



# Bid coordination in split-award procurement: The buyer need not know anything



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## HIGHLIGHTS

- Anton and Yao (1989) show anti-competitive bid-coordination in split-award procurement.
- Anton–Yao's result assumes that the buyer knows suppliers' costs.
- The current analysis dispenses the knowledgeable buyer assumption.
- With the weaker assumption, Anton–Yao's result is re-established.

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## ABSTRACT

Anton and Yao (1989) show that in split-award procurement auctions bidders coordinate their bids to sustain high buyer price. We relax their assumption that the buyer has full information about the suppliers' production costs and restore the coordination outcome.

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## 1. Introduction

Anton and Yao (1989), hereafter AY, established the surprising result that under complete information about each other's costs if two bidders were allowed to bid for a continuum of splits of a given production requirement of a buyer, the bidders would coordinate their bids leading to a high price for the buyer.<sup>2</sup> Since for any split  $\alpha \in (0, 1)$ , with  $\alpha$  fraction of production awarded to a developer  $D$  and  $1 - \alpha$  fraction awarded to a second source  $S$ , either bidder can veto the split by submitting a high own bid; the main discipline on the equilibrium price and the viability of an interior split comes

from the bidders' sole-source bids. At an interior split a high overall price and individual bidder profits are maintained using sole-source profits as thresholds and threat points.

One notable aspect of AY's setup is the assumption that the buyer has full information about the suppliers' costs. This assumption plays an important role when a tie in minimal total bids occurs: among the tied splits the buyer should select one that involves the minimal production cost. But with such knowledge there is no reason for the buyer to hold an auction. Instead, he can make a take-it-or-leave-it joint offer of a price equal to the minimum total production cost which the suppliers cannot refuse, thus avoiding the coordination outcome. Furthermore, in practice, it is very unlikely for a buyer to be fully aware of the suppliers' costs.

In this note, we assume instead that the buyer has no information about the suppliers' costs. To accommodate this assumption, we use an intuitively plausible tie-breaking rule that works independently of the buyer's information. This tie-breaker first looks at all splits associated with the minimum total bid, and then picks

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<sup>2</sup> There are significant follow-on works based on Anton and Yao (1989). See, for example, Alcalde and Dahm (2013) and the references therein.

the split that is closest to the equal-share split. If this process results in two different splits equidistant from  $\alpha = 1/2$ , then the tie-breaker can favor either bidder, say bidder  $D$ , giving him the option to choose between the two splits and if he does not exercise his option then bidder  $S$  selects her desired split. We re-establish the bid coordination outcome under this weaker assumption.

## 2. Two bidders game

Two potential suppliers,  $D$  (developer) and  $S$  (second source), submit sealed bids for a continuum of splits,  $\alpha \in [0, 1]$ , of a total production contract  $x$ . A pair of bids  $(P_D(\alpha), P_S(\alpha))$  implies that at split  $\alpha$ ,  $D$  produces  $\alpha$  share for a payment of  $P_D(\alpha)$  while  $S$  produces  $1 - \alpha$  share for  $P_S(\alpha)$ , leading to a buyer price  $P_D(\alpha) + P_S(\alpha)$ . The bid functions are not required to be smooth. Let  $(C_D(\alpha), C_S(\alpha))$  be the respective production costs of  $D$  and  $S$  at split  $\alpha$  with  $C_D(0) = C_S(1) = 0$ . As in AY, there are no additional restrictions on the cost functions, and the suppliers are assumed to be fully informed about each other's costs when they bid. The total production cost at split  $\alpha$  is

$$B(\alpha) = C_D(\alpha) + C_S(\alpha).$$

The bidders' profits are given by

$$\Pi_i(\alpha) = P_i(\alpha) - C_i(\alpha), \quad i = D, S.$$

■ *The buyer's selection of the production split.* For any pair of bids, the buyer chooses a split  $\alpha$  to minimize its procurement cost:

$$\min_{\alpha \in [0,1]} G(\alpha) := P_D(\alpha) + P_S(\alpha).$$

If the solution is unique, the buyer chooses this production split. If the minimization yields more than one solution, a tie-breaking rule is needed to pick one split. AY assume that the buyer knows the cost-minimizing split and chooses that split when a tie occurs. As mentioned earlier, if a buyer has such information about the costs, there is no point in holding an auction—it can simply make a take-it-or-leave-it offer, minimizing its procurement cost and avoiding the coordination problem.

We assume, instead, that the buyer has no information about the suppliers' production costs. To accommodate this new assumption, we propose the following tie-breaking rule:

First determine  $\alpha$  value(s) closest to  $1/2$ .

1. If this  $\alpha$  value is unique, choose the corresponding production split.
2. If there are two  $\alpha$  values equidistant from  $1/2$ , let bidder  $D$  get the priority to declare his preference ordering over these two splits.

If  $D$  declares a strict preference for one  $\alpha$  over another, pick  $D$ 's preferred  $\alpha$  as the final split. If  $D$  expresses an indifference, then  $S$  gets to pick her preferred  $\alpha$  from the two values which then becomes the final split. If  $S$  is also indifferent then the buyer selects the higher of the two  $\alpha$ 's.||

### ■ Equilibrium analysis

**Lemma 1** (AY, 1989). Let  $(P_D^*, P_S^*)$  be a Nash equilibrium and  $g^*$  be the corresponding price to the buyer. Then,

$$g^* = P_D^*(1) = P_S^*(0).$$

**Lemma 1** defines the ceiling on the equilibrium price through sole-source bids. We omit the proof because it is the same as in AY.

