \sum_{i=1}^t w_i$. Then*

$$n \leq \gamma m^{\lceil \frac{N}{(u-1)} \rceil},$$

where $\gamma = (w_1 w_2 + u - w_1 - w_2)$, and w_1 and w_2 are the smallest two of the integers w_i .

Note that the constant γ in Theorem 4 depends on w_1, w_2, \dots, w_t . If we take $\gamma = \binom{u}{2}$ for the theorem, we obtain a bound derived from the graph theoretical method [2], and if we take $\gamma = 2(u - w_1)w_1 - w_1$, where w_1 is the smallest of the integers w_i , we have the bound in [14].

It should be noted that there exist further bounds for type $\{w_1, w_2\}$ and for general type $\{w_1, w_2, \dots, w_t\}$ [14,15]. However as those bounds have been improved by the bound of Theorem 4, they are not included here.

To date, Theorem 4 presents the best known bound for SHF of general type $\{w_1, \dots, w_t\}$.

In this paper we present new strong bounds for SHF which improve the Blackburn–Etzion–Stinson–Zaverucha bound of Theorem 4.

2. Bounds for SHF of type $\{w_1, \dots, w_t\}$

We aim to prove the following results.

Theorem 5. *Suppose there exists an SHF($N; n, m, \{w_1, w_2\}$). Let $u = w_1 + w_2$. Then*

$$n \leq (u - 1)m^{\lceil \frac{N}{(u-1)} \rceil}.$$

Theorem 6. *Let $t \geq 3$ be an integer. Suppose there exists an SHF($N; n, m, \{w_1, w_2, \dots, w_t\}$). Let $u = \sum_{i=1}^t w_i$. Then*

$$n \leq (u - 1)(m - 1)^{\lceil \frac{N}{(u-1)} \rceil} + 1.$$

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