



Optimality certificates for convex minimization and Helly numbers



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ABSTRACT

We consider the problem of minimizing a convex function over a subset of \mathbb{R}^n that is not necessarily convex (minimization of a convex function over the integer points in a polytope is a special case). We define a family of duals for this problem and show that, under some natural conditions, strong duality holds for a dual problem in this family that is more restrictive than previously considered duals.

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1. Introduction

Insights obtained through duality theory have spawned efficient optimization algorithms (combinatorial and numerical) which simultaneously work on a pair of primal and dual problems. Striking examples are Edmonds' seminal work in combinatorial optimization, and interior-point algorithms for numerical/continuous optimization.

Compared to duality theory for continuous optimization, duality theory for mixed-integer optimization is still underdeveloped. Although the linear case has been extensively studied, see, e.g., [4,5,13,14], nonlinear integer optimization duality was essentially unexplored until recently. An important step was taken by Morán et al. for conic mixed-integer problems [12], followed up by Baes et al. [2] who presented a duality theory for general convex mixed-integer problems. The approach taken by Moran et al. was essentially algebraic, drawing on the theory of subadditive functions. Baes et al. took a more geometric viewpoint and developed a duality theory based on lattice-free polyhedra. We follow the latter approach.

Given $S \subseteq \mathbb{R}^n$ and a convex function $f : \mathbb{R}^n \rightarrow \mathbb{R}$, we consider the problem

$$\inf_{s \in S} f(s). \quad (1)$$

In this paper, we restrict ourselves to the case that there exists an $x_0 \in \mathbb{R}^n$ such that $f(x_0) \leq f(s)$ for all $s \in S$. We describe a geometric

dual object that can be used to certify optimality of a solution to (1). To this end, fix such a point x_0 and define a closed set C to be an S -free neighborhood of x_0 if $x_0 \in \text{int}(C)$ and $\text{int}(C) \cap S = \emptyset$. Using the convexity of f , it follows that for any $\bar{s} \in S$ and any C that is an S -free neighborhood of x_0 , the following holds:

$$f(\bar{s}) \geq \inf_{z \in \text{bd}(C)} f(z) =: L(C), \quad (2)$$

where $\text{bd}(C)$ denotes the boundary of C (to see this, consider the line segment connecting \bar{s} and x_0 and a point at which this line segment intersects $\text{bd}(C)$). Thus, an S -free neighborhood of x_0 can be interpreted as a “dual object” that provides a lower bound of the type (2). As a consequence, the following is true.

Proposition 1. *If there exist $\bar{s} \in S$ and $C \subseteq \mathbb{R}^n$ that is an S -free neighborhood of x_0 , such that equality holds in (2), then \bar{s} is an optimal solution to (1).*

2. The dual problem

This motivates the definition of a dual optimization problem to (1). For any family \mathcal{F} of S -free neighborhoods of x_0 , define the \mathcal{F} -dual of (1) as

$$\sup_{C \in \mathcal{F}} L(C). \quad (3)$$

We say that strong duality holds with the \mathcal{F} -dual, if there exists $C^* \in \mathcal{F}$ such that $L(C^*) = \sup_{C \in \mathcal{F}} L(C) = \inf_{s \in S} f(s)$. For instance, if the hypothesis of Proposition 1 holds for some $C \in \mathcal{F}$, then

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one has strong duality with the \mathcal{F} -dual. Assuming very mild conditions on S and f (e.g., when S is a closed subset of \mathbb{R}^n disjoint from $\arg \inf_{x \in \mathbb{R}^n} f(x)$), it is straightforward to show that if \mathcal{F} is the family of all S -free neighborhoods of x_0 , then strong duality holds, i.e., there exist $\bar{s} \in S$ and $C \in \mathcal{F}$ such that the condition in Proposition 1 holds. However, the entire family of S -free neighborhoods is too unstructured to be useful as a dual problem. Moreover, the inner optimization problem (2) of minimizing on the boundary of C can be very hard if C has no structure other than being S -free. Thus, we would like to identify subfamilies \mathcal{F} of S -free neighborhoods that still maintain strong duality, while at the same time, are much easier to work with inside a primal–dual framework. We list below three subclasses that we expect to be useful in this line of research. First, we need the concept of a gradient polyhedron:

