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# Partitioning the Cartesian product of a tree and a cycle<sup>☆</sup>

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## ABSTRACT

Let G = (V, E) be a graph of order n, and  $\lambda = (\lambda_1, \lambda_2, ..., \lambda_p)$  a sequence of positive integers. The sequence  $\lambda$  is admissible for G if  $\lambda_1 + \cdots + \lambda_p = n$ . Such an admissible sequence  $\lambda$  is said to be realizable in G if there exists a partition  $(V_1, V_2, ..., V_p)$  of the vertex set V such that  $V_i$  induces a connected subgraph of order  $n_i$  in G for each i. If every admissible sequence is realizable in G, then we say that G is arbitrarily partitionable (AP, for short). We show that if a tree T of maximum degree at most n + 1 has a path containing all the vertices of degree n + 1, then  $T \square C_n$  has a Hamiltonian path. In particular, for any caterpillar T with maximum degree at most n + 1,  $T \square C_n$  is AP. In addition, if T is a caterpillar with  $\Delta(T) \ge n + 4$ , then  $T \square C_n$  is not AP. For the cases  $n + 2 \le \Delta(T) \le n + 3$ , we present some sufficient conditions for a caterpillar T such that  $T \square C_n$  is AP.

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#### 1. Introduction

Let G = (V, E) be a simple graph of order n, and  $\lambda = (\lambda_1, \lambda_2, ..., \lambda_p)$  a sequence of positive integers. The sequence  $\lambda$  is *admissible* for G if  $\lambda_1 + \cdots + \lambda_p = n$ . Such an admissible sequence  $\lambda$  is said to be *realizable* in G if there exists a partition  $(V_1, V_2, ..., V_p)$  of the vertex set V such that each  $V_i$  induces a connected subgraph of order  $n_i$  in G. If every admissible sequence is realizable in G, then we say that G is *arbitrarily partitionable* (AP, for short). For a given integer  $k \in \{1, 2, ..., n\}$ , a graph G is called k-AP if every admissible sequence  $\lambda$  of G of length at most k is realizable in G. The notion of AP graphs was first introduced by Barth et al. [1], and independently, by Horňák and Woźniak [17]. Observe that having a (almost) perfect matching is necessary for a graph to be AP. So, the independence number of an n-vertex AP graph is at most  $\lceil \frac{n}{2} \rceil$ . Questions related to arbitrarily partitions of graphs are old. In 1978 and 1977 respectively, Győri [14] and Lovász [23] independently proved that k-connected graphs are k-AP.

It is clear that a connected graph is AP if one of its spanning trees is AP. The problem of deciding whether a given admissible sequence is realizable in a given graph *G* is NP-complete [1]. We refer to Refs. [2,11,12] for algorithmic aspect of AP graphs. Various properties on AP trees have been shown. It is not known whether it is NP-complete for the class of all trees. In [1], Barth et al. showed that this problem is polynomial for the class of tripodes. In [17], Horňák and Woźniak showed that the maximum degree of a AP tree is at most 6. Later, this bound has been decreased to 4 in [2]. Cichacz et al. [13] gave a complete characterization of AP caterpillars with four leaves. They also described two infinite families of AP trees with maximum degree three or four. Ravaux [26] focused on trees with a large diameter. There are also some results on AP star-like trees, unicyclic AP graphs and the shape of AP trees [3,18,21].

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**Fig. 1.** A Hamiltonian path  $P_0$  of  $T_0 \Box C_n$ .

A graph is said to be *traceable* if it contains a Hamiltonian path. It is clear that each traceable graph is AP. Therefore, a condition for a graph being traceable implies that the graph is AP. So, there are some relevant works [15,20,24,25] trying to replace some known conditions for traceability by the weaker ones implying that a graph is AP.

There are other families of AP graphs, such as OL-AP, R-AP and AP+k. For results on these aspect, we refer to [5–7,9,16,19,22]. Some structure of AP graphs and minimal AP graphs are investigated in [8,10].

Given two graphs *G* and *H*, the *Cartesian product* of *G* and *H*, denoted by  $G \Box H$ , is the graph whose vertex set is  $V(G) \times V(H)$  and two vertices  $(v_1, u_1)$  and  $(v_2, u_2)$  are adjacent in  $G \Box H$  if and only if either  $v_1 = v_2$  and  $u_1 u_2 \in E(H)$  or  $u_1 = u_2$  and  $v_1 v_2 \in E(G)$ . In [6], Baudon et al. conjectured that the Cartesian product of an AP graph and a traceable graph is AP and proved that the conjecture holds if the order of the traceable graph is at most four.

