

Contents lists available at ScienceDirect

## **European Journal of Combinatorics**

journal homepage: www.elsevier.com/locate/ejc



# Packing and covering immersion-expansions of planar sub-cubic graphs\*



Archontia C. Giannopoulou<sup>a</sup>, O-joung Kwon<sup>a</sup>, Jean-Florent Raymond<sup>b,c</sup>, Dimitrios M. Thilikos<sup>c,d</sup>

- <sup>a</sup> Logic and Semantics, Technische Universität Berlin, Berlin, Germany
- <sup>b</sup> Institute of Computer Science, University of Warsaw, Poland
- <sup>c</sup> AlGCo project team, CNRS, LIRMM, Montpellier, France
- <sup>d</sup> Department of Mathematics, National and Kapodistrian University of Athens, Greece

#### ARTICLE INFO

Article history: Received 1 August 2016 Accepted 24 May 2017

#### ABSTRACT

A graph H is an immersion of a graph G if H can be obtained by some subgraph G after lifting incident edges. We prove that there is a polynomial function  $f: \mathbb{N} \times \mathbb{N} \to \mathbb{N}$ , such that if H is a connected planar sub-cubic graph on h > 0 edges, G is a graph, and k is a non-negative integer, then either G contains k vertex/edge-disjoint subgraphs, each containing H as an immersion, or G contains a set F of f(k,h) vertices/edges such that  $G \setminus F$  does not contain H as an immersion.

© 2017 Elsevier Ltd. All rights reserved.

#### 1. Introduction

All graphs in this paper are finite, undirected, loopless, and may have multi-edges. Whenever we call a graph *simple* we also assume that it does not have multi-edges. Let  $\mathcal{C}$  be a class of graphs. A  $\mathcal{C}$ -vertex/edge cover of G is a set S of vertices/edges such that each subgraph of G that is isomorphic to a graph in G contains some element of G. A G-vertex/edge packing of G is a collection of vertex/edge-disjoint subgraphs of G, each isomorphic to some graph in G.

 $\label{lem:email$ 

An extended abstract of this paper appeared in the proceedings of WG 2016. Archontia C. Giannopoulou and O-joung Kwon have been supported by the European Research Council (ERC) under the European Union's Horizon 2020 research and innovation programme (ERC consolidator grant DISTRUCT, agreement No 648527). Part of the research of Archontia C. Giannopoulou took place while she was a post-doc at the Warsaw Center of Mathematics and Computer Science. O-joung Kwon has been supported by ERC Starting Grant PARAMTIGHT (No. 280152). Jean-Florent Raymond has been supported by the (Polish) National Science Centre grant PRELUDIUM 2013/11/N/ST6/02706. Dimitrios M. Thilikos has been supported by the ANR project DEMOGRAPH (ANR-16-CE40-0028).

We say that a graph class  $\mathcal C$  has the  $\operatorname{vertex/edge} \operatorname{Erdős-Pósa} \operatorname{property}$  (shortly  $\operatorname{v/e-E} \operatorname{\&P} \operatorname{property}$ ) for some graph class  $\mathcal G$  if there is a function  $f:\mathbb N\to\mathbb N$ , called a  $\operatorname{gap} \operatorname{function}$ , such that, for every graph G in  $\mathcal G$  and every non-negative integer k, either G has a  $\operatorname{vertex/edge} \mathcal C$ -packing of size k or G has a  $\operatorname{vertex/edge} \mathcal C$ -cover of size f(k). In the case where  $\mathcal G$  is the class of all graphs we simply say that  $\mathcal C$  has the  $\operatorname{v/e-E} \operatorname{\&P} \operatorname{property}$ . An interesting topic in Graph Theory, related to the notion of duality between graph parameters, is to detect instantiations of  $\mathcal C$  and  $\mathcal G$  such that  $\mathcal C$  has the  $\operatorname{v/e-E} \operatorname{\&P} \operatorname{property}$  for  $\mathcal G$  and, optimize the corresponding gap. Certainly, the first result of this type was the celebrated result of Erdős and Pósa in [11] who proved that the class of all cycles has the  $\operatorname{v-E} \operatorname{\&P} \operatorname{property}$  with gap function  $O(k \cdot \log k)$ . This result has triggered a lot of research on its possible extensions. One of the most general ones was given in [24] where it was proven that the class of graphs that are contractible to some graph H has the  $\operatorname{v-E} \operatorname{\&P} \operatorname{property}$  iff H is planar (see also [4,5,8] for improvements on the gap function).