**Lemma 2** (Production Costs). Suppose an inefficient split,  $\alpha^{\text{in}} \in [0, 1]$ , is supported in an equilibrium  $(P_D^*, P_S^*)$ . Then,

$$\min\{B(0), B(1)\} \geq B(\alpha^{\text{in}}). \quad (1)$$

**Proof.** By Lemma 1,

$$g^* = P_D^*(1) = P_S^*(0) = P_D^*(\alpha^{\text{in}}) + P_S^*(\alpha^{\text{in}}).$$

Without loss of generality, suppose  $B(1) \leq B(0)$ . Suppose contrary to (1),  $C_D(1) < C_D(\alpha^{\text{in}}) + C_S(\alpha^{\text{in}})$ . Then

$$0 \leq \Pi_D^*(\alpha^{\text{in}}) \leq g^* - [C_D(\alpha^{\text{in}}) + C_S(\alpha^{\text{in}})] < g^* - C_D(1),$$

where  $\Pi_D^*(\alpha^{\text{in}})$  is  $D$ 's profit in the posited equilibrium involving  $\alpha^{\text{in}}$ -split. But then  $D$  can lower his bid slightly below  $g^*$  at the sole-source and realize a profit arbitrarily close to  $g^* - C_D(1)$  that exceeds  $\Pi_D^*(\alpha^{\text{in}})$ , contradicting that  $\alpha^{\text{in}}$ -split is an equilibrium outcome. Hence, (1) must hold. □

**Lemma 2** implies that no strictly inefficient split can be supported in a Nash equilibrium if sole-source production is cost efficient.

**Proposition 1** (Equilibrium Characterization). Bidding strategies  $(P_D^*, P_S^*)$  constitute a Nash equilibrium resulting in an equilibrium split  $\alpha^*$  if and only if the following complete set of conditions under [1]–[3] are satisfied:

1. Price ceiling condition:

$$g^* = P_D^*(1) = P_S^*(0). \quad (2)$$

2. No profitable deviation in bidding: Neither bidder finds it profitable to deviate unilaterally to an alternative bidding strategy, i.e.,

$$\Pi_i^*(\alpha^*) + B(\alpha^*) \leq \Pi_i^*(\alpha) + B(\alpha) \quad \text{for all } \alpha \in [0, 1], i = D, S. \quad (3)$$

3. Picking the winning split  $\alpha^*$  using the buyer's selection rule and the tie-breaker, given submitted bids  $(P_D^*, P_S^*)$ :

$$(i) \text{ If } |\alpha - \frac{1}{2}| < |\alpha^* - \frac{1}{2}|, \text{ then} \\ g^* < P_D^*(\alpha) + P_S^*(\alpha); \quad (4)$$

$$(ii) \text{ If } |\alpha^* - \frac{1}{2}| < |\alpha - \frac{1}{2}|, \text{ then} \\ g^* \leq P_D^*(\alpha) + P_S^*(\alpha); \quad (5)$$

$$(iii) \text{ If } |\alpha^* - \frac{1}{2}| = |\alpha - \frac{1}{2}|, \text{ then} \\ \text{– either (a) :} \\ g^* < P_D^*(\alpha) + P_S^*(\alpha), \quad (6)$$

$$\text{– or (b) :} \\ g^* = P_D^*(\alpha) + P_S^*(\alpha), \quad \text{and} \quad (7)$$

$$\left\{ \begin{array}{l} \Pi_D^*(\alpha^*) > \Pi_D^*(\alpha); \\ \text{or } \Pi_D^*(\alpha^*) = \Pi_D^*(\alpha) \\ \text{and } \Pi_S^*(\alpha^*) > \Pi_S^*(\alpha); \\ \text{or } \Pi_D^*(\alpha^*) = \Pi_D^*(\alpha) \\ \Pi_S^*(\alpha^*) = \Pi_S^*(\alpha) \\ \text{and } \alpha^* > \alpha. \end{array} \right. \quad (8)$$

**Proof.** (Necessity) The necessity of item [1] follows from Lemma 1. The derivation of condition (3) in item [2] is exactly the same as in AY.

To verify the necessity of item [3], first observe that  $\alpha^*$  being the winner, it must pick itself when faced with all alternative values  $\alpha \neq \alpha^*$ . The conditions are exhaustively listed by partitioning the range of production splits  $[0, 1]$ . In the range under (i), if condition (4) fails for some  $\alpha$  then the tie-breaker would discard  $\alpha^*$  as the winner, so (4) must hold. For  $\alpha$  in the range listed under (ii), even if the overall bid price equals  $g^*$  the tie-breaker will pick  $\alpha^*$ , implying condition (5). For the unique  $\alpha$  under (iii), either the overall price must be higher than  $g^*$  implying (6), or in the case of a tie between  $\alpha^*$  and  $\alpha$  the second tie-breaking provision is implemented implying conditions (7) and (8).

(Sufficiency) The proof is straightforward and omitted. □

**Proposition 2** (Sole-Source Outcome).

- (i) If  $B(0) < B(\alpha)$  for all  $\alpha \in (0, 1]$ , the sole-source contract awarded to the cost-efficient supplier  $S$  is the unique Nash

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