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Maximum matchings in regular graphs

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ABSTRACT

It was conjectured by Mkrtchyan, Petrosyan and Vardanyan that every graph *G* with $\Delta(G) - \delta(G) \leq 1$ has a maximum matching *M* such that any two *M*-unsaturated vertices do not share a neighbor. The results obtained in Mkrtchyan et al. (2010), Petrosyan (2014) and Picouleau (2010) leave the conjecture unknown only for *k*-regular graphs with $4 \leq k \leq 6$. All counterexamples for *k*-regular graphs ($k \geq 7$) given in Petrosyan (2014) have multiple edges. In this paper, we confirm the conjecture for all *k*-regular simple graphs and also *k*-regular multigraphs with $k \leq 4$.

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1. Introduction

Graphs considered in this paper may have multi-edges, but no loops. A graph without multi-edges is called a *simple* graph. A *matching* M of a graph G is a set of independent edges. A vertex is M-saturated if it is incident with an edge of M, and M-unsaturated otherwise. A matching M is said to be maximum if for any other matching M', $|M| \ge |M'|$. A matching M is perfect if it covers all vertices of G. If G has a perfect matching, the every maximum matching is a perfect matching. The maximum and minimum degrees of a graph G are denoted by $\Delta(G)$ and $\delta(G)$, respectively. Mkrtchyan, Petrosyan and Vardanyan [4,5] made the following conjecture.

Conjecture 1.1 (*Mkrtchyan et al.* [4,5]). Let G be a graph with $\Delta(G) - \delta(G) \leq 1$. Then G contains a maximum matching M such that any two M-unsaturated vertices do not share a neighbor.

This conjecture is verified for subcubic graphs (i.e. $\Delta(G) = 3$) with multi-edges by Mkrtchyan, Petrosyan and Vardanyan [5]. Later, Picouleau [7] find a counterexample to the conjecture, which is a bipartite simple graph with $\delta(G) = 4$ and $\Delta(G) = 5$. Petrosyan [6] constructs counterexamples to the conjecture for all *k*-regular graphs with $k \ge 7$ and for graphs *G* with $\Delta(G) - \delta(G) = 1$ and $\Delta(G) \ge 4$. Note that, most of counterexamples of Conjecture 1.1 for graphs *G* with $\Delta(G) - \delta(G) = 1$ are simple, but all *k*-regular graphs with $k \ge 7$ given by Petrosyan [6] have multi-edges. As affirmative answer to Conjecture 1.1 is known only for graphs with $\Delta(G) \le 3$, Mkrtchyan et al. [5] asked whether the conjecture holds for any *k*-regular graphs with $k \ge 4$.

In this note, we consider the conjecture for both *k*-regular simple graphs and *k*-regular graphs with multi-edges. First we show that Conjecture 1.1 does hold for all *k*-regular simple graphs.

Theorem 1.2. Let *G* be a *k*-regular simple graph. Then *G* has a maximum matching *M* such that any two *M*-unsaturated vertices do not share a neighbor.

Further, we show that Conjecture 1.1 holds for *k*-regular graphs with multi-edges for $k \le 4$.

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Theorem 1.3. Let G be a k-regular graph with $k \le 4$. Then G has a maximum matching M such that any two M-unsaturated vertices do not share a neighbor.

Our results together with examples given by Petrosyan [6] leave Conjecture 1.1 unknown for 5 and 6-regular graphs with multi-edges.

2. Preliminaries

Let *G* be a graph and *v* be a vertex of *G*. The *neighborhood* of *v* is set of all vertices adjacent to *v*, denoted by N(v). The degree of *v*, denoted by $d_G(v)$ (or d(v) if there is no confusion), is the number of edges incident to *v*. For $X \subseteq V(G)$, let $\delta(X) := \min\{d(v)|v \in X\}$ and $\Delta(X) := \max\{d(v)|v \in X\}$. The neighborhood of *X* is defined as $N(X) := \{y|y \text{ is a neighbor of a vertex <math>x \in X\}$. For two subsets X_1 and X_2 of V(G), use $[X_1, X_2]$ to denote the all edges with one endvertex in X_1 and another endvertex in X_2 . For two subgraphs G_1 and G_2 of *G*, the symmetric difference of $G_1 \oplus G_2$ is defined as a subgraph with vertex set $V(G_1) \cup V(G_2)$ and edge set $(E(G_1) \cup E(G_2)) \setminus (E(G_1) \cap E(G_2))$.

