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Journal of Combinatorial Theory,  
Series B

www.elsevier.com/locate/jctb

Layered separators in minor-closed graph classes  
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## ARTICLE INFO

## Article history:

Received 14 April 2015

Available online xxxx

## Keywords:

Separator

Planar graph

Surface

Minor

Topological minor

Layered separator

Layered treewidth

Nonrepetitive colouring

Queue layout

3-dimensional grid drawing

## ABSTRACT

Graph separators are a ubiquitous tool in graph theory and computer science. However, in some applications, their usefulness is limited by the fact that the separator can be as large as  $\Omega(\sqrt{n})$  in graphs with  $n$  vertices. This is the case for planar graphs, and more generally, for proper minor-closed classes. We study a special type of graph separator, called a *layered separator*, which may have linear size in  $n$ , but has bounded size with respect to a different measure, called the *width*. We prove, for example, that planar graphs and graphs of bounded Euler genus admit layered separators of bounded width. More generally, we characterise the minor-closed classes that admit layered separators of bounded width as those that exclude a fixed apex graph as a minor.

We use layered separators to prove  $\mathcal{O}(\log n)$  bounds for a number of problems where  $\mathcal{O}(\sqrt{n})$  was a long-standing previous best bound. This includes the *nonrepetitive chromatic number* and *queue-number* of graphs with bounded Euler genus. We extend these results with a  $\mathcal{O}(\log n)$  bound on the nonrepetitive chromatic number of graphs excluding a fixed topological minor, and a  $\log^{\mathcal{O}(1)} n$  bound on the queue-

☆ A short version of this paper and reference [22] was presented at the 54th Annual IEEE Symposium on Foundations of Computer Science (FOCS '13).

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<sup>1</sup> Research supported by NSERC and the Ontario Ministry of Research and Innovation.

<sup>2</sup> Research supported by NSERC.

<sup>3</sup> Research supported by the Australian Research Council.

<http://dx.doi.org/10.1016/j.jctb.2017.05.006>

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number of graphs excluding a fixed minor. Only for planar graphs were  $\log^{\mathcal{O}(1)} n$  bounds previously known. Our results imply that every  $n$ -vertex graph excluding a fixed minor has a 3-dimensional grid drawing with  $n \log^{\mathcal{O}(1)} n$  volume, whereas the previous best bound was  $\mathcal{O}(n^{3/2})$ .

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## 1. Introduction

Graph separators are a ubiquitous tool in graph theory and computer science since they are key to many divide-and-conquer and dynamic programming algorithms. Typically, the smaller the separator the better the results obtained. For instance, many problems that are  $\mathcal{NP}$ -complete for general graphs have polynomial time solutions for classes of graphs that have bounded size separators—that is, graphs of bounded treewidth.

By the classical result of Lipton and Tarjan [53], every  $n$ -vertex planar graph has a separator of size  $\mathcal{O}(\sqrt{n})$ . More generally, the same is true for every proper minor-closed graph class,<sup>4</sup> as proved by Alon et al. [3]. While these results have found widespread use, separators of size  $\Theta(\sqrt{n})$ , or non-constant separators in general, are not small enough to be useful in some applications.

<sup>4</sup> A graph  $H$  is a *topological minor* of a graph  $G$  if a subdivision of  $H$  is a subgraph of  $G$ . A graph  $H$  is a *minor* of a graph  $G$  if a graph isomorphic to  $H$  can be obtained from a subgraph of  $G$  by contracting edges. A class  $\mathcal{G}$  of graphs is *minor-closed* if  $H \in \mathcal{G}$  for every minor  $H$  of  $G$  for every graph  $G \in \mathcal{G}$ . A minor-closed class is *proper* if it is not the class of all graphs.

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