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## Minimum density of identifying codes of king grids  $1$

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## **Abstract**

A set  $C \subseteq V(G)$  is an *identifying code* in a graph G if for all  $v \in V(G)$ ,  $C[v] \neq \emptyset$ , and for all distinct  $u, v \in V(G)$ ,  $C[u] \neq C[v]$ , where  $C[v] = N[v] \cap C$  and  $N[v]$  denotes the closed neighbourhood of  $v$  in  $G$ . The minimum density of an identifying code in G is denoted by  $d^*(G)$ . In this paper, we study the density of king grids which are strong product of two paths. We show that for every king grid  $G, d^*(G) \geq 2/9$ . In addition, we show this bound is attained only for king grids which are strong products of two infinite paths. Given  $k \geq 3$ , we denote by  $\mathcal{K}_k$  the (infinite) king strip with k rows. We prove that  $d^*(K_3)=1/3$ ,  $d^*(K_4)=5/16$ ,  $d^*(K_5)=4/15$  and  $d^*(\mathcal{K}_6) = 5/18$ . We also prove that  $\frac{2}{9} + \frac{8}{81k} \leq d^*(\mathcal{K}_k) \leq \frac{2}{9} + \frac{4}{9k}$  for every  $k \geq 7$ .

*Keywords:* Identifying code, King grid, Discharging Method.

## **1 Introduction**

Let G be a graph. The *neighbourhood* of a vertex v of G, denoted by  $N(v)$ , is the set of vertices adjacent to  $v$  in  $G$ , and the *closed neighbourhood* of  $v$  is the

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set  $N[v] = N(v) \cup \{v\}$ . Given a set  $C \subseteq V(G)$ , let  $C[v] = N[v] \cap C$ . We say that C is an *identifying code* of G if  $C[v] \neq \emptyset$  for all  $v \in V(G)$ , and  $C[u] \neq C[v]$ for all distinct  $u, v \in V(G)$ . Clearly, a graph G has an identifying code if and only if it contains no *twins* (vertices  $u, v \in V(G)$  with  $N[u] = N[v]$ ).

Let G be a (finite or infinite) graph with bounded maximum degree. For any non-negative integer r and vertex v, we denote by  $B_r(v)$  the ball of radius r in G centered at v, that is  $B_r(v) = \{x \mid dist(v, x) \leq r\}$ . For any set of vertices  $C \subseteq V(G)$ , the *density* of C in G, denoted by  $d(C, G)$ , is defined by

$$
d(C, G) = \limsup_{r \to +\infty} \frac{|C \cap B_r(v_0)|}{|B_r(v_0)|} ,
$$

where  $v_0$  is an arbitrary vertex in G. The infimum of the density of an identifying code in G is denoted by  $d^*(G)$ . Observe that if G is finite, then  $d^*(G) = |C^*|/|V(G)|$ , where  $C^*$  is a minimum-size identifying code of G.

The problem of finding low-density identifying codes was introduced in [\[13\]](#page--1-0) in relation to fault diagnosis in arrays of processors. Particular interest was dedicated to grids as many processor networks have a grid topology. Many results have been obtained on square grids  $[4,1,10,2,12]$ , triangular grids  $[13,11]$ , and hexagonal grids  $[6,8,9]$ . In this paper, we study king grids, which are strong products of two paths. The *strong product* of two graphs  $G$  and  $H$ , denoted by  $G \boxtimes H$ , is the graph with vertex set  $V(G) \times V(H)$  and edge set :

$$
E(G \boxtimes H) = \{(a, b)(a, b') \mid a \in V(G) \text{ and } bb' \in E(H)\}
$$
  

$$
\cup \{(a, b)(a', b) \mid aa' \in E(G) \text{ and } b \in V(H)\}
$$
  

$$
\cup \{(a, b)(a', b') \mid aa' \in E(G) \text{ and } bb' \in E(H)\}.
$$

The two-way infinite path, denoted by  $P_{\mathbb{Z}}$ , is the graph with vertex set  $\mathbb{Z}$ and edge set  $\{\{i, i+1\} \in \mathbb{Z}\}\$ , and the *one-way infinite path*, denoted by  $P_{\mathbb{N}}$ , is the graph with vertex set N and edge set  $\{\{i, i+1\} \mid i \in \mathbb{N}\}\$ . A path is a connected subgraph of  $P_{\mathbb{Z}}$ . For every positive integer k,  $P_k$  is the subgraph of  $P_{\mathbb{Z}}$  induced by  $\{1, 2, \ldots, k\}$ . A king grid is the strong product of two (finite or infinite) paths. The plane king grid is  $\mathcal{G}_K = P_{\mathbb{Z}} \boxtimes P_{\mathbb{Z}}$ , the half-plane king grid is  $\mathcal{H}_K = P_{\mathbb{Z}} \boxtimes P_{\mathbb{N}}$ , the quarter-plane king grid is  $\mathcal{Q}_K = P_{\mathbb{N}} \boxtimes P_{\mathbb{N}}$ , and the  $king \ strip \ of \ height \ k \ is \ \mathcal{K}_k = P_{\mathbb{Z}} \boxtimes P_k.$ 

In 2001, Cohen et al. [\[7\]](#page--1-0) proved that  $d^*(\mathcal{G}_K) \geq 2/9$ . In 2002, Charon et al. [\[3\]](#page--1-0) obtained an optimal identifying code with density 2/9. They provided the tile depicted in Fig. [1,](#page--1-0) which generates a periodic tiling of the plane with periods (0, 6) and (6, 0), yielding an identifying code  $C_{\infty}$  of the bidimensional infinite king grid with density  $\frac{2}{9}$ .

In this paper, using the Discharging Method (see Section 3 of [\[11\]](#page--1-0) for a

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