# Injective choosability of subcubic planar graphs with girth 6 

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#### Abstract

An injective coloring of a graph $G$ is an assignment of colors to the vertices of $G$ so that any two vertices with a common neighbor have distinct colors. A graph $G$ is injectively $k$-choosable if for any list assignment $L$, where $|L(v)| \geq k$ for all $v \in V(G), G$ has an injective $L$-coloring. Injective colorings have applications in the theory of error-correcting codes and are closely related to other notions of colorability. In this paper, we show that subcubic planar graphs with girth at least 6 are injectively 5 -choosable. This strengthens the result of Lužar, Škrekovski, and Tancer that subcubic planar graphs with girth at least 7 are injectively 5 -colorable. Our result also improves several other results in particular cases.


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## 1. Introduction

A proper coloring of a graph $G$ is an assignment of colors to the vertices of $G$ so that any neighboring vertices receive distinct colors. The chromatic number of $G, \chi(G)$, is the minimum number of colors needed for a proper coloring of $G$. An injective coloring of a graph $G$ is an assignment of colors to the vertices of $G$ so that any two vertices with a common neighbor receive distinct colors. The injective chromatic number, $\chi_{i}(G)$, is the minimum number of colors needed for an injective coloring of $G$. An injective coloring of $G$ is not necessarily a proper coloring of $G$. Define the neighboring graph $G^{(2)}$ by $V\left(G^{(2)}\right)=V(G)$ and $E\left(G^{(2)}\right)=\{u v: u$ and $v$ have a common neighbor in $G\}$. Note that $\chi_{i}(G)=\chi\left(G^{(2)}\right)$.

Injective colorings were first introduced by Hahn, Kratochvíl, Širáň, and Sotteau [14], where the authors showed injective colorings can be used in coding theory, by relating the injective chromatic number of the hypercube to the theory of errorcorrecting codes. The authors showed that for a graph $G$ with maximum degree $\Delta, \chi_{i}(G) \leq \Delta(\Delta-1)+1$. They also showed that computing the injective chromatic number is NP-complete and gave bounds and structural results for the injective chromatic numbers of graphs with special properties. It is easy to see that $\Delta(G) \leq \chi_{i}(G) \leq|V(G)|$.

For each $v \in V(G)$, let $L(v)$ be a set of colors assigned to $v$. Then $L=\{L(v) \mid v \in V(G)\}$ is a list assignment of G. Given a list assignment $L$ of $G$, an injective coloring $\varphi$ of $G$ is called an injective L-coloring of $G$ if $\varphi(v) \in L(v)$ for every $v \in V(G)$. A graph is injectively $k$-choosable if for any list assignment $L$, where $|L(v)| \geq k$ for all $v \in V(G), G$ has an injective $L$-coloring. The injective choosability number of $G$, denoted as $\chi_{i}^{\ell}(G)$, is the minimum $k$ needed such that $G$ is injectively $k$-choosable. It is clear that $\chi_{i}(G) \leq \chi_{i}^{\ell}(G)$.

Graphs with low injective chromatic numbers have been studied extensively. A number of authors have studied the injective chromatic number of graphs $G$ in relation to their maximum degree, $\Delta(G)$, or their maximum average degree, $\operatorname{mad}(G)=\max _{H \subseteq G}\{2|E(H)| /|V(H)|\}$, for instance $[3-5,16]$. As $\operatorname{mad}(G)<\frac{2 g(G)}{g(G)-2}$ for all planar graphs, we can compute a

[^0]Table 1
Known results on the injective chromatic number and injective list chromatic number. A 'Yes' in the 'Planar' column indicates that the result holds only for planar graphs, and a 'No' indicates that the result holds for both planar and non-planar graphs. $A^{*}$ in the ' $g(G)$ ' column indicates that the girth was obtained using the bound $\operatorname{mad}(G)<\frac{2 g(G)}{g(G)-2}$. Results in $[9,10]$ are stated for injective coloring. However, the same proofs work also for injective list coloring (See Refs. [2,3,6,7,9,10,12,13,15,16]).

