# On palindromic factorization of words 

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#### Abstract

Given a finite word $u$, we define its palindromic length $|u|_{\mathrm{pal}}$ to be the least number $n$ such that $u=v_{1} v_{2} \ldots v_{n}$ with each $v_{i}$ a palindrome. We address the following open question: let $P$ be a positive integer and $w$ an infinite word such that $|u|_{\text {pal }} \leqslant P$ for every factor $u$ of $w$. Must $w$ be ultimately periodic? We give a partial answer to this question by proving that for each positive integer $k$, the word $w$ must contain a $k$-power, i.e., a factor of the form $u^{k}$. In particular, $w$ cannot be a fixed point of a primitive morphism. We also prove more: for each pair of positive integers $k$ and $l$, the word $w$ must contain a position covered by at least $l$ distinct $k$-powers. In particular, $w$ cannot be a Sierpinski-like word.


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## 1. Introduction

Let $A$ be a finite non-empty set, and let $A^{+}$denote the set of all finite non-empty words in $A$. A word $u=u_{1} u_{2} \ldots u_{n} \in A^{+}$is called a palindrome if $u_{i}=u_{n-i+1}$ for each $i=1, \ldots, n-1$. In particular each $a \in A$ is a palindrome. We also regard the empty word as a palindrome.

Palindrome factors of finite or infinite words have been studied from different points of view. In particular, Droubay, Justin and Pirillo [4] proved that a word of length $n$ can contain at most $n+1$ distinct palindromes, which gave rise to the theory of rich words (see [5]). The number of palindromes of a given length occurring in an infinite word is called its palindrome complexity and is bounded by a function of its usual subword complexity [1]. However, in this paper we study palindromes in an infinite word from the point of view of decompositions.

[^0]For each word $u \in A^{+}$we define its palindromic length, denoted by $|u|_{\text {pal }}$, to be the least number $P$ such that $u=v_{1} v_{2} \ldots v_{P}$ with each $v_{i}$ a palindrome. As each letter is a palindrome, we have $|u|_{\text {pal }} \leqslant|u|$, where $|u|$ denotes the length of $u$. For example, $|01001010010|_{\text {pal }}=1$ while $|010011|_{\text {pal }}=3$. Note that 010011 may be expressed as a product of 3 palindromes in two different ways: $(0)(1001)(1)$ and $(010)(0)(11)$. In [10], O. Ravsky obtains an intriguing formula for the supremum of the palindromic lengths of all binary words of length $n$. The question considered in this paper is

Question 1. Do there exist an infinite non-ultimately periodic word $w$ and a positive integer $P$ such that $|u|_{\text {pal }} \leqslant P$ for each factor $u$ of $w$ ?

We conjecture that such aperiodic words do not exist, but at the moment we can prove it only partially. Namely, in this paper we prove that if such a word exists, then it is not $k$-power-free for any $k$ and moreover, for all $k>1, l \geqslant 0$ it does not satisfy the ( $k, l$ )-condition defined in Section 4. A discussion what exactly the condition means and which class of words should be studied now to give a complete answer to the question is given in Section 5.

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## 2. The case of $\boldsymbol{k}$-power-free words

Let $k$ be a positive integer. A word $v \in A^{+}$is called a $k$-power if $v=u^{k}$ for some word $u \in A^{+}$. An infinite word $w=w_{1} w_{2} \ldots \in A^{\mathbb{N}}$ is said to be $k$-power-free if no factor $u$ of $w$ is a $k$-power. For instance, the Thue-Morse word $0110100110010110 \ldots$ fixed by the morphism $0 \mapsto 01,1 \mapsto 10$ is 3 -power-free (see for example [7]).

Theorem 1. Let $k$ be a positive integer and $w=w_{1} w_{2} \ldots \in A^{\mathbb{N}}$. If $w$ is $k$-power-free, then for each positive integer $P$ there exists a prefix $u$ of $w$ with $|u|_{\text {pal }}>P$.

Recall that a word $u_{1} \ldots u_{n}$ is called $t$-periodic if $u_{i}=u_{i+t}$ for all $i$ such that $1 \leqslant i \leqslant n-t$. The proof of Theorem 1 will make use of the following lemmas.

Lemma 2. Let $u$ be a palindrome. Then for every palindromic proper prefix $v$ of $u$, we have that $u$ is $(|u|-|v|)$ periodic.

Proof. If $u$ and $v$ are palindromes with $v$ a proper prefix of $u$, then $v$ is also a suffix of $u$ and hence $u$ is $(|u|-|v|)$-periodic.

In what follows, the notation $w[i . . j]$ can mean the factor $w_{i} w_{i+1} \ldots w_{j}$ of a word $w=$ $w_{1} \ldots w_{n} \ldots$ as well as its precise occurrence starting at the position numbered $i$; we always specify it when necessary.

Lemma 3. Suppose the infinite word $w$ is $k$-power-free. If $w\left[i_{1} . . i_{2}\right]$ and $w\left[i_{1} . . i_{3}\right]$ are palindromes with $i_{3}>i_{2}$, then

$$
\frac{\left|w\left[i_{1} . . i_{3}\right]\right|}{\left|w\left[i_{1} . . i_{2}\right]\right|}>1+\frac{1}{k-1} .
$$

Proof. By Lemma 2, the word $w\left[i_{1} . . i_{3}\right]$ is $\left(i_{3}-i_{2}\right)$-periodic; at the same time, it cannot contain a $k$-power, so, $\left|w\left[i_{1} . . i_{3}\right]\right|<k\left(i_{3}-i_{2}\right)$. Thus,

$$
\frac{\left|w\left[i_{1} . . i_{3}\right]\right|}{\left|w\left[i_{1} . . i_{2}\right]\right|}=\frac{\left|w\left[i_{1} . . i_{3}\right]\right|}{\left|w\left[i_{1} . . i_{3}\right]\right|-\left(i_{3}-i_{2}\right)}>\frac{\left|w\left[i_{1} . . i_{3}\right]\right|}{\left(1-\frac{1}{k}\right)\left(\left|w\left[i_{1} . . i_{3}\right]\right|\right)}=1+\frac{1}{k-1} .
$$

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