A matching of a graph *G* is a *near-perfect matching* if it covers all vertices except one. If a graph *G* has a near perfect matching, then *G* has odd number of vertices. A graph is *factor-critical* if, for any vertex *v*, the subgraph $G \setminus \{v\}$ has a perfect matching. Every maximum matching of a factor-critical graph is a near-perfect matching.

Let *D* be the set of all vertices of a graph *G* which are not covered by at least one maximum matching, and *A*, the set of all vertices in V(G) - D adjacent to at least one vertex in *D*. Denote C = V(G) - A - D. The graph induced by all vertices in *D* (resp. *A* and *C*) is denoted by *G*[*D*] (resp. *G*[*A*] and *G*[*C*]). The following theorem characterizes the structures of maximum matchings of graphs, which is due to Gallai [2] and Edmonds [1].

Theorem 2.1 (*Gallai–Edmonds Structure Theorem, Theorem 3.2.1 in* [3]). Let *G* be a graph, and *A*, *D* and *C* are defined as above. Then:

(1) the components of the subgraph induced by D are factor-critical;

(2) the subgraph induced by C has a perfect matching;

(3) if *M* is a maximum matching of *G*, it contains a near-perfect matching of each component of G[D], a perfect matching of G[C] and matches all vertices of *A* with vertices in distinct components of G[D].

Let *M* be a maximum matching of a graph *G*. By Gallai–Edmonds Structure Theorem, *M* does not contain edges from *G*[*A*] and all *M*-unsaturated vertices of *G* belong to *D*. Contract every component of *G*[*D*] to a vertex and let *B* be the set of all these vertices. Let *G*(*A*, *B*) be the bipartite graph with bipartition (*A*, *B*) and all edges of *G* in [*A*, *D*]. Because every component of *G*[*D*] is factor-critical, a maximum matching of *G*(*A*, *B*) is corresponding to a maximum matching of *G*, and vice versa. In fact, the proofs of our results mainly focus on maximum matchings of the bipartite graph *G*(*A*, *B*).

Before processing to prove our main results, we need some results for maximum matchings of bipartite graphs G(A, B).

Theorem 2.2 (Hall's Theorem, Theorem 1.13 in [3]). Let G(A, B) be a bipartite graph. If $|N(S)| \ge |S|$ for any $S \subseteq A$, then G has a matching M covering all vertices of A.

The following technical lemma is needed in the proofs of our main results.

Lemma 2.3. Let G(A, B) be a bipartite graph such that every maximum matching of G(A, B) covers all vertices of A. Let $W \subseteq B$ such that $\delta(W) \ge \Delta(A) \neq 0$. Then G(A, B) has a maximum matching M covering all vertices of W.

Proof. Let *M* be a maximum matching of G(A, B) such that the number of vertices of *W* covered by *M* is maximum. If *M* covers all vertices of *W*, the lemma follows. So assume that there exists an *M*-unsaturated vertex $x \in W$.

For any $U \subseteq W$, we have $\delta(U) \geq \delta(W)$ and $N(U) \subset A$. Further,

$$\delta(W)|U| \leq \delta(U)|U| \leq |[U, N(U)]| \leq \sum_{v \in N(U)} d(v) \leq \Delta(A)|N(U)|.$$

It follows that $|N(U)| \ge |U|$ because $\delta(W) \ge \Delta(A) \ne 0$. By applying Hall's Theorem on the subgraph induced by W and N(W), it follows that G has a matching M' covering all vertices of W.

Let $M \oplus M'$ be the symmetric difference of M and M'. Every component of $M \oplus M'$ is either a path or a cycle. Since x is not covered by M but is covered by M', it follows that x is an end-vertex of some path-component P of $M \oplus M'$. Let y be another end-vertex of P. Note that every vertex of A is covered by an edge of M and every vertex of W is covered by an edge of M'. So $y \in B \setminus W$.

Then let $M'' = M \oplus P$. Then M'' is a maximum matching of *G* which covers *x* and all vertices covered by *M* except *y*. Note that $y \in B \setminus W$ and $x \in W$. Hence M'' covers more vertices of *W* than *M*, a contradiction to the maximality of the number of vertices of *W* covered by *M*. This completes the proof. \Box

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