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# Universal solvability of interval max-plus matrix equations

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#### ABSTRACT

This paper deals with the solvability of interval matrix equations in max-plus algebra. Maxplus algebra is the algebraic structure in which classical addition and multiplication are replaced by  $\oplus$  and  $\otimes$ , where  $a \oplus b = \max\{a, b\}$  and  $a \otimes b = a + b$ .

The notation  $A \otimes X \otimes C = B$ , where A, B, and C are given interval matrices, represents an interval max-plus matrix equation. We define three types of solvability of interval max-plus matrix equations, namely the strong universal, universal, and weak universal solvability. We derive the necessary and sufficient conditions which can be verified in polynomial times

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### 1. Introduction

Behaviour of discrete event systems, in which the individual components move from event to event rather than varying continuously through time, is often described by systems of linear equations or by matrix equations. Discrete dynamic systems and related algebraic structures were studied using max-plus matrix operations in [2,3,16]. In the last decades, significant effort has been developed to study systems of max-plus linear equations in the form  $A \otimes x = b$ , where A is a matrix, b and x are vectors of compatible dimensions. Systems of linear equations over max-plus algebra are used in several branches of applied mathematics. Among interesting real-life applications let us mention e.g. a large scale model of Dutch railway network or synchronizing traffic lights in Delfts [13]. In the last two decades, interval systems of the form  $A \otimes x = b$  have been studied, for details see [2,4–6,9,8,12].

In this paper, we shall deal with interval matrix equations of the form  $A \otimes X \otimes C = B$ , where A, B, and C are given interval matrices of suitable sizes a X is an unknown matrix. In the following example we show one of the possible applications.

**Example 1.1.** Consider the manufacturing company that carries its three types of products in three places  $P_1$ ,  $P_2$ , and  $P_3$ . These products are placed in two warehouses  $W_1$  and  $W_2$ . Afterwards, they are loaded into two trucks  $T_1$  and  $T_2$  which deliver products to three shops  $S_1$ ,  $S_2$  and  $S_3$  (see Fig. 1).

In Fig. 1, there is an arrow from  $P_i$  to  $W_j$  if the products made in place  $P_i$  are stored in warehouse  $W_j$ . If products from warehouse  $W_j$  are loaded on a truck  $T_l$ , then there is an arrow from  $W_j$  to  $T_l$ . And, there is an arrow from  $T_l$  to  $S_k$ , if the truck  $T_l$  expedites products to shop  $S_k$  (i = 1, 2, 3; j = 1, 2; k = 1, 2, 3; l = 1, 2).

The symbols above or below the arrows,  $a_{ij}$  ( $c_{lk}$ ), express known times needed to transport goods from place  $P_i$  to warehouse  $W_j$  (from truck  $T_l$  to shop  $S_k$ ). The time needed to transport goods from warehouse  $W_j$  to truck  $T_k$  via  $W_j$  and  $T_l$  is denoted by  $x_{ik}$ . Then, total time needed to transport goods from place  $P_i$  to shop  $S_k$  is  $a_{ij} + x_{il} + c_{ik}$ .

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Fig. 1. Transport diagram.

Denote by  $b_{ik}$  the time scheduled to transport products from place  $P_i$  to shop  $S_k$ . To ensure the transportation for products made in  $P_1$  to shops  $S_1$ ,  $S_2$ , and  $S_3$ , the following equations must be satisfied:

$$\max\{a_{11}+x_{11}+c_{11},\ a_{11}+x_{12}+c_{21},\ a_{12}+x_{21}+c_{11},\ a_{12}+x_{22}+c_{21}\}=b_{11},\\ \max\{a_{11}+x_{11}+c_{12},\ a_{12}+x_{21}+c_{12}\}=b_{12},\\ \max\{a_{11}+x_{12}+c_{23},\ a_{12}+x_{22}+c_{23}\}=b_{13}.$$

Similar equalities must be satisfied to arrange the transportation for all products from  $P_2$  and  $P_3$  to shops  $S_1$ ,  $S_2$  and  $S_3$ .

