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## Planar graphs without intersecting 5-cycles are 4-choosable

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#### a r t i c l e i n f o

### A B S T R A C T

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A graph *G* is *k*-choosable if it can be properly colored whenever every vertex has a list of at least *k* available colors. In the paper, it is proved that all planar graphs without intersecting 5-cycles are 4-choosable.

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### **1. Introduction**

All graphs considered in this paper are simple, finite and undirected, and we follow  $[2]$  for the terminologies and notations not defined here. Let *G* be a graph with the vertex set  $V(G)$  and the edge set  $E(G)$ . For a vertex  $v \in V$ , let  $N(v)$  denote the set of vertices adjacent to v, and let *d*(v) = |*N*(v)| denote the degree of v. A *k*-*vertex*, a *k* <sup>+</sup>-*vertex* or a *k* <sup>−</sup>-*vertex* is a vertex of degree *k* , at least *k* or at most *k* respectively. We use δ(*G*) to denote the minimum degree of *G*. A *k*-cycle is a cycle of length *k*, and a 3-cycle is usually called a *triangle*. Two cycles are *adjacent*( or *intersecting*) if they share at least one edge (or vertex, respectively).

A *proper k*-*coloring* of a graph *G* is a mapping  $\phi$  from *V*(*G*) to the color set [*k*] = {1, 2, ..., *k*} such that  $\phi(x) \neq \phi(y)$ for every two adjacent vertices *x* and *y* of *G*. We say that *L* is an *assignment* for the graph *G* if it assigns a list *L*(v) of possible colors to each vertex v of *G*. If *G* has a proper coloring  $\phi$  such that  $\phi(v) \in L(v)$  for any vertex v, then we say that *G* is *L*-colorable or  $\phi$  is an *L*-coloring of *G*. The graph *G* is *k*-*choosable* if it is *L*-colorable for every assignment *L* satisfying  $|L(v)| > k$  for any vertex v.

The concept of a list coloring was introduced by Vizing [\[10\]](#page--1-1) and Erdős, Rubin and Taylor [\[4\]](#page--1-2), respectively. Thomassen [\[9\]](#page--1-3) showed that every planar graph is 5-choosable. Examples of plane graphs which are not 4-choosable and plane graphs of girth 4 which are not 3-choosable were given by Voigt [\[11](#page--1-4)[,12\]](#page--1-5). Since every planar graph *G* without 3-cycles has a vertex of degree at most 3, it is 4-choosable. Wang and Lih [\[13\]](#page--1-6) extended the result to all planar graphs without intersecting 3-cycles. Lam, Xu and Liu [\[8\]](#page--1-7) proved that all planar graphs without 4-cycles are 4-choosable. Lam, Shiu and Xu [\[7\]](#page--1-8)(Wang and Lih [\[14\]](#page--1-9), respectively) proved that all planar graphs without 5-cycles are 4-choosable. Fijavž, Juvan, Mohar, Škrekovski [\[6\]](#page--1-10) proved that all planar graphs without 6-cycles are 4-choosable. Farzad [\[5\]](#page--1-11) proved that all planar graphs without 7-cycles are 4-choosable. Cheng, Chen and Wang [\[3\]](#page--1-12) proved that planar graphs without 4-cycles adjacent to triangles are 4-choosable. In the paper, we prove that all planar graphs without intersecting 5-cycles are 4-choosable.