| Bounds | Planar | $\Delta(G)$ | $\operatorname{mad}(\mathrm{G})$ | $g(G)$ | Authors |
| :--- | :--- | :--- | :--- | :--- | :--- |
| $\chi_{i}(G) \leq \Delta+1$ | Yes | $\geq 18$ |  | $\geq 6$ | Borodin and Ivanova [3] |
| $\chi_{i}(G) \leq \Delta+3$ | Yes |  |  | $\geq 6$ | Dong and Lin [12] |
| $\chi_{i}(G) \leq \Delta+3$ | No |  | $<\frac{14}{5}$ | $\geq 7^{*}$ | Doyon, Hahn, and Raspaud [13] |
| $\chi_{i}(G) \leq \Delta+4$ | No |  | $<3$ | $\geq 6^{*}$ | Doyon, Hahn, and Raspaud [13] |
| $\chi_{i}(G) \leq \Delta+8$ | No |  | $<\frac{10}{3}$ | $\geq 5^{*}$ | Doyon, Hahn, and Raspaud [13] |
| $\chi_{i}(G) \leq 5$ | Yes | $\leq 3$ |  | $\geq 7$ | Lužar, Skrekovski, and Tancer [16] |
| $\chi_{i}^{\ell}(G) \leq \Delta+1$ | No |  | $<\frac{5}{2}$ | $\geq 10^{*}$ | Cranston, Kim and Yu [9] |
| $\chi_{i}^{\ell}(G) \leq \Delta+1$ | Yes | $\geq 4$ |  | $\geq 9$ | Cranston, Kim and Yu [9] |
| $\chi_{i}^{\ell}(G)=\Delta$ | Yes | $\geq 4$ |  | $\geq 13$ | Cranston, Kim and Yu [9] |
| $\chi_{i}^{\ell}(G)=\Delta$ | No |  | $<\frac{42}{19}$ | $\geq 21^{*}$ | Cranston, Kim and Yu [9] |
| $\chi_{i}^{\ell}(G) \leq 5$ | No | $\geq 3$ | $<\frac{36}{13}$ | $\geq 8^{*}$ | Cranston, Kim and Yu [10] |
| $\chi_{i}^{\ell}(G) \leq \Delta+2$ | No | $\geq 4$ | $<\frac{14}{5}$ | $\geq 7^{*}$ | Cranston, Kim and Yu [10] |
| $\chi_{i}^{\ell}(G) \leq \Delta+1$ | Yes | $\geq 24$ |  | $\geq 6$ | Borodin and Ivanova [2] |
| $\chi_{i}^{\ell}(G) \leq \Delta+2$ | Yes | $\geq 12$ |  | $\geq 6$ | Li and Xu [15] |
| $\chi_{i}^{\ell}(G) \leq \Delta+2$ | Yes | $\geq 8$ |  | $\geq 6$ | Bu and Lu [6] |
| $\chi_{i}^{\ell}(G) \leq \Delta+3$ | Yes |  |  | $\geq 6$ | Chen and Wu [7] |
| $\chi_{i}^{\ell}(G) \leq \Delta+4$ | Yes | $\geq 30$ |  | $\geq 5$ | Li and Xu [15] |
| $\chi_{i}^{\ell}(G) \leq \Delta+5$ | Yes | $\geq 18$ |  | $\geq 5$ | Li and Xu [15] |
| $\chi_{i}^{\ell}(G) \leq \Delta+6$ | Yes | $\geq 14$ |  | $\geq 5$ | Li and Xu [15] |
| $\chi_{i}^{\ell}(G) \leq 5$ | Yes | $\leq 3$ |  | $\geq 6$ | This paper |

bound for the girth of $G, g(G)$, given mad $(G)$. Table 1 consists of results for the injective chromatic number and the injective choosability number of graphs which depend on planarity, the maximum degree, the maximum average degree, and the girth.

For a planar graph $G$ with girth at least 6 and any maximum degree $\Delta$, the best known result about the injective chromatic number is $\chi_{i}(G) \leq \Delta+3$ [12]. In this paper, we improve this result for the case $\Delta=3$. Moreover, we improve the result of Lužar, Škrekovski, and Tancer [16] by decreasing the girth condition and by changing to injective list coloring. We also improve the other two highlighted bounds in Table 1 in special cases.

## Theorem 1. Every planar graph $G$ with $\Delta(G) \leq 3$ and $g(G) \geq 6$ is injectively 5-choosable.

This theorem is a step towards the conjecture of Chen, Hahn, Raspaud, and Wang [8] that all planar subcubic graphs are injectively 5-colorable. In order to prove Theorem 1 we prove a slightly stronger result in Theorem 2. Let $G$ be a graph and let $L$ be a list assignment. A precolored path in $G$ is a path $P_{k}$ on $k$ vertices where $|L(v)|=1$ for all $v \in V\left(P_{k}\right)$ and there is at most one vertex $v \in V\left(P_{k}\right)$ with a neighbor in $G-P_{k}$. Moreover, $\operatorname{deg}_{P_{k}}(v)$ is maximal among the other vertices in $P_{k}$ and $v$ has at most one neighbor in $G-P_{k}$. Vertices with lists of size one are called precolored. The set of all precolored vertices $\mathcal{P}$ in $G$ is proper if the lists of precolored vertices give a proper coloring of $G^{(2)}$ when restricted to $\mathcal{P}$. That is, the precolored vertices do not conflict among themselves.

Theorem 2. Let $G$ be a plane graph with $\Delta(G) \leq 3$ and $g(G) \geq 6$. Let $\mathcal{P} \subseteq V(G)$. Let $L$ be a list assignment of $G$ such that $|L(v)| \geq 5$ for $v \in V(G) \backslash \mathcal{P}$ and $|L(v)|=1$ for $v \in \mathcal{P}$. If the precolored vertices are proper, are all in the same face, form at most two precolored paths, each of which is on at most three vertices, then $G$ is injectively L-colorable.

## 2. Preliminaries

The following notation shall be used in the sequel. A $k$-vertex is a vertex of degree $k$. We denote the degree of a vertex $v$ by $\operatorname{deg}(v)$. We denote the set of neighbors of $v$ by $N(v)$ and $N(v) \cup\{v\}$ by $N[v]$. If we want to stress that the degree or neighborhood is in a particular graph $G$, we use subscript $G$, e.g. $\operatorname{deg}_{G}(v)$. A cut edge or bridge is an edge which, when

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