In general, let us suppose that there are m places  $P_1, P_2, \ldots, P_m$ , n warehouses  $W_1, W_2, \ldots, W_n$ , s trucks  $T_1, T_2, \ldots, T_s$ , and r shops  $S_1, S_2, \ldots, S_r$ . If there is no connection from  $P_i$  to  $W_j$  (from  $T_l$  to  $S_k$ ), we put  $a_{ij} = -\infty$  ( $c_{lk} = -\infty$ ). Let us denote the following index sets  $M = \{1, 2, \ldots, m\}$ ,  $N = \{1, 2, \ldots, n\}$ ,  $N = \{1, 2, \ldots, n\}$ , and  $N = \{1, 2, \ldots, n\}$ .

We would like to determine times  $x_{jl}$  for any  $j \in N$  and for any  $l \in S$  such that the maximum of total transport time from  $P_i$  to  $S_k$  is equal to a given number  $b_{ik}$  for any  $i \in M$  and for any  $k \in R$ . It leads to solving the system of equations of the form

$$\max_{i \in N, l \in S} \{a_{ij} + x_{jl} + c_{lk}\} = b_{ik}. \tag{1}$$

#### 2. Preliminaries

Max-plus algebra is the triple  $(\overline{\mathbb{R}}, \oplus, \otimes)$ , where

$$\overline{\mathbb{R}} = \mathbb{R} \cup \{-\infty\}, \ a \oplus b = \max\{a, b\} \text{ and } a \otimes b = a + b.$$

The set of all  $m \times n$  matrices over  $\overline{\mathbb{R}}$  is denoted by  $\overline{\mathbb{R}}(m, n)$  and the set of all column n-vectors over  $\overline{\mathbb{R}}$  by  $\overline{\mathbb{R}}(n)$ .

Operations  $\oplus$  and  $\otimes$  are extended to matrices and vectors in the same way as in the classical algebra. We consider the *ordering* < on the sets  $\overline{\mathbb{R}}(m, n)$  and  $\overline{\mathbb{R}}(n)$  defined as follows:

- for  $A, C \in \overline{\mathbb{R}}(m, n)$ :  $A \leq C$  if  $a_{ij} \leq c_{ij}$  for each  $i \in M$  and for each  $j \in N$ ,
- for  $x, y \in \overline{\mathbb{R}}(n)$ :  $x \le y$  if  $x_i \le y_i$  for each  $j \in N$ .

We will use the *monotonicity of*  $\otimes$ , which means that for each  $A, C \in \overline{\mathbb{R}}(m, n)$  and for each  $B, D \in \overline{\mathbb{R}}(n, s)$  the implication

if 
$$A \leq C$$
 and  $B \leq D$  then  $A \otimes B \leq C \otimes D$ 

holds. Let  $A \in \overline{\mathbb{R}}(m,n)$  and  $b \in \overline{\mathbb{R}}(m)$ . We can write the system of max-plus linear equations in the matrix form

$$A \otimes x = b$$
. (2)

It is known (see [2,15]) that system (2) is solvable if and only if the vector  $x^*(A, b)$ , defined by

$$x_j^*(A, b) = \min_{i \in M} \{b_i - a_{ij}\}$$
 (3)

for any  $j \in N$ , where min  $\emptyset = I$ , is its solution. The vector  $x^*(A, b)$  is called a *principal solution* of system (2).

#### 3. Matrix equations

Let  $A=(a_{ij})\in \overline{\mathbb{R}}(m,n)$  and  $B=(b_{kl})\in \overline{\mathbb{R}}(r,s)$  be given. The *tensor product* of A and B is the following matrix of size  $mr\times ns$ :

$$A \boxtimes B = \begin{pmatrix} A \otimes b_{11} & A \otimes b_{12} & \dots & A \otimes b_{1s} \\ A \otimes b_{21} & A \otimes b_{22} & \dots & A \otimes b_{2s} \\ & \dots & & \dots & \dots \\ A \otimes b_{r1} & A \otimes b_{r2} & \dots & A \otimes b_{rs} \end{pmatrix}$$

2

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