Definition 2. Given a set of points $z_1, \dots, z_k \in \mathbb{R}^n$,

$$Q := \{x \in \mathbb{R}^n : \langle a_i, x - z_i \rangle \leq 0, \quad i = 1, \dots, k\}$$

is said to be a gradient polyhedron of z_1, \dots, z_k if for every $i = 1, \dots, k$, $a_i \in \partial f(z_i)$, i.e., a_i is a subgradient of f at z_i .

Remark 3. For every gradient polyhedron Q of points z_1, \dots, z_k we have $L(Q) = \inf_{x \in \text{bd}(Q)} f(x) = \min_{i \in [k]} f(z_i)$.

Indeed, note that $a \in \partial f(z)$ means that $f(x) \geq f(z) + \langle a, x - z \rangle$ holds for all $x \in \mathbb{R}^n$.

Thus, for every $x \in \text{bd}(Q)$ we must have $\langle a_i, x - z_i \rangle = 0$ for some $i \in [k]$, which implies $f(x) \geq f(z_i)$.

We consider the following families.

- The family \mathcal{F}_{\max} of maximal convex S -free neighborhoods of x_0 , i.e., those S -free neighborhoods that are convex, and are not strictly contained in a larger convex S -free neighborhood.
- The family \mathcal{F}_∂ of convex S -free neighborhoods of x_0 that are also gradient polyhedra for some finite set of points in \mathbb{R}^n .
- The family $\mathcal{F}_{\partial,S} \subseteq \mathcal{F}_\partial$ of convex S -free neighborhoods of x_0 that are also gradient polyhedra for some finite set of points in S .

We propose the above families so as to leverage a recent surge of activity analyzing their structure; the surveys [3] and Chapter 6 of [6] provide good overviews and references for this whole line of work. This well-developed theory provides powerful mathematical tools to work with these families. As an example, this prior work shows that for most sets S that occur in practice (which includes the integer and mixed-integer cases), the family \mathcal{F}_{\max} only contains polyhedra. This is good from two perspectives:

- polyhedra are easier to represent and compute with than general S -free neighborhoods,
- the inner optimization problem (2) of computing $L(C)$ becomes the problem of solving finitely many continuous convex optimization problems, corresponding to the facets of C .

Of course, the first question to settle is whether these three families actually enjoy strong duality, i.e., does strong duality hold with the \mathcal{F}_{\max} -dual, \mathcal{F}_∂ -dual and $\mathcal{F}_{\partial,S}$ -dual? It turns out that the main result in [2] shows that for the mixed-integer case, i.e., $S = C \cap (\mathbb{Z}^{n_1} \times \mathbb{R}^{n_2})$ for some convex set C , the \mathcal{F}_∂ -dual enjoys strong duality under conditions of the Slater type from continuous optimization. It is not hard to strengthen their result to also show that strong duality holds with the $\mathcal{F}_{\max} \cap \mathcal{F}_\partial$ -dual, under some additional assumptions.

In this paper, we give conditions on S and f such that strong duality holds with the family $\mathcal{F}_{\max} \cap \mathcal{F}_{\partial,S}$. Below we give an explanation as to why this family is very desirable. If these conditions on S and f are met, our result is stronger than Baes et al. [2]. For

example, when S is the set of integer points in a compact convex set and f is any convex function, our certificate is a stronger one. However, our conditions on S and f do not cover certain mixed-integer problems; whereas, the certificate from Baes et al. still exists in these settings. Having said that, it is not immediately clear to us whether strong certificates like ours exist for all mixed-integer problems.

