# Maximum matchings in regular graphs 

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#### Abstract

It was conjectured by Mkrtchyan, Petrosyan and Vardanyan that every graph $G$ with $\Delta(G)-\delta(G) \leq 1$ has a maximum matching $M$ such that any two $M$-unsaturated vertices do not share a neighbor. The results obtained in Mkrtchyan et al. (2010), Petrosyan (2014) and Picouleau (2010) leave the conjecture unknown only for $k$-regular graphs with $4 \leq k \leq 6$. All counterexamples for $k$-regular graphs ( $k \geq 7$ ) given in Petrosyan (2014) have multiple edges. In this paper, we confirm the conjecture for all $k$-regular simple graphs and also $k$-regular multigraphs with $k \leq 4$.


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## 1. Introduction

Graphs considered in this paper may have multi-edges, but no loops. A graph without multi-edges is called a simple graph. A matching $M$ of a graph $G$ is a set of independent edges. A vertex is $M$-saturated if it is incident with an edge of $M$, and $M$-unsaturated otherwise. A matching $M$ is said to be maximum if for any other matching $M^{\prime},|M| \geq\left|M^{\prime}\right|$. A matching $M$ is perfect if it covers all vertices of $G$. If $G$ has a perfect matching, the every maximum matching is a perfect matching. The maximum and minimum degrees of a graph $G$ are denoted by $\Delta(G)$ and $\delta(G)$, respectively. Mkrtchyan, Petrosyan and Vardanyan $[4,5]$ made the following conjecture.

Conjecture 1.1 (Mkrtchyan et al. [4,5]). Let $G$ be a graph with $\Delta(G)-\delta(G) \leq 1$. Then $G$ contains a maximum matching $M$ such that any two $M$-unsaturated vertices do not share a neighbor.

This conjecture is verified for subcubic graphs (i.e. $\Delta(G)=3$ ) with multi-edges by Mkrtchyan, Petrosyan and Vardanyan [5]. Later, Picouleau [7] find a counterexample to the conjecture, which is a bipartite simple graph with $\delta(G)=4$ and $\Delta(G)=5$. Petrosyan [6] constructs counterexamples to the conjecture for all $k$-regular graphs with $k \geq 7$ and for graphs $G$ with $\Delta(G)-\delta(G)=1$ and $\Delta(G) \geq 4$. Note that, most of counterexamples of Conjecture 1.1 for graphs $G$ with $\Delta(G)-$ $\delta(G)=1$ are simple, but all $k$-regular graphs with $k \geq 7$ given by Petrosyan [6] have multi-edges. As affirmative answer to Conjecture 1.1 is known only for graphs with $\Delta(G) \leq 3$, Mkrtchyan et al. [5] asked whether the conjecture holds for any $k$-regular graphs with $k \geq 4$.

In this note, we consider the conjecture for both $k$-regular simple graphs and $k$-regular graphs with multi-edges. First we show that Conjecture 1.1 does hold for all $k$-regular simple graphs.

Theorem 1.2. Let $G$ be a k-regular simple graph. Then $G$ has a maximum matching $M$ such that any two $M$-unsaturated vertices do not share a neighbor.

Further, we show that Conjecture 1.1 holds for $k$-regular graphs with multi-edges for $k \leq 4$.

[^0]Theorem 1.3. Let $G$ be a $k$-regular graph with $k \leq 4$. Then $G$ has a maximum matching $M$ such that any two $M$-unsaturated vertices do not share a neighbor.

Our results together with examples given by Petrosyan [6] leave Conjecture 1.1 unknown for 5 and 6-regular graphs with multi-edges.

## 2. Preliminaries

Let $G$ be a graph and $v$ be a vertex of $G$. The neighborhood of $v$ is set of all vertices adjacent to $v$, denoted by $N(v)$. The degree of $v$, denoted by $d_{G}(v)$ (or $d(v)$ if there is no confusion), is the number of edges incident to $v$. For $X \subseteq V(G)$, let $\delta(X):=\min \{d(v) \mid v \in X\}$ and $\Delta(X):=\max \{d(v) \mid v \in X\}$. The neighborhood of $X$ is defined as $N(X):=\{y \mid y$ is a neighbor of a vertex $x \in X\}$. For two subsets $X_{1}$ and $X_{2}$ of $V(G)$, use $\left[X_{1}, X_{2}\right]$ to denote the all edges with one endvertex in $X_{1}$ and another endvertex in $X_{2}$. For two subgraphs $G_{1}$ and $G_{2}$ of $G$, the symmetric difference of $G_{1} \oplus G_{2}$ is defined as a subgraph with vertex set $V\left(G_{1}\right) \cup V\left(G_{2}\right)$ and edge set $\left(E\left(G_{1}\right) \cup E\left(G_{2}\right)\right) \backslash\left(E\left(G_{1}\right) \cap E\left(G_{2}\right)\right)$.