For a vertex v of a graph G, let d(v) denote the degree of v, and  $\Delta(G)$  is the maximum degree of G. In this paper, we focus on the Cartesian product of a tree and a cycle. In Section 2, we show that if a tree T with  $\Delta(T) \le n + 1$  has a path containing all the vertices of degree n + 1, then  $T \square C_n$  is traceable. In particular,  $T \square C_n$  is AP for any caterpillar T with  $\Delta(T) \le n + 1$ . In Section 3, we show that for a caterpillar T, the Cartesian product  $T \square C_n$  is not AP if  $\Delta(T) \ge n + 4$ . For the cases  $n + 2 \le \Delta(T) \le n + 3$ , we present some sufficient conditions for a caterpillar T such that  $T \square C_n$  is AP.

#### 2. Cartesian product of a tree and a cycle

Let *G* be a graph. For a vertex v, let N(v) denote the set of neighbors of v in *G*, and for a set  $S \subseteq V(G)$ , let  $N(S) = \bigcup_{v \in S} N(v)$ . The distance of vertices u and v in *G*, denoted by d(u, v), is the length of a shortest path joining u and v in *G*. In the sequel, we denote the vertices of the cycle  $C_n$  in the clockwise direction by  $u_1, u_2, \ldots, u_n$  and the vertex  $(v, u_j)$  of  $T \square C_n$  by  $v^j$ . In 1982, Batagelj and Pisanski proved the following elegant theorem.

**Theorem 2.1.** (Batagelj, Pisanski [4]) For an integer  $n \ge 3$  and a tree T, the Cartesian product  $T \square C_n$  is Hamiltonian if and only if  $\Delta(T) \le n$ .

**Corollary 2.2.** Let *G* be a connected graph, and *H* be a Hamiltonian graph of order *n*. If *G* has a spanning tree of maximum degree at most *n*, then  $G \square H$  is AP.

**Theorem 2.3.** Let *T* be a tree with  $\Delta(T) \le n + 1$  for an integer  $n \ge 3$ . If a path of *T* contains all the vertices of degree n + 1, then  $T \square C_n$  is traceable.

**Proof.** In view of Theorem 2.1, we assume that  $\Delta(T) = n + 1$ . Let  $v_0v_1, \ldots, v_s$  be a maximal path in T which contains all the vertices of maximum degree. Clearly,  $d(v_0) = d(v_s) = 1$ ,  $d(v_i) \le n$  for  $v_i \in V(T) \setminus V_0$ , where  $V_0 = \{v_0, v_1, \ldots, v_s\}$ . Let  $t = \max\{d_T(v, u) : v \in V_0, u \in V(T) \setminus V_0\}$ . Let  $V_1 = N(V_0) \setminus V_0$  and  $V_{i+1} = N(V_i) \setminus V_{i-1}$  for any  $i \in \{1, \ldots, t-1\}$ . Since T is a tree and  $G[V_0]$  is connected,  $V_i$  is an independent set for each  $i \ge 1$ . Let  $T_i = T[\cup_{j=0}^i V_j]$  for  $i \in \{0, 1, \ldots, t\}$ . Clearly,  $T_t \cong T$ , and  $T_i \square C_n$  is a subgraph of  $T_i \square C_n$  for  $i \le j$ . We construct a Hamiltonian path of  $T \square C_n$  recursively.

**Step 1:** Take a Hamiltonian path  $P_0$  (in bold lines) of  $T_0 \Box C_n$  as shown in Fig. 1.

$$P_0 = \begin{cases} \frac{v_0^1 v_0^2 \dots v_0^n}{v_0^1 v_0^2 \dots v_0^n} \frac{v_1^n v_1^{n-1} \dots v_1^1}{v_1^n v_1^{n-1} \dots v_1^1} \frac{v_2^1 v_2^2 \dots v_2^n v_3^n \dots v_{s-1}^1}{v_2^1 v_2^2 \dots v_s^n v_3^n \dots v_{s-1}^n} \frac{v_s^1 v_s^2 \dots v_s^n}{v_s^n v_s^{n-1} \dots v_s^1}, & \text{if } s \text{ is odd.} \end{cases}$$

**Step 2:** Extend the Hamiltonian path  $P_0$  of  $T_0 \Box C_n$  to a Hamiltonian path  $P_1$  of  $T_1 \Box C_n$ .

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