

## On palindromic factorization of words

### A.E. Frid<sup>a,\*,1</sup>, S. Puzynina<sup>a,b,2</sup>, L.Q. Zamboni<sup>c,b,3</sup>

<sup>a</sup> Sobolev Institute of Mathematics, 4 Koptyug av., 630090, Novosibirsk, Russia

<sup>b</sup> Department of Mathematics and Turku Centre for Computer Science, University of Turku, 20014 Turku, Finland

<sup>c</sup> Institut Camille Jordan, Université Claude Bernard Lyon 1, 43 boulevard du 11 novembre 1918, F69622 Villeurbanne Cedex, France

#### ARTICLE INFO

Article history: Received 12 November 2012 Accepted 4 January 2013 Available online 29 January 2013

MSC: 68R15

*Keywords:* Palindrome Periodicity of words Complexity of words

#### ABSTRACT

Given a finite word u, we define its *palindromic length*  $|u|_{pal}$  to be the least number n such that  $u = v_1 v_2 \dots v_n$  with each  $v_i$  a palindrome. We address the following open question: let P be a positive integer and w an infinite word such that  $|u|_{pal} \leq P$  for every factor u of w. Must w be ultimately periodic? We give a partial answer to this question by proving that for each positive integer k, the word w must contain a k-power, i.e., a factor of the form  $u^k$ . In particular, w cannot be a fixed point of a primitive morphism. We also prove more: for each pair of positive integers k and l, the word w must contain a position covered by at least l distinct k-powers. In particular, w cannot be a Sierpinski-like word.

© 2013 Elsevier Inc. All rights reserved.

#### 1. Introduction

Let *A* be a finite non-empty set, and let  $A^+$  denote the set of all finite non-empty words in *A*. A word  $u = u_1 u_2 \dots u_n \in A^+$  is called a *palindrome* if  $u_i = u_{n-i+1}$  for each  $i = 1, \dots, n-1$ . In particular each  $a \in A$  is a palindrome. We also regard the empty word as a palindrome.

Palindrome factors of finite or infinite words have been studied from different points of view. In particular, Droubay, Justin and Pirillo [4] proved that a word of length n can contain at most n + 1 distinct palindromes, which gave rise to the theory of *rich* words (see [5]). The number of palindromes of a given length occurring in an infinite word is called its *palindrome complexity* and is bounded by a function of its usual subword complexity [1]. However, in this paper we study palindromes in an infinite word from the point of view of decompositions.

\* Corresponding author.

0196-8858/\$ - see front matter © 2013 Elsevier Inc. All rights reserved. http://dx.doi.org/10.1016/j.aam.2013.01.002

E-mail addresses: anna.e.frid@gmail.com (A.E. Frid), svepuz@utu.fi (S. Puzynina), lupastis@gmail.com (L.Q. Zamboni).

<sup>&</sup>lt;sup>1</sup> Supported in part be RFBR grant 12-01-00089 and by the Presidential grant MK-4075.2012.1.

<sup>&</sup>lt;sup>2</sup> Supported in part by grant No. 251371 from the Academy of Finland and by RFBR grant 12-01-00448.

<sup>&</sup>lt;sup>3</sup> Supported in part by a FiDiPro grant from the Academy of Finland and by ANR grant SUBTILE.

For each word  $u \in A^+$  we define its *palindromic length*, denoted by  $|u|_{pal}$ , to be the least number *P* such that  $u = v_1 v_2 \dots v_P$  with each  $v_i$  a palindrome. As each letter is a palindrome, we have  $|u|_{pal} \leq |u|$ , where |u| denotes the length of *u*. For example,  $|01001010010|_{pal} = 1$  while  $|010011|_{pal} = 3$ . Note that 010011 may be expressed as a product of 3 palindromes in two different ways: (0)(1001)(1) and (010)(0)(11). In [10], O. Ravsky obtains an intriguing formula for the supremum of the palindromic lengths of all binary words of length *n*. The question considered in this paper is

**Question 1.** Do there exist an infinite non-ultimately periodic word *w* and a positive integer *P* such that  $|u|_{\text{pal}} \leq P$  for each factor *u* of *w*?

We conjecture that such aperiodic words do not exist, but at the moment we can prove it only partially. Namely, in this paper we prove that if such a word exists, then it is not *k*-power-free for any *k* and moreover, for all k > 1,  $l \ge 0$  it does not satisfy the (k, l)-condition defined in Section 4. A discussion what exactly the condition means and which class of words should be studied now to give a complete answer to the question is given in Section 5.

A preliminary version of this paper has been reported at Journées Montoises 2012.

#### 2. The case of *k*-power-free words

Let *k* be a positive integer. A word  $v \in A^+$  is called a *k*-power if  $v = u^k$  for some word  $u \in A^+$ . An infinite word  $w = w_1 w_2 \ldots \in A^{\mathbb{N}}$  is said to be *k*-power-free if no factor *u* of *w* is a *k*-power. For instance, the Thue–Morse word 0110100110010110... fixed by the morphism  $0 \mapsto 01$ ,  $1 \mapsto 10$  is 3-power-free (see for example [7]).

**Theorem 1.** Let k be a positive integer and  $w = w_1 w_2 \ldots \in A^{\mathbb{N}}$ . If w is k-power-free, then for each positive integer P there exists a prefix u of w with  $|u|_{pal} > P$ .

Recall that a word  $u_1 \dots u_n$  is called *t*-periodic if  $u_i = u_{i+t}$  for all *i* such that  $1 \le i \le n-t$ . The proof of Theorem 1 will make use of the following lemmas.

**Lemma 2.** Let u be a palindrome. Then for every palindromic proper prefix v of u, we have that u is (|u| - |v|)-periodic.

**Proof.** If *u* and *v* are palindromes with *v* a proper prefix of *u*, then *v* is also a suffix of *u* and hence *u* is (|u| - |v|)-periodic.  $\Box$ 

In what follows, the notation w[i..j] can mean the factor  $w_i w_{i+1} \dots w_j$  of a word  $w = w_1 \dots w_n \dots$  as well as its precise occurrence starting at the position numbered *i*; we always specify it when necessary.

**Lemma 3.** Suppose the infinite word w is k-power-free. If  $w[i_1..i_2]$  and  $w[i_1..i_3]$  are palindromes with  $i_3 > i_2$ , then

$$\frac{|w[i_1..i_3]|}{|w[i_1..i_2]|} > 1 + \frac{1}{k-1}.$$

**Proof.** By Lemma 2, the word  $w[i_1..i_3]$  is  $(i_3 - i_2)$ -periodic; at the same time, it cannot contain a k-power, so,  $|w[i_1..i_3]| < k(i_3 - i_2)$ . Thus,

$$\frac{|w[i_1..i_3]|}{|w[i_1..i_2]|} = \frac{|w[i_1..i_3]|}{|w[i_1..i_3]| - (i_3 - i_2)} > \frac{|w[i_1..i_3]|}{(1 - \frac{1}{k})(|w[i_1..i_3]|)} = 1 + \frac{1}{k - 1}.$$

Download English Version:

# https://daneshyari.com/en/article/6419680

Download Persian Version:

https://daneshyari.com/article/6419680

Daneshyari.com