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<span id="page-1-1"></span>

**Fig. 2.** An orientation of the configurations in Figure 1.

#### <span id="page-1-2"></span>**2. Main result and its proof**

First, we introduce some notations and definitions used in this section. Let *G* be a plane graph. We use *F* or *F* (*G*) to denote the face set of *G*. For  $f \in F(G)$ , we use  $V(f)$  to denote the set of vertices on the boundary of f and write  $f = [u_1u_2 \cdots u_n]$  if  $u_1, u_2, \ldots, u_n$  are the boundary vertices of f in a cyclic order. A face of G is said to be *incident* with all edges and vertices in its boundary. The *degree* of a face *f* , denoted by *dG*(*f* ), is the number of edges incident with it, where a cut edge is counted twice. A *k*-*face*, *k* <sup>−</sup>-*face* or a *k* <sup>+</sup>-*face* is a face of degree *k*, at most *k* or at least *k*, respectively. For convenience, a *k*-face  $f = [v_1v_2 \cdots v_k]$  is often said to be a  $(d(v_1), d(v_2), \ldots, d(v_k))$ -face. For a face f, let  $n_i(f)$  and  $n_{i^+}(f)$  denote the number of  $i$ -vertices and  $i^+$ -vertices incident with $f$ , respectively. Denote by  $f_d(v)$  and  $f_{d^+}(v)$  the number of  $d$ -faces and  $d^+$ -faces incident with  $v$ , respectively.

<span id="page-1-0"></span>**Theorem 1.** *All planar graphs without intersecting* 5*-cycles are* 4*-choosable.*

**Proof.** Suppose, to the contrary, that [Theorem 1](#page-1-0) is false. Let *G* be a counterexample to [Theorem 1](#page-1-0) with the fewest vertices. Then  $\delta(G) > 4$  (see [\[8\]](#page--1-7)).

First, we introduce a well-known theorem proved by Alon and Tarsi [\[1\]](#page--1-13). This intricate theorem reveals the connection between the list coloring of a graph *G* and its orientations. A digraph (directed graph) *D* is an ordered pair (*V*(*D*), *A*(*D*)) consisting of the vertex set  $V(D)$  and arc set  $A(D)$ . For any arc  $a = \langle u, v \rangle$ , we say that *u* is the tail of *a* and *v* its head. The indegree  $d_D^-(v)$  of a vertex v in *D* is the number of arcs with head v, and the outdegree  $d_D^+(v)$  of v is the number of arcs with tail v. A directed cycle is denoted by a cyclic sequence  $u_1u_2 \cdots u_ku_1$  in which each vertex dominates its successor. A subdigraph *H* of a directed graph *D* is called *Eulerian* if the indegree *d* − *H* (v) of every vertex v of *H* is equal to its outdegree *d* + *H* (v). Note that we do not assume that *H* is connected. *H* is *e*v*en* if it has an even number of edges. Otherwise it is *odd*. Let *EE*(*D*) and *EO*(*D*) denote the numbers of even and odd Eulerian subgraphs of *D*, respectively. (For convenience we agree that the empty subgraph is an even Eulerian subgraph).

<span id="page-1-3"></span>**Theorem 2.** [\[1\]](#page--1-13) Let D be a digraph. For each vertex  $v \in V(D)$ , let  $f(v)$  be a set of  $d_D^+(v)+1$  distinct integers, where  $d_D^+(v)$  is the *outdegree of* v. If  $E(D) \neq EO(D)$ , then there is a proper coloring c :  $V(D) \rightarrow \mathbb{Z}$  such that  $c(v)$  is in  $f(v)$  for every vertex v. That *is, if L is a list assignment such that*  $|L(v)| = d_D^+(v) + 1$  *for all vertices*  $v$  *in D, then D is L-colorable.* 

**Lemma 3.** *G contains no subgraph isomorphic to one of the configurations in [Fig.](#page-1-1)* 1*, where the vertices marked by* • *have degree of* 4 *in G and those vertices marked by* ◦ *have degree of* 4 *or* 5 *in G.*

**Proof.** Let *L* be an arbitrary list assignment of *G* such that each vertex is assigned precisely 4 colors. Suppose that *G* contains a subgraph *H* isomorphic to one of the configurations in [Fig. 1.](#page-1-1) By the minimality of *G*, the graph *G* − *V*(*H*) has an *L*-coloring  $\varphi$ . An orientation of *H* and available colors of every vertex are shown in [Fig. 2.](#page-1-2) Since  $EE(G_1) = 1 < EO(G_1) = 2$ ,  $EE(G_2) = 1 < EO(G_2) = 2$  and  $EE(G_3) = 5 > EO(G_3) = 4$ , any  $G_i(1 \le i \le 3)$  satisfies [Theorem 2.](#page-1-3) So  $\varphi$  can be extended to *G*.  $\blacksquare$ 

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