3. Strong optimality certificates

Definition 4. A strong optimality certificate of size k for (1) is a set of points $z_1, \dots, z_k \in S$ together with subgradients $a_i \in \partial f(z_i)$ such that

$$Q := \{x \in \mathbb{R}^n : \langle a_i, x - z_i \rangle \leq 0, \quad i = 1, \dots, k\} \text{ is } S\text{-free}, \quad (4)$$

$$\langle a_i, z_j - z_i \rangle < 0 \text{ for all } i \neq j. \quad (5)$$

Remark 5. If a strong optimality certificate exists, then the infimum of f over S is attained and we have $\min_{s \in S} f(s) = \min_{i \in [k]} f(z_i)$. In other words, given a strong optimality certificate, we can compute (1) by simply evaluating $f(z_1), \dots, f(z_k)$.

Indeed, recall that $a \in \partial f(z)$ means that $f(x) \geq f(z) + \langle a, x - z \rangle$ holds for all $x \in \mathbb{R}^n$.

Since Q is S -free, for every $s \in S$ there is some $i \in [k]$ such that $\langle a_i, s - z_i \rangle \geq 0$ and hence $f(s) \geq f(z_i)$.

In order to verify that z_1, \dots, z_k together with a_1, \dots, a_k form a strong optimality certificate, one has to check whether the polyhedron Q is S -free. Deciding whether a general polyhedron is S -free might be a difficult task. However, Property (5) ensures that Q is maximal S -free, i.e., Q is not properly contained in any other S -free closed convex set: Indeed, Property (5) implies that Q is a full-dimensional polyhedron and that $\{x \in Q : \langle a_i, x - z_i \rangle = 0\}$ is a facet of Q containing $z_i \in S$ in its relative interior for every $i \in [k]$. Since every closed convex set C that properly contains Q contains the relative interior of at least one facet of Q in its interior, C cannot be S -free.

For particular sets S , the properties of S -free sets that are maximal have been extensively studied and are much better understood than general S -free sets. For instance, if $S = (\mathbb{R}^d \times \mathbb{Z}^n) \cap C$ where C is a closed convex subset of \mathbb{R}^{n+d} , maximal S -free sets are polyhedra with at most 2^n facets [11]. In particular, if $S = \mathbb{Z}^2$ the characterizations in [8,10] yield a very simple algorithm to detect whether a polyhedron is maximal \mathbb{Z}^2 -free.

In order to state our main result, we need the notion of the Helly number $h(S)$ of the set S , which is the largest number m such that there exist convex sets $C_1, \dots, C_m \subseteq \mathbb{R}^n$ satisfying

$$\bigcap_{i \in [m]} C_i \cap S = \emptyset \quad \text{and} \quad \bigcap_{i \in [m] \setminus \{j\}} C_i \cap S \neq \emptyset \text{ for every } j \in [m]. \quad (6)$$

For an introduction to Helly numbers we refer to [9]. We are now ready to state the main theorem of this paper.

Theorem 6. Let $S \subseteq \mathbb{R}^n$ and $f : \mathbb{R}^n \rightarrow \mathbb{R}$ be a convex function such that

- (i) $0 \notin \partial f(s)$ for all $s \in S$,
- (ii) $h(S)$ is finite, and
- (iii) for every polyhedron $P \subseteq \mathbb{R}^n$ with $\text{int}(P) \cap S \neq \emptyset$ there exists an $s^* \in \text{int}(P) \cap S$ with $f(s^*) = \inf_{s \in \text{int}(P) \cap S} f(s)$.

Then there exists a strong optimality certificate of size at most $h(S)$.

Let us comment on the assumptions in Theorem 6. First, if $0 \in \partial f(s^*)$ for some $s^* \in S$, then s^* is an optimal solution to (1) as well as $s^* \in \arg \inf_{x \in \mathbb{R}^n} f(x)$. An easy certificate of optimality in this case is the subgradient 0 .

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