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# Partitioning a triangle-free planar graph into a forest and a forest of bounded degree

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## ABSTRACT

An  $(\mathcal{F}, \mathcal{F}_d)$ -partition of a graph is a vertex-partition into two sets  $F$  and  $F_d$  such that the graph induced by  $F$  is a forest and the one induced by  $F_d$  is a forest with maximum degree at most  $d$ . We prove that every triangle-free planar graph admits an  $(\mathcal{F}, \mathcal{F}_5)$ -partition. Moreover we show that if for some integer  $d$  there exists a triangle-free planar graph that does not admit an  $(\mathcal{F}, \mathcal{F}_d)$ -partition, then it is an NP-complete problem to decide whether a triangle-free planar graph admits such a partition.

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## 1. Introduction

We only consider finite simple graphs, with neither loops nor multiple edges. Planar graphs we consider are supposed to be embedded in the plane. Consider  $i$  classes of graphs  $\mathcal{G}_1, \dots, \mathcal{G}_i$ . A  $(\mathcal{G}_1, \dots, \mathcal{G}_i)$ -partition of a graph  $G$  is a vertex-partition into  $i$  sets  $V_1, \dots, V_i$  such that, for all  $1 \leq j \leq i$ , the graph  $G[V_j]$  induced by  $V_j$  belongs to  $\mathcal{G}_j$ . In the following we will consider the following classes of graphs:

- $\mathcal{F}$  the class of forests,
- $\mathcal{F}_d$  the class of forests with maximum degree at most  $d$ ,
- $\mathcal{D}_d$  the class of  $d$ -degenerate graphs (recall that a  $d$ -degenerate graph is a graph such that all subgraphs have a vertex of degree at most  $d$ ),
- $\Delta_d$  the class of graphs with maximum degree at most  $d$ ,
- $\mathcal{I}$  the class of empty graphs (i.e. graphs with no edges).

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**Table 1**  
Known results.

Classes	Vertex-partitions	References
Planar graphs	$(\mathcal{I}, \mathcal{I}, \mathcal{I}, \mathcal{I})$	The Four Color Theorem [1,2]
	$(\mathcal{I}, \mathcal{F}, \mathcal{F})$	Borodin [3]
	$(\mathcal{F}_2, \mathcal{F}_2, \mathcal{F}_2)$	Poh [8]
	$(\mathcal{F}, \mathcal{D}_2)$	Thomassen [10]
	$(\mathcal{I}, \mathcal{D}_3)$	Thomassen [11]
Planar graphs with girth 4	$(\mathcal{I}, \mathcal{I}, \mathcal{I})$	Grötzsch [6]
	$(\mathcal{F}, \mathcal{F})$	Folklore
	$(\mathcal{F}_5, \mathcal{F})$	Present paper (Theorem 3)
	$(\mathcal{I}, \mathcal{F})$	Open question (Question 1)
Planar graphs with girth 5	$(\mathcal{I}, \mathcal{F})$	Borodin and Glebov [4]

For example, an  $(\mathcal{I}, \mathcal{F}, \mathcal{D}_2)$ -partition of  $G$  is a vertex-partition into three sets  $V_1, V_2, V_3$  such that  $G[V_1]$  is an empty graph,  $G[V_2]$  is a forest, and  $G[V_3]$  is a 2-degenerate graph.

The Four Colour Theorem [1,2] states that every planar graph  $G$  admits a proper 4-colouring, that is  $G$  can be partitioned into four empty graphs, i.e.  $G$  has an  $(\mathcal{I}, \mathcal{I}, \mathcal{I}, \mathcal{I})$ -partition. Borodin [3] proved that every planar graph admits an acyclic colouring with at most five colours (an acyclic colouring is a proper colouring in which every two colour classes induce a forest). This implies that every planar graph admits an  $(\mathcal{I}, \mathcal{F}, \mathcal{F})$ -partition. Poh [8] proved that every planar graph admits an  $(\mathcal{F}_2, \mathcal{F}_2, \mathcal{F}_2)$ -partition. Thomassen proved that every planar graph admits an  $(\mathcal{F}, \mathcal{D}_2)$ -partition [10], and an  $(\mathcal{I}, \mathcal{D}_3)$ -partition [11]. However, there are planar graphs that do not admit any  $(\mathcal{F}, \mathcal{F})$ -partition [5]. Borodin and Glebov [4] proved that every planar graph of girth at least 5 (that is every planar graph with no triangles nor cycles of length 4) admits an  $(\mathcal{I}, \mathcal{F})$ -partition.

We focus on triangle-free planar graphs. Raspaud and Wang [9] proved that every planar graph with no triangles at distance at most 2 (and thus in particular every triangle-free planar graph) admits an  $(\mathcal{F}, \mathcal{F})$ -partition. However, it is not known whether every triangle-free planar graph admits an  $(\mathcal{I}, \mathcal{F})$ -partition. We pose the following questions:

**Question 1.** Does every triangle-free planar graph admit an  $(\mathcal{I}, \mathcal{F})$ -partition?

**Question 2.** More generally, what is the lowest  $d$  such that every triangle-free planar graph admits an  $(\mathcal{F}, \mathcal{F}_d)$ -partition?

Note that proving  $d = 0$  in Question 2 would prove Question 1. The main result of this paper is the following:

**Theorem 3.** Every triangle-free planar graph admits an  $(\mathcal{F}, \mathcal{F}_5)$ -partition.

This implies that  $d \leq 5$  in Question 2. Our proof uses the discharging method. It is constructive and immediately yields an algorithm for finding an  $(\mathcal{F}, \mathcal{F}_5)$ -partition of a triangle-free planar graph in quadratic time.

Note that Montassier and Ochem [7] proved that not every triangle-free planar graph can be partitioned into two graphs of bounded degree (which shows that our result is tight in some sense).

Finally, we show that if for some  $d$ , there exists a triangle-free planar graph that does not admit an  $(\mathcal{F}, \mathcal{F}_d)$ -partition, then deciding whether a triangle-free planar graph admits such a partition is NP-complete. That is, if the answer to Question 2 is some  $k > 0$ , then for all  $0 \leq d < k$ , deciding whether a triangle-free planar graph admits an  $(\mathcal{F}, \mathcal{F}_d)$ -partition is NP-complete. We prove this by reduction to PLANAR 3-SAT.

All presented results on vertex-partition of planar graphs are summarized in Table 1.

Theorem 3 will be proved in Section 2. Section 3 is devoted to complexity results.

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