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5-choosability of graphs with crossings far apart [☆]Zdeněk Dvořák ^{a,1}, Bernard Lidický ^{b,2}, Bojan Mohar ^{c,3,4}^a Computer Science Institute of Charles University, Prague, Czechia^b Iowa State University, Ames IA, USA^c Department of Mathematics, Simon Fraser University, Burnaby, B.C. V5A 1S6, Canada

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ABSTRACT

We give a new proof of the fact that every planar graph is 5-choosable, and use it to show that every graph drawn in the plane so that the distance between every pair of crossings is at least 15 is 5-choosable. At the same time we may allow some vertices to have lists of size four only, as long as they are far apart and far from the crossings.

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Thomassen [5] showed that every planar graph is 5-choosable. We prove a strengthening of this result, allowing some crossings. Suppose that a graph G is drawn in the plane with some crossings, and for $i \in \{1, 2\}$, let $u_i v_i$ and $x_i y_i$ be two edges of G crossing each other. The *distance* between the crossing of $u_1 v_1$ with $x_1 y_1$ and the crossing of $u_2 v_2$ with

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x_2y_2 is the length of the shortest path in G with one end in the set $\{u_1, v_1, x_1, y_1\}$ and the other end in the set $\{u_2, v_2, x_2, y_2\}$. Our main result is the following.

Theorem 1. *Every graph drawn in the plane so that the distance between every pair of crossings is at least 15 is 5-choosable.*

Let us recall that a *list assignment* L for G is a function that assigns to each vertex of G a set $L(v)$, called the *list of admissible colors* for v . An L -coloring is a choice of a color $\varphi(v) \in L(v)$ for each $v \in V(G)$ such that no two adjacent vertices receive the same color. The graph is k -choosable if it admits an L -coloring for every list assignment L with $|L(v)| \geq k$ for every $v \in V(G)$.

The main idea of Thomassen's beautiful proof of 5-choosability of planar graphs is to establish the following more general statement.

Theorem 2 (Thomassen [5]). *Let G be a plane graph with the outer face F , xy an edge of F , and L a list assignment such that $|L(v)| \geq 5$ for $v \in V(G) \setminus V(F)$, $|L(v)| \geq 3$ for $v \in V(F) \setminus \{x, y\}$, $|L(x)| = |L(y)| = 1$ and $L(x) \neq L(y)$. Then G is L -colorable.*

Let us note that the lists of x and y of size 1 give a precoloring of a path of length 1 in the outer face of G . Naturally, one might try to prove [Theorem 1](#) by showing a variant of [Theorem 2](#) allowing distant crossings. Unfortunately, almost any attempt to alter the statement of [Theorem 2](#) (e.g., by allowing more than two vertices to be precolored, allowing lists of size 2 subject to some constraints, allowing some crossings in the drawing, etc.) fails with infinitely many counterexamples. To overcome this obstacle, we give a new proof of 5-choosability of planar graphs, see [Theorem 6](#) in [Section 1](#), which turns out to be more robust with respect to some strengthenings of the planar 5-choosability theorem. The proof of [Theorem 6](#) is inspired by Thomassen's proof [\[6\]](#) of 3-choosability of planar graphs of girth 5.

In the course of the proof of [Theorem 1](#), it is convenient to allow in addition to crossings also some other irregularities, such as vertices with fewer than 5 available colors. Hence, we actually obtain the following stronger statement (the distance between a vertex z and a crossing of edges uv and xy is the length of the shortest path from z to $\{u, v, x, y\}$).

Theorem 3. *Let G be a graph drawn in the plane with some crossings and let $N \subseteq V(G)$ be a set of vertices such that the distance between any pair of crossed edges is at least 15, the distance between any crossing and a vertex in N is at least 13, and the distance between any two vertices in N is at least 11. Then G is L -colorable for any list assignment L such that $|L(v)| = 4$ for $v \in N$ and $|L(v)| \geq 5$ for $v \in V(G) \setminus N$.*

Some distance condition on the crossings in [Theorem 1](#) is necessary, even if we allow only three crossings, as shown by K_6 . On the other hand, it was proved in [\[3\]](#) and independently also in [\[2\]](#) that the distance requirement is not needed, if we have at most two crossings. The inductive proof of [Theorem 3](#) involves a stronger inductive hypothesis

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