A matching of a graph $G$ is a near-perfect matching if it covers all vertices except one. If a graph $G$ has a near perfect matching, then $G$ has odd number of vertices. A graph is factor-critical if, for any vertex $v$, the subgraph $G \backslash\{v\}$ has a perfect matching. Every maximum matching of a factor-critical graph is a near-perfect matching.

Let $D$ be the set of all vertices of a graph $G$ which are not covered by at least one maximum matching, and $A$, the set of all vertices in $V(G)-D$ adjacent to at least one vertex in $D$. Denote $C=V(G)-A-D$. The graph induced by all vertices in $D$ (resp. $A$ and $C$ ) is denoted by $G[D]$ (resp. $G[A]$ and $G[C]$ ). The following theorem characterizes the structures of maximum matchings of graphs, which is due to Gallai [2] and Edmonds [1].

Theorem 2.1 (Gallai-Edmonds Structure Theorem, Theorem 3.2.1 in [3]). Let $G$ be a graph, and $A, D$ and $C$ are defined as above. Then:
(1) the components of the subgraph induced by D are factor-critical;
(2) the subgraph induced by C has a perfect matching;
(3) if $M$ is a maximum matching of $G$, it contains a near-perfect matching of each component of $G[D]$, a perfect matching of $G[C]$ and matches all vertices of $A$ with vertices in distinct components of $G[D]$.

Let $M$ be a maximum matching of a graph $G$. By Gallai-Edmonds Structure Theorem, $M$ does not contain edges from $G[A]$ and all $M$-unsaturated vertices of $G$ belong to $D$. Contract every component of $G[D]$ to a vertex and let $B$ be the set of all these vertices. Let $G(A, B)$ be the bipartite graph with bipartition $(A, B)$ and all edges of $G$ in $[A, D]$. Because every component of $G[D]$ is factor-critical, a maximum matching of $G(A, B)$ is corresponding to a maximum matching of $G$, and vice versa. In fact, the proofs of our results mainly focus on maximum matchings of the bipartite graph $G(A, B)$.

Before processing to prove our main results, we need some results for maximum matchings of bipartite graphs $G(A, B)$.
Theorem 2.2 (Hall's Theorem, Theorem 1.13 in [3]). Let $G(A, B)$ be a bipartite graph. If $|N(S)| \geq|S|$ for any $S \subseteq A$, then $G$ has a matching $M$ covering all vertices of $A$.

The following technical lemma is needed in the proofs of our main results.
Lemma 2.3. Let $G(A, B)$ be a bipartite graph such that every maximum matching of $G(A, B)$ covers all vertices of $A$. Let $W \subseteq B$ such that $\delta(W) \geq \Delta(A) \neq 0$. Then $G(A, B)$ has a maximum matching $M$ covering all vertices of $W$.

Proof. Let $M$ be a maximum matching of $G(A, B)$ such that the number of vertices of $W$ covered by $M$ is maximum. If $M$ covers all vertices of $W$, the lemma follows. So assume that there exists an $M$-unsaturated vertex $x \in W$.

For any $U \subseteq W$, we have $\delta(U) \geq \delta(W)$ and $N(U) \subset A$. Further,

$$
\delta(W)|U| \leq \delta(U)|U| \leq|[U, N(U)]| \leq \sum_{v \in N(U)} d(v) \leq \Delta(A)|N(U)|
$$

It follows that $|N(U)| \geq|U|$ because $\delta(W) \geq \Delta(A) \neq 0$. By applying Hall's Theorem on the subgraph induced by $W$ and $N(W)$, it follows that $G$ has a matching $M^{\prime}$ covering all vertices of $W$.

Let $M \oplus M^{\prime}$ be the symmetric difference of $M$ and $M^{\prime}$. Every component of $M \oplus M^{\prime}$ is either a path or a cycle. Since $x$ is not covered by $M$ but is covered by $M^{\prime}$, it follows that $x$ is an end-vertex of some path-component $P$ of $M \oplus M^{\prime}$. Let $y$ be another end-vertex of $P$. Note that every vertex of $A$ is covered by an edge of $M$ and every vertex of $W$ is covered by an edge of $M^{\prime}$. So $y \in B \backslash W$.

Then let $M^{\prime \prime}=M \oplus P$. Then $M^{\prime \prime}$ is a maximum matching of $G$ which covers $x$ and all vertices covered by $M$ except $y$. Note that $y \in B \backslash W$ and $x \in W$. Hence $M^{\prime \prime}$ covers more vertices of $W$ than $M$, a contradiction to the maximality of the number of vertices of $W$ covered by $M$. This completes the proof.

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