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# On deep holes of Gabidulin codes

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### A R T I C L E I N F O

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#### ABSTRACT

In this paper, we study deep holes of Gabidulin codes in both rank and Hamming metrics. Specifically, first, we give a tight lower bound for the distance of any word to a Gabidulin code and a sufficient and necessary condition for achieving this lower bound as well. Then, a class of deep holes of a Gabidulin code are discovered. Furthermore, we obtain some other deep holes for certain Gabidulin codes.

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## 1. Introduction

Let  $\mathbb{F}_{q^m}^n$  be an *n*-dimensional vector space over a finite field  $\mathbb{F}_{q^m}$  where *q* is a prime power, and *n*, *m* are positive integers. In this paper we only consider the case when  $n \leq m$ . Let  $\beta = (\beta_1, \ldots, \beta_m)$  be a basis of  $\mathbb{F}_{q^m}$  over  $\mathbb{F}_q$ . Let  $\mathcal{F}_i$  be the map from  $\mathbb{F}_{q^m}$  to  $\mathbb{F}_q$  where  $\mathcal{F}_i(u)$  is the *i*-th coordinate of an element  $u \in \mathbb{F}_{q^m}$  in the basis representation with  $\beta$ . To any  $\mathbf{u} = (u_1, \ldots, u_n)$  in  $\mathbb{F}_{q^m}^n$ , we may associate the matrix  $\bar{\mathbf{u}} = (\bar{u}_{i,j})_{1 \leq i \leq m, 1 \leq j \leq n} \in \mathcal{M}_{m,n}(\mathbb{F}_q)$  in which  $\bar{u}_{i,j} = \mathcal{F}_i(u_j)$ . The rank weight of the vector  $\mathbf{u}$  can be defined by the rank of the associated matrix  $\bar{\mathbf{u}}$ , denoted by  $w_R(\mathbf{u})$ . Thus, we can define the rank distance between two vectors  $\mathbf{u}$  and  $\mathbf{v}$  in  $\mathbb{F}_{q^m}^n$  as  $d_R(\mathbf{u}, \mathbf{v}) = w_R(\mathbf{u} - \mathbf{v})$ . We refer to [18] for more details on codes for the rank distance.

For integers  $1 \leq k \leq n$ , a linear rank-metric code C of length n and dimension k over  $\mathbb{F}_{q^m}$  is a subspace of dimension k of  $\mathbb{F}_{q^m}^n$  embedded with the rank metric. The minimum rank distance of the code C, denoted by  $d_R(C)$ , is the minimum rank weight of the non-zero codewords in C. A linear rank-metric code C of length n and dimension k over  $\mathbb{F}_{q^m}$  is called a maximum rank distance (MRD) code if  $d_R(C) = n - k + 1$ . A  $k \times n$  matrix is called a generator matrix of C if its rows span the code.

The rank distance of any word  $\mathbf{u} \in \mathbb{F}_{q^m}^n$  to C is defined as

$$d_R(\mathbf{u}, C) = \min\{d_R(\mathbf{u}, \mathbf{c}) \mid \mathbf{c} \in C\}.$$

It plays an important role in decoding of rank-metric codes. The maximum rank distance

$$\rho_R(C) = \max\{d_R(\mathbf{u}, C) \mid \mathbf{u} \in \mathbb{F}_{q^m}^n\}$$

is called the covering radius of C. If the rank distance from a word to the code C achieves the covering radius of the code, the word is called a deep hole of the code C.

The covering radius and deep holes of a linear code embedded with Hamming metric were studied extensively [1-5,10,12,14,16,22-27], in which MDS codes such as generalized Reed–Solomon codes, standard Reed–Solomon codes and projective Reed–Solomon codes were explored deeply. Gabidulin codes were introduced by Gabidulin in [7] and independently by Delsarte in [6]. Gabidulin codes can be seen as the q-analog of Reed– Solomon codes. Furthermore, Gabidulin codes are MRD codes. Over the last decade there has been increased interest in Gabidulin codes, mainly because of their relevance to network coding [15,19]. The covering radius for a Gabidulin code was also studied in [8,9,20]. However, little is known about deep holes for such a code. In this paper, we give a tight lower bound for the distance of any word to a Gabidulin code in both rank and Hamming metrics, and a sufficient and necessary condition for attaining this lower bound as well. Then, a class of deep holes of a Gabidulin code are discovered. Furthermore, we study the distance of a special class of words to a Gabidulin code and so obtain some other deep holes for certain Gabidulin codes. Note that we refer to rank metric if Hamming metric is not explicitly pointed out in this paper. Download English Version:

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