Other instantiations of  $\mathcal C$  for which the v-E &P property has been proved concern odd cycles [18,21], long cycles [2], and graphs containing cliques as minors [9] (see also [14,16,23] for results on more general combinatorial structures).

As noticed in [8], cycles have the e-E &P property as well. Interestingly, only few more results exist for the cases where the e-E &P property is satisfied. It is known for instance that graphs contractible to  $\theta_r$  (i.e. the graph consisting of two vertices and an edge of multiplicity r between them) have the e-E &P property [3]. Moreover it was proven that odd cycles have the e-E &P property for planar graphs [19] and for 4-edge-connected graphs [18].

Given two graphs G and H, we say that H is an *immersion* of G if H can be obtained from some subgraph of G by lifting incident edges (see Section 2 for the definition of the lift operation). Given a graph H, we denote by  $\mathcal{I}(H)$  the set of all graphs that contain H as an immersion. Using this terminology, the edge variant of the original result of Erdős and Pósa in [11] implies that the class  $\mathcal{I}(\theta_2)$  has the v-E &P property (and, according to [8], the e-E &P property as well). A natural question is whether this can be extended for  $\mathcal{I}(H)$ , for other graphs H, different than  $\theta_2$ . This is the question that we consider in this paper. A distinct line of research is to identify the graph classes G such that for every graph H,  $\mathcal{I}(H)$  has the e-E &P property for G. In this direction, it was recently proved in [20] that for every graph H,  $\mathcal{I}(H)$  has the e-E &P property for 4-edge-connected graphs.

In this paper we show that if H is non-trivial (i.e., has at least one edge), connected, planar, and sub-cubic, i.e., each vertex is incident with at most 3 edges, then  $\mathcal{I}(H)$  has the v/e-E &P property (with polynomial gap in both cases). More concretely, our main result is the following.

**Theorem 1.** Let  $k \in \mathbb{N}$ . If H is a connected planar sub-cubic graph and G is a graph without any  $\mathcal{I}(H)$ -vertex/edge packing of size greater than k then G has a  $\mathcal{I}(H)$ -vertex/edge cover of size bounded by a polynomial function of h and k, where h = |E(H)|.

The main tools of our proof are the graph invariants of tree-cut width and tree-partition width, defined in [28] and [10] respectively (see Section 2 for the formal definitions). Our proof uses the fact that every graph of polynomially (on k) big tree-cut width contains a wall of height k as an immersion (as proved in [28]). This permits us to consider only graphs of bounded tree-cut width and, by applying suitable reductions, we finally reduce the problem to graphs of bounded tree-partition width (Theorem 2). The result follows as we next prove that for every H, the class  $\mathcal{I}(H)$  has the e-E &P property for graphs of bounded tree-partition width (Theorem 3).

One might conjecture that the result in Theorem 1 is tight in the sense that both being planar and sub-cubic are necessary for H in order for  $\mathcal{I}(H)$  to have the e-E &P property. In this direction, in Section 7, we give counterexamples for the cases where H is planar but not sub-cubic and is sub-cubic but not planar.

#### 2. Definitions and preliminary results

We use  $\mathbb{N}^+$  for the set of all positive integers and we set  $\mathbb{N} = \mathbb{N}^+ \cup \{0\}$ . Given a function  $f: A \to B$  and a set  $C \subseteq A$ , we denote by  $f|_C = \{(x, f(x)) \mid x \in C\}$ .

### Download English Version:

# https://daneshyari.com/en/article/5777407

Download Persian Version:

https://daneshyari.com/article/5777407

<u>Daneshyari.com</u>