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Barvinok's naive algorithm in Distance Geometry[☆]

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Abstract

In 1997, A. Barvinok gave a probabilistic algorithm to derive a near-feasible solution of a quadratically (equation) constrained problem from its semidefinite relaxation. We generalize this algorithm to handle matrix variables instead of vectors, and to two-sided inequalities instead of equations. We derive a heuristic for the distance geometry problem, and showcase its computational performance on a set of instances related to protein conformation.

Keywords: distance geometry, concentration of measure, protein structure

1. Introduction

We consider the following

Distance Geometry Problem (DGP). Given an integer $K > 0$ and a simple graph $G = (V, E, d)$ where $d : E \rightarrow \mathbb{R}_+$, decide whether there is a realization $x : V \rightarrow \mathbb{R}^K$ such that

$$\forall \{i, j\} \in E \quad \|x_i - x_j\|_2^2 = d_{ij}^2. \quad (1)$$

Let $n = |V|$ and $m = |E|$. The DGP is well known in the literature [1, 2], as it serves as a model for several applications (e.g. to sensor networks [3], molecular conformation [4, 5] and more [6]). A natural extension of the DGP is the interval version (denoted by *iDGP*), where $d : E \rightarrow \mathbb{IR}_+$ associates intervals (instead of scalars) to edges. This variant is the one often used in applications, since intervals naturally model data uncertainty and noise [7].

We adapt the *naive algorithm* proposed by A. Barvinok in [8, §5] to the DGP setting. Barvinok's main insight is that the solutions of a Semidefinite Programming (SDP) relaxation are "not too far" from the feasible set of the Quadratically Constrained Problem (QCP) which gives rise to the SDP relaxation: it suffices to factor the SDP solution and multiply it by a random vector having components sampled from a normal distribution. The very natural idea we propose is to use the naive algorithm first, then use this approximate solution as a starting point for a local Nonlinear Programming (NLP) solver deployed on the

original QCP, hoping it will converge to a realization satisfying the feasibility of Eq. (1).

Barvinok's naive algorithm is a randomized algorithm based on the concentration of measure phenomenon. It applies to SDP relaxations of systems of quadratic equations, i.e. pure feasibility, equation-only QCPs having vector solutions, or, equivalently, $n \times 1$ matrix solutions. This algorithm cannot be natively applied to the DGP, since a realization x is naturally represented by an $n \times K$ matrix $x = (x_{ik} \mid i \leq n, k \leq K)$ the i -th row of which is the position vector of vertex $i \in V$ in \mathbb{R}^K . This matrix may in general have rank greater than one. In [8], Barvinok gives a proof sketch which only applies to $n \times 1$ matrices. Of course, an $n \times K$ matrix can also be represented as an $n' \times 1$ matrix where $n' = nK$, but this would entail a QCP (and hence also an SDP relaxation) with $nK \times nK$ data matrices, which is practically prohibitive to solve. Moreover, the naive algorithm cannot be applied to the *iDGP* since it is constrained by quadratic inequalities rather than equations.

We make two contributions in this paper, one theoretical and the other computational. (i) We propose a generalization of Barvinok's result [8] in two directions: the $n \times K$ case, and the case of QCP inequalities of the form $d^L \leq x^T Q x \leq d^U$ (where x is $n \times K$). (ii) We establish the practical usefulness of Barvinok's result on a set of medium and large-scale DGP instances extracted from the Protein Data Bank (PDB) [9]. We also remark that our proof is detailed, and fills many gaps in Barvinok's "proof sketch" found in [8]. The rest of this paper is organized as follows. In Sect. 2 we present Barvinok's naive algorithm. In Sect. 3 we state and prove our generalization of Barvinok's concentration of measure result. In Sect. 5 we discuss